

COS 301: Programming Languages

Syntax Analysis

Parsing

- Syntactic analysis = *parsing*
- Goal of parser:
 - Find all syntax errors \Rightarrow diagnostic message
 - If no syntax errors \Rightarrow parse tree

Types of parsers

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Computational complexity

- Brute-force approach:
 - Try every possible rule (exhaustive search)
 - *Exponential* in length of the program
 - So don't do that!
- Better:
 - Several algorithms $\Rightarrow \mathcal{O}(n^3)$
 - Still too expensive for commercial compilers
- Can reduce the generality of the language to be parsed \Rightarrow linear $\mathcal{O}(n)$

Top-Down Parsing

Top-down parsing

- Given $xA\alpha$, where
 - x is a string of terminal symbols
 - A is the leftmost nonterminal symbol
 - α is a string of terminals and non-terminals
- Goal: find next sentential form in leftmost derivation
- Thus: choose rule such that A is the LHS; e.g.,
 - rules: $A \Rightarrow bB$, $A \Rightarrow cBb$, $A \Rightarrow a$
 - then we choose between:
 - $xA\alpha \Rightarrow xbB\alpha$, $xA\alpha \Rightarrow xcBb\alpha$, and $xA\alpha \Rightarrow xa\alpha$

Top-down parsing

- Which to choose for $xA\alpha$?
 - choices: $A \Rightarrow xbB\alpha$, $A \Rightarrow xcBb\alpha$, and $A \Rightarrow xa\alpha$
 - Look at the input
 - If the next token after x is a , b , or c , then it's clear
- What if some rule's RHS has nonterminals first?
 - E.g., $A \rightarrow Bb$, choice includes $xBb\alpha$
 - Which to choose?
- Much harder in this case – depends on what B expands to, etc.

Most common top-down parsers

- Two most common: *recursive-descent* and *table-driven*
- Work on a *subset* of CF grammars
- Algorithms are called **LL**
 - First L = left-to-right scan of input
 - Second L = leftmost derivation

Recursive-descent parser

- Simple algorithm to understand
- **Mutually-recursive set of procedures**
 - One for each (EBNF) production of language
 - I.e., one for each nonterminal
- Top-down parsing
- To begin, call procedure for start nonterminal
- Call procedure for leftmost nonterminal, etc.
- Procedure tries \Rightarrow parse tree rooted at itself, matches input
- Use scanner (lexer, tokenizer) \Rightarrow next token as needed

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→ Call $\langle \text{factor} \rangle$

→ Call $\langle \text{id} \rangle \Rightarrow B$

← Return $\langle \text{id} \rangle \rightarrow B$

← Return $\langle \text{factor} \rangle \rightarrow \langle \text{id} \rangle \rightarrow B$

Next token is *, so...

→ Call $\langle \text{factor} \rangle$

→ Call $\langle \text{id} \rangle \Rightarrow C,$

← Return $\langle \text{id} \rangle \rightarrow C$

← Return $\langle \text{factor} \rangle \rightarrow \langle \text{id} \rangle \rightarrow C$

RD parser: example

Given: **A+B*C**

- Grammar:

<expr> → **<term>** { (+ | -) **<term>** }

<term> → **<factor>** { (* | /) **<factor>** }

<factor> → **<id>** | **int_const** | (**<expr>**)

→ Call **<expr>**

→ Call **<term>**

→ Call **<factor>**

→ Call **<id>**: next token is an id (A), so success; next token not * or /

← Return **<id>** → A

← Return **<factor>** → **<id>** → A

Next token is +, so...

← Return **<term>** → **<factor>** → **<id>** → A

Next token is +, so...

→ Call **<term>**

→ Call **<factor>**

→ Call **<id>** ⇒ B

← Return **<id>** → B

← Return **<factor>** → **<id>** → B

Next token is *, so...

→ Call **<factor>**

→ Call **<id>** ⇒ C,

← Return **<id>** → C

← Return **<factor>** → **<id>** → C

← Return **<term>** → (**<factor>** → **<id>** → B) * (**<factor>** → **<id>** → C)

RD parser: example

Given: **A+B*C**

- Grammar:

<expr> → <term> { (+ | -) <term> }

<term> → <factor> { (* | /) <factor> }

<factor> → <id> | int_const | (<expr>)

→ Call <expr>

→ Call <term>

→ Call <factor>

→ Call <id>: next token is an id (A), so success; next token not * or /

← Return <id> → A

← Return <factor> → <id> → A

Next token is +, so...

← Return <term> → <factor> → <id> → A

Next token is +, so...

→ Call <term>

→ Call <factor>

→ Call <id> ⇒ B

← Return <id> → B

← Return <factor> → <id> → B

Next token is *, so...

→ Call <factor>

→ Call <id> ⇒ C,

← Return <id> → C

← Return <factor> → <id> → C

← Return <term> → (<factor> → <id> → B) * (<factor> → <id> → C)

← Return <expr> → (<term> → <factor> → <id> → A) + (<factor> → <id> → B) * <factor> → <id> → C)

RD parser: <expr>

```
void expr() {
    /* parses <expr> => <term> { (+|-) <term> } */
    printf("enter <expr>\n");
    term();
    while (nextToken == ADD_OP ||
           nextToken == SUB_OP) {
        lex();
        term();
    }
    printf("exit <expr>\n");
}
```

/* Q: Where does nextToken come from?

A: each function leaves the next unconsumed token in nextToken
each function assumes on entry that it is available in
nextToken */

RD parser: <term>

```
void term() {
    /* parses <term> => <factor> { (*|/) <factor> } */

    printf("enter <term>\n");
    factor();
    while (nextToken == MULT_OP ||
           nextToken == DIV_OP) {
        lex();
        factor();
    }
    printf("exit <term>\n");
}
```

RD parser: $\langle \text{factor} \rangle$

- Has to choose among several alternatives:
 $\langle \text{factor} \rangle \rightarrow \langle \text{id} \rangle \mid \langle \text{int_const} \rangle \mid (\langle \text{expr} \rangle)$
- May be able to detect a syntax error here:
 - E.g., $A + B * / C$
 - Fail in $\langle \text{factor} \rangle \Rightarrow$ “/” isn’t any allowed symbol

RD parser: <factor>

```
void factor() {
/* parses <factor> => <id> | int_constant | ( expr ) */

printf("enter <factor>\n");
if (nextToken == IDENT || nextToken == INT_LIT)
    lex();
else {
    if (nextToken == LEFT_PAREN) {
        lex();
        expr(); /* recursion! */
        if (nextToken == RIGHT_PAREN)
            lex();
        else
            error();
    } else
        error();
}
printf("exit <factor>\n");
}
```

RD parser: Example

- Given (sum + 47)/total

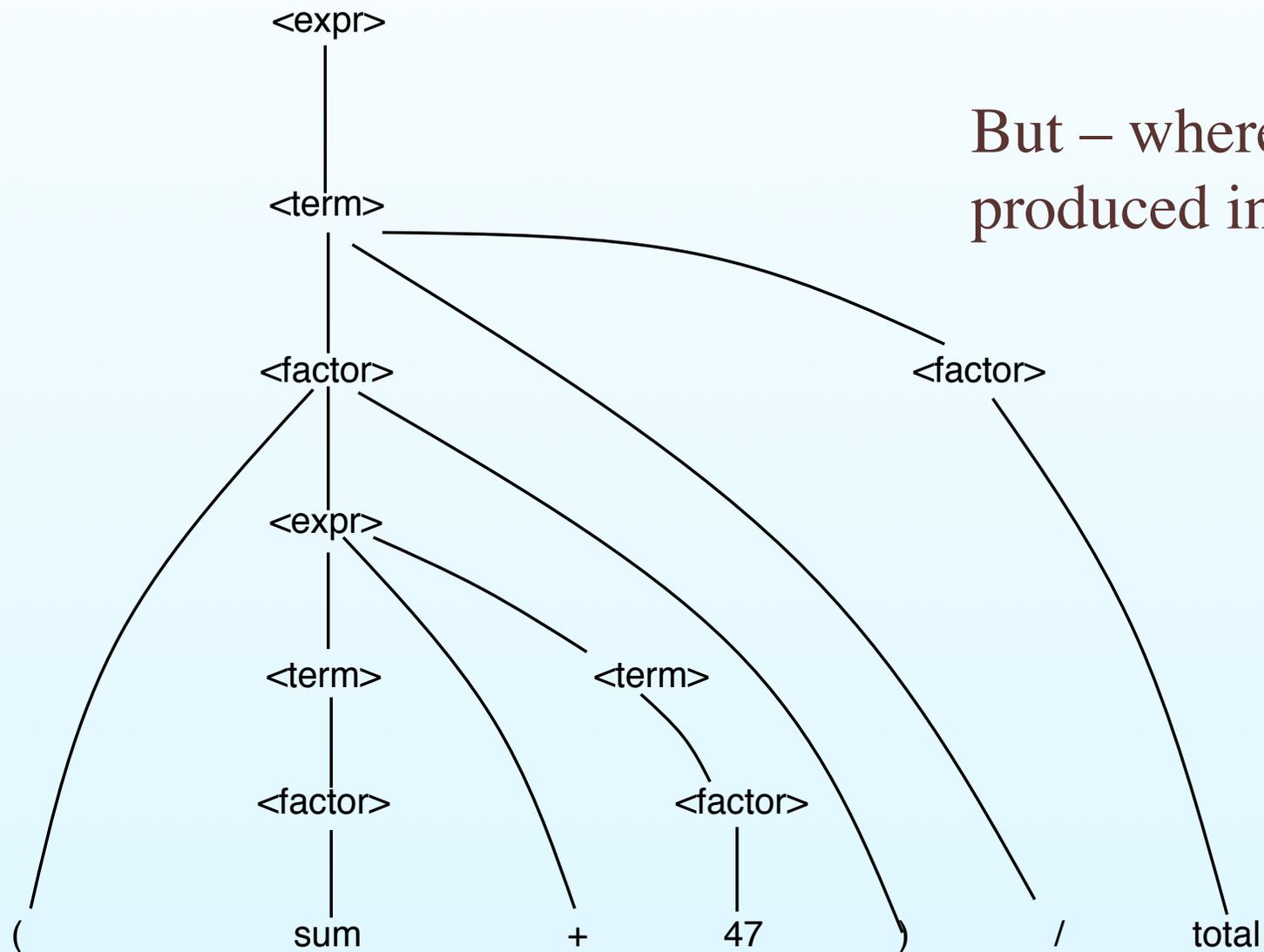
```
Next token is: 25 lexeme is (  
Enter <expr>  
Enter <term>  
Enter <factor>  
Next token is: 11 lexeme is sum  
Enter <expr>  
Enter <term>  
Enter <factor>  
Next token is: 21 lexeme is +  
Exit <factor>  
Exit <term>  
Next token is: 10 lexeme is 47  
Enter <term>  
Enter <factor>
```

RD parser: Example

```
Next token is: 26 lexeme is )  
Exit <factor>  
Exit <term>  
Exit <expr>  
Next token is: 24 lexeme is /  
Exit <factor>  
Next token is: 11 lexeme is total  
Enter <factor>  
Next token is: -1 lexeme is EOF  
Exit <factor>  
Exit <term>  
Exit <expr>
```

Parse tree

- For (sum + 47)/total



But – where is this produced in the functions?

Another example: if statement

```
<ifStmt> → if ( <boolExpr> ) <stmt> [else <stmt> ]
```

- Program for `<ifStmt>` nonterminal:
 - Check if current token is `if`
 - Call `lex()`, check if current token is `(`
 - Call `lex()`, then call `<boolExpr>`
 - Check if current token is `)`
 - Call `lex()` and call `<stmt>`
 - Check if current token = `else`
 - if so, call `lex()`, then call `<stmt>`

<ifStmt>

```
void ifstmt(){
    if (nextToken != IF_CODE)
        error();
    else {
        lex();
        if (nextToken != LEFT_PAREN)
            error();
        else {
            lex(); /* error in text; this was omitted */
            boolexpr();
            if (nextToken != RIGHT_PAREN)
                error();
            else {
                lex(); /* error in text; this was omitted */
                stmt();

                if (nextToken == ELSE_CODE){
                    lex();
                    stmt();
                } /* end if (nextToken == ELSE_CODE) */
            } /* end if (nextToken != RIGHT_PAREN) */
        } /* end if (nextToken != LEFT_PAREN) */
    } /* end if (nextToken != IF_CODE) */
} /* end IF_STMT */
```

LL grammars

- Recursive-descent parsers: parse left-to-right, leftmost derivations (LL)

- **Classes of grammars limited**

- Can't parse grammars with rules of form:

$$A \rightarrow A + B$$

- Direct left-recursion \Rightarrow infinite recursion!

- Can also have indirect recursion:

$$A \rightarrow Ba$$

$$B \rightarrow Ab$$

- Also, need to be able to choose correct RHS based on next input token – not always possible

Eliminating left-recursion

- For every nonterminal A :

- Group the rules for A as:

$$A \rightarrow A\alpha_1 \mid A\alpha_2 \mid \dots \mid A\alpha_m \mid \beta_1 \mid \beta_2 \mid \dots \mid \beta_n$$

- Then replace these rules with:

$$A \rightarrow \beta_1A' \mid \beta_2A' \mid \dots \mid \beta_nA'$$

$$A' \rightarrow \alpha_1A' \mid \alpha_2A' \mid \dots \mid \alpha_mA' \mid \varepsilon$$

- Left-recursive grammar \Rightarrow right-recursive grammar

Example

Example

- Grammar:

$$\begin{aligned} \langle e \rangle &\rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle \\ \langle t \rangle &\rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle \\ \langle f \rangle &\rightarrow (\langle e \rangle) \mid \langle \text{id} \rangle \end{aligned}$$

Example

- Grammar:

$$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$$

$$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$$

$$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle \text{id} \rangle$$

- For $\langle e \rangle$, $\alpha_1 = + \langle t \rangle$, $\beta_1 = \langle t \rangle$, so:

$$\langle e \rangle \rightarrow \langle t \rangle \langle e' \rangle$$

$$\langle e' \rangle \rightarrow + \langle t \rangle \langle e' \rangle \mid \varepsilon$$

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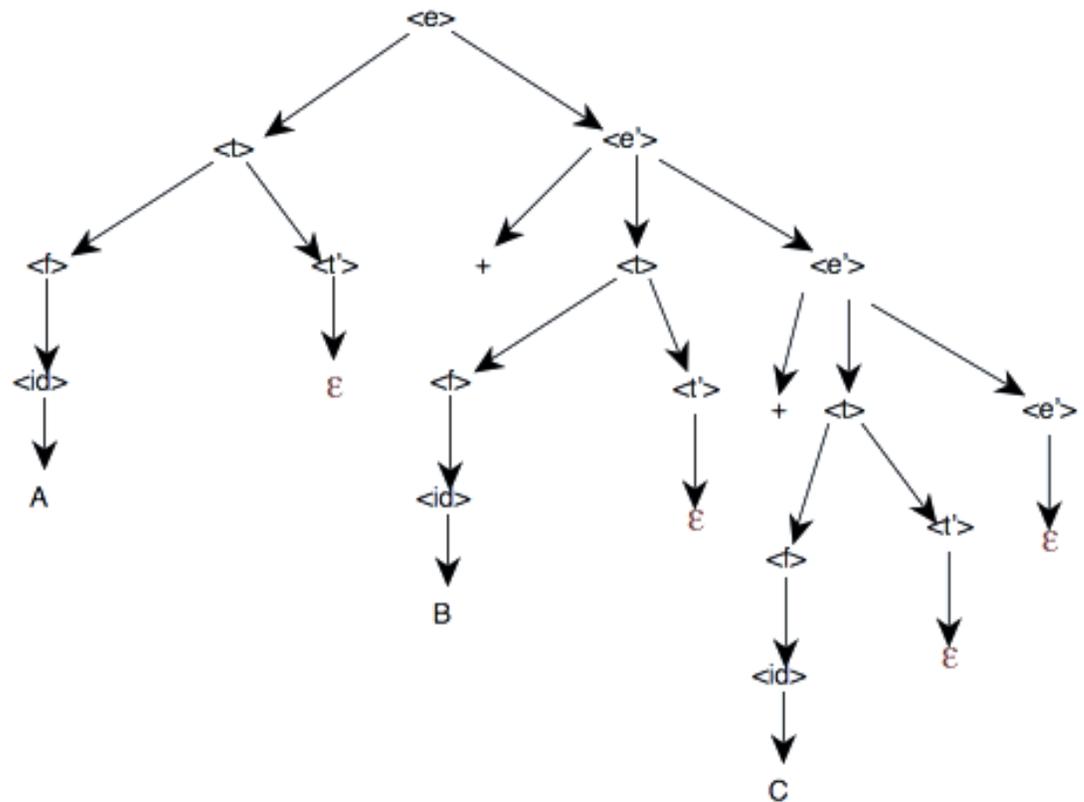
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Parse: $A + B * C$



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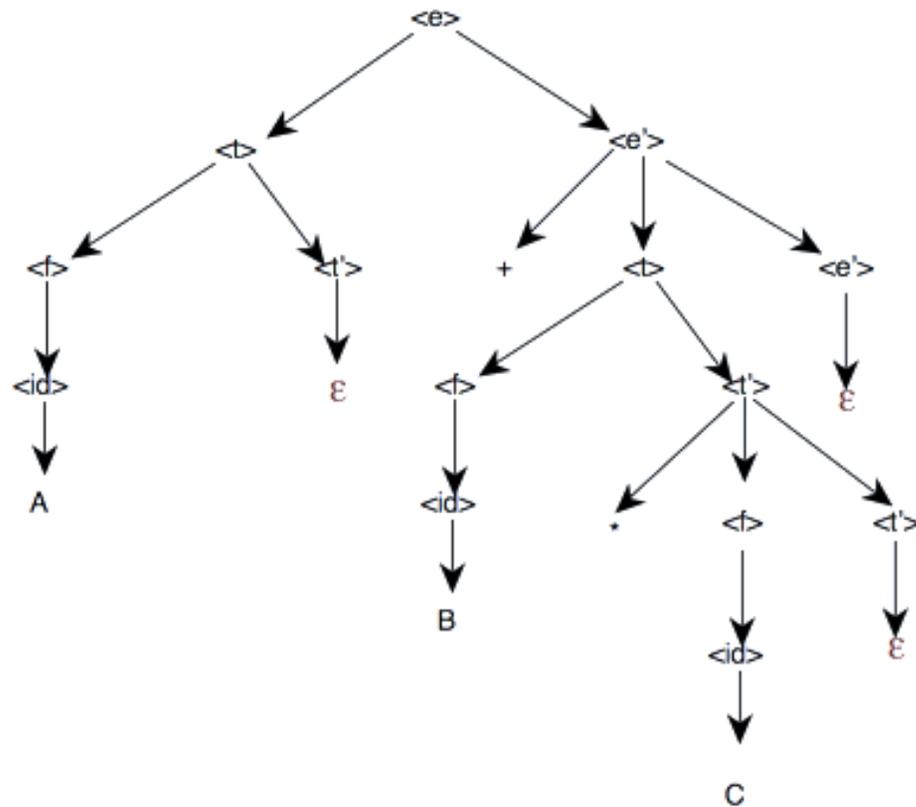
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Indirect recursion

- Indirect recursion also a problem for recursive-descent parsers
 - E.g.,

$$A \rightarrow B c$$
$$B \rightarrow A d$$

- Algorithm exists to get rid of it, too
- May be able to write grammar to avoid both kinds of recursion

Choosing RHS

- Lookahead
- Best case:
 - single token, no backtracking
 - called a **LL(1)** grammar (and algorithm)
- **Pairwise disjointness test**: checks if grammar is LL(1)

Pairwise disjointedness test

- General idea:
 - If non-terminal has > 1 RHS and...
 - ...the different RHSs ultimately produce same terminal...
 - ...then:
 - grammar is not LL(1) and...
 - ...can't be parsed w/ recursive-descent

Pairwise disjointedness test

- Compute set $\text{First}(r)$, where r is a rule RHS

$$\text{First}(r) = \{t \mid r \Longrightarrow^* t \beta\}$$

where t is a terminal, β is rest of sentential form

- If non-terminal N has two rules, i and j
 - Compute $\text{First}(i)$, $\text{First}(j)$
 - If $\text{First}(i) \cap \text{First}(j) \neq \{\}$ \Rightarrow grammar is not LL(1)

Left factoring

- If not pairwise disjoint, sometimes can fix grammar

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$$\text{Ex: } \langle e \rangle \rightarrow \langle t \rangle \mid \langle t \rangle + \langle e \rangle \mid \langle t \rangle - \langle e \rangle$$

$$\Rightarrow \langle e \rangle \rightarrow \langle t \rangle \langle e' \rangle$$

$$\langle e' \rangle \rightarrow + \langle e \rangle \mid - \langle e \rangle$$

- Or use EBNF: $A \rightarrow a [B]$

Bottom-Up Parsing

Bottom-up parsing

- Called **LR algorithms**
 - L = left-to-right scan
 - R = rightmost derivation
- Creates a parse tree starting with the leaves (the input)
- Creates tree in reverse order of rightmost derivation
- Sentential forms for this are called *right sentential forms*

Basic process

- Start with the input sentence (a right-sentential form)
- Find a RHS in the sentential form
- Replace it with the LHS
- Repeat until only the start symbol left.
- Record of the *reductions* \Rightarrow the parse tree

Example

Example

- Grammar:

Example

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$$\mathbf{E} \rightarrow \mathbf{E} + \mathbf{T} \mid \mathbf{T}$$

Example

- Grammar:

$$\begin{aligned} E &\rightarrow E + T \mid T \\ T &\rightarrow T * F \mid F \end{aligned}$$

Example

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$$E \rightarrow E + T \mid T$$
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- Parse: $A + B + C$
 - $A + B + C$ — assume lex can give $\text{id} + \text{id} + \text{id}$

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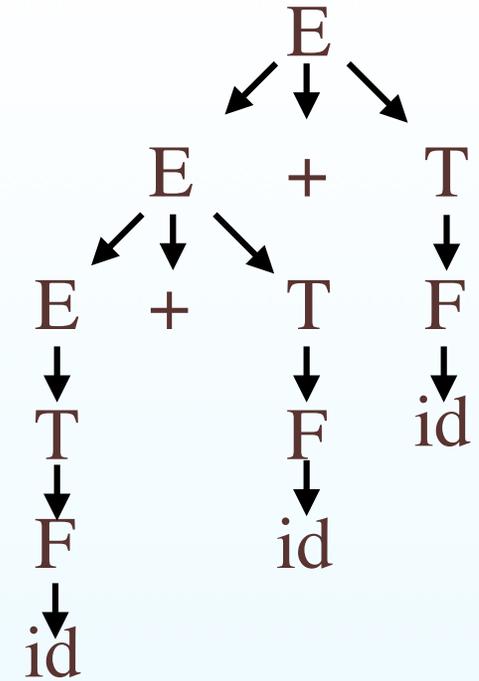
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- $\mathbf{F} + \mathbf{id} + \mathbf{id}$
- $\mathbf{T} + \mathbf{id} + \mathbf{id}$
- $\mathbf{E} + \mathbf{id} + \mathbf{id}$
- $\mathbf{E} + \mathbf{F} + \mathbf{id}$
- $\mathbf{E} + \mathbf{T} + \mathbf{id}$
- $\mathbf{E} + \mathbf{id}$
- $\mathbf{E} + \mathbf{F}$
- $\mathbf{E} + \mathbf{T}$
- \mathbf{E}

Example

- Grammar:

$$\begin{aligned} E &\rightarrow E + T \mid T \\ T &\rightarrow T * F \mid F \\ F &\rightarrow \text{id} \mid (E) \end{aligned}$$


- Parse: $A + B + C$

- $A + B + C$ — assume lex can give $\text{id} + \text{id} + \text{id}$
- $\text{id} + \text{id} + \text{id}$
- $F + \text{id} + \text{id}$
- $T + \text{id} + \text{id}$
- $E + \text{id} + \text{id}$
- $E + F + \text{id}$
- $E + T + \text{id}$
- $E + \text{id}$
- $E + F$
- $E + T$
- E

Problem: Which RHS to choose?

$$E \rightarrow E + T \mid T$$

$$T \rightarrow T * F \mid F$$

$$F \rightarrow \text{id} \mid (E)$$

- If current sentential form is: $E + T + \text{id}$
- Choose:
 - $E \rightarrow T$?
 - $E \rightarrow E + T$?
 - $T \rightarrow F$?

Handles

- Problem: Which RHS to choose, if several?
- Find *handle* in current sentential form
 - Handle = *correct* pattern in RHS to replace s.t.:
 - \Rightarrow new sentential form
 - & sentential form in right-derivation of current
- Only 1 handle/sentential form
- So...how to find handle?

Finding the handle

- Right derivation of $A + B * (C + D)$ given grammar:

$$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$$
$$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$$
$$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$$

Finding the handle

Parse: A + B * (C + D)

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: A + B * (C + D)

<e>

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: A + B * (C + D)

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$

$\langle e \rangle + \langle t \rangle * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$

$\langle e \rangle + \langle t \rangle * (C + D)$

$\langle e \rangle + \langle f \rangle * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$

$\langle e \rangle + \langle t \rangle * (C + D)$

$\langle e \rangle + \langle f \rangle * (C + D)$

$\langle e \rangle + \langle id \rangle * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B^* (C + D)$

$\langle e \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle + \langle t \rangle * \langle f \rangle$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$

$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$

$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$

$\langle e \rangle + \langle t \rangle * (C + D)$

$\langle e \rangle + \langle f \rangle * (C + D)$

$\langle e \rangle + \langle id \rangle * (C + D)$

$\langle e \rangle + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B^* (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B^* (C + D)$
 $\langle t \rangle + B^* (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B^* (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B^* (C + D)$
 $\langle t \rangle + B^* (C + D)$
 $\langle f \rangle + B^* (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B^* (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B^* (C + D)$
 $\langle t \rangle + B^* (C + D)$
 $\langle f \rangle + B^* (C + D)$
 $\langle id \rangle + B^* (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B^* (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B^* (C + D)$
 $\langle t \rangle + B^* (C + D)$
 $\langle f \rangle + B^* (C + D)$
 $\langle id \rangle + B^* (C + D)$
 $A + B^* (C + D)$

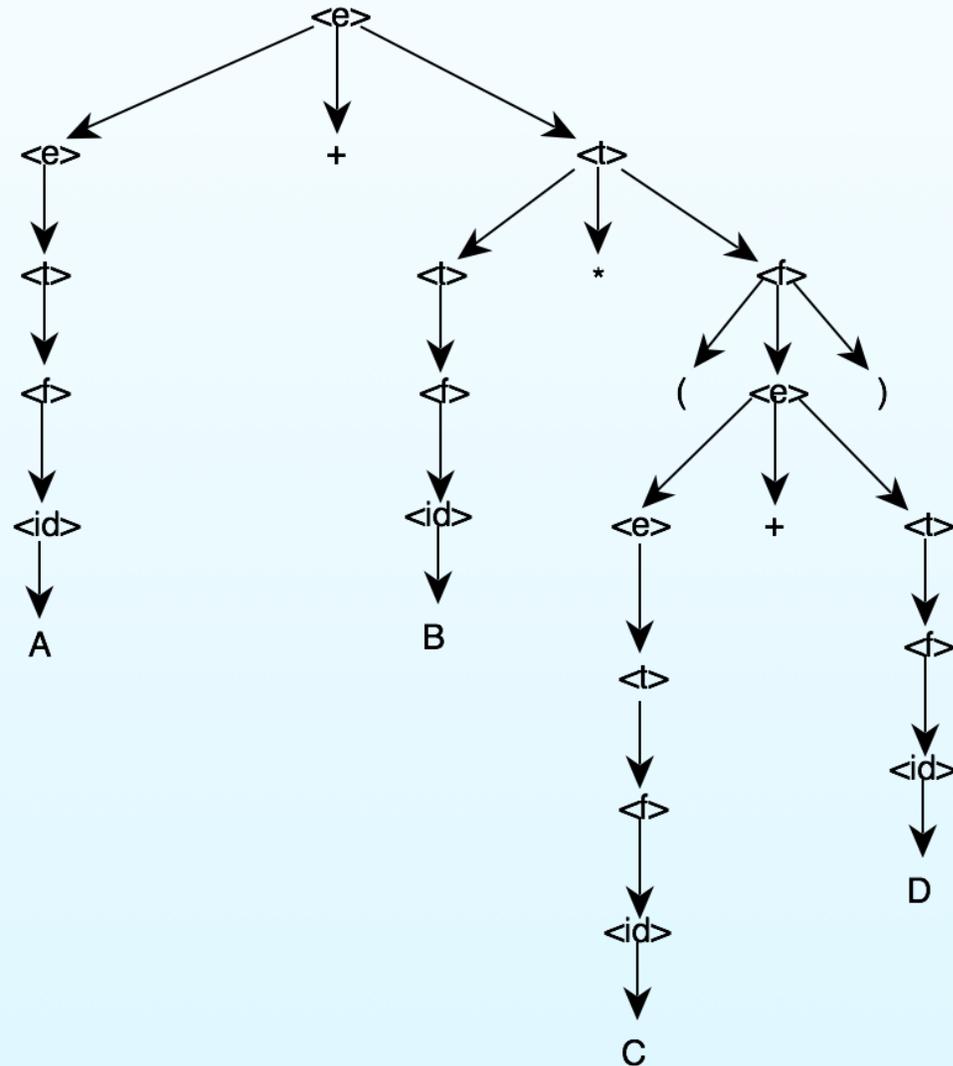
$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B * (C + D)$
 $\langle t \rangle + B * (C + D)$
 $\langle f \rangle + B * (C + D)$
 $\langle id \rangle + B * (C + D)$
 $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

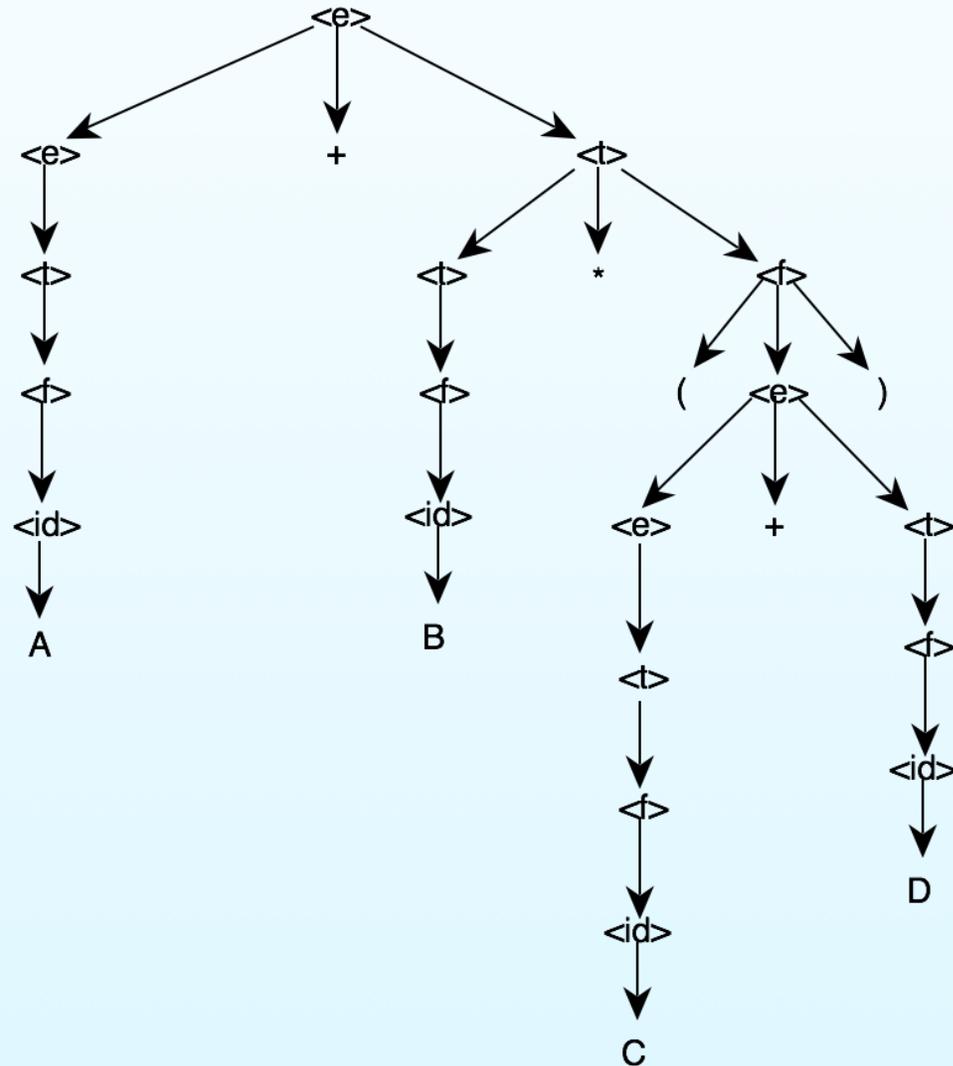


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B * (C + D)$
 $\langle t \rangle + B * (C + D)$
 $\langle f \rangle + B * (C + D)$
 $\langle id \rangle + B * (C + D)$
 $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

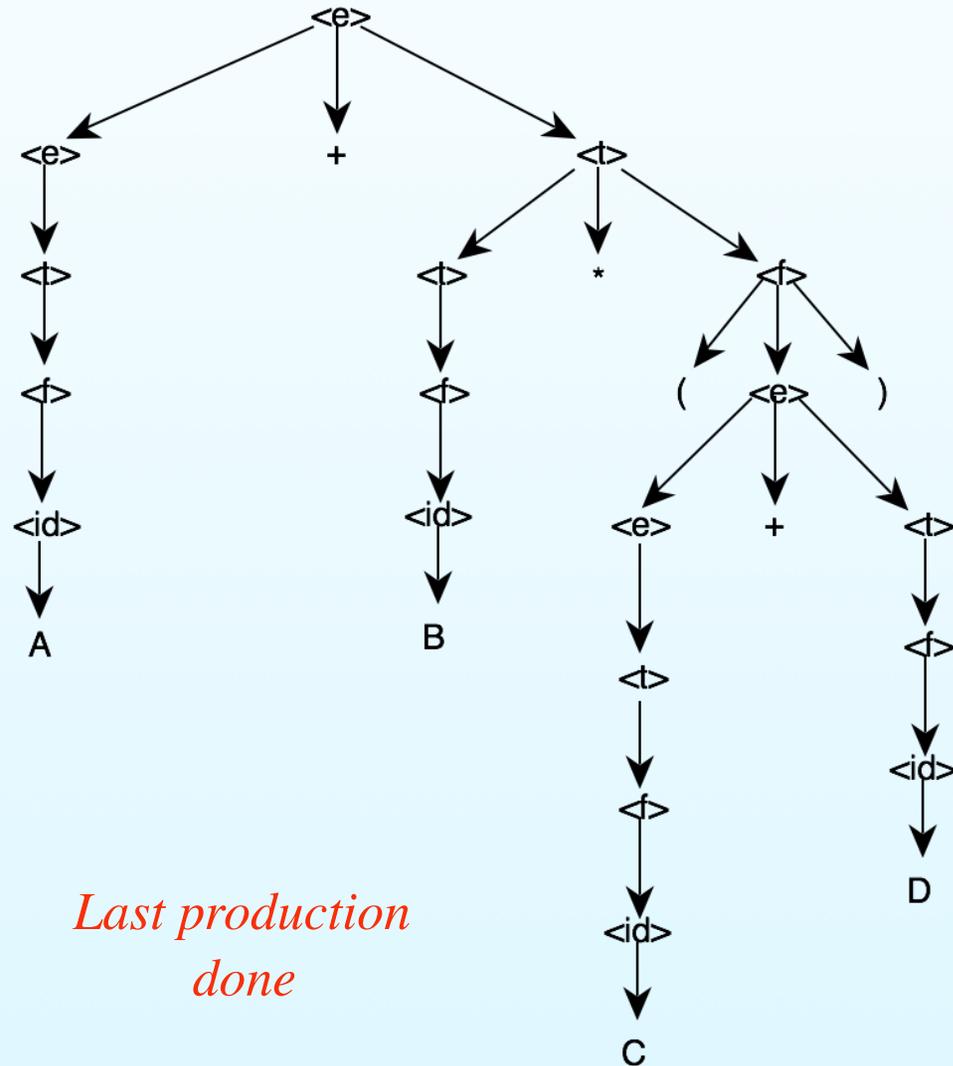


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B * (C + D)$
 $\langle t \rangle + B * (C + D)$
 $\langle f \rangle + B * (C + D)$
 $\langle id \rangle + B * (C + D)$
 $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

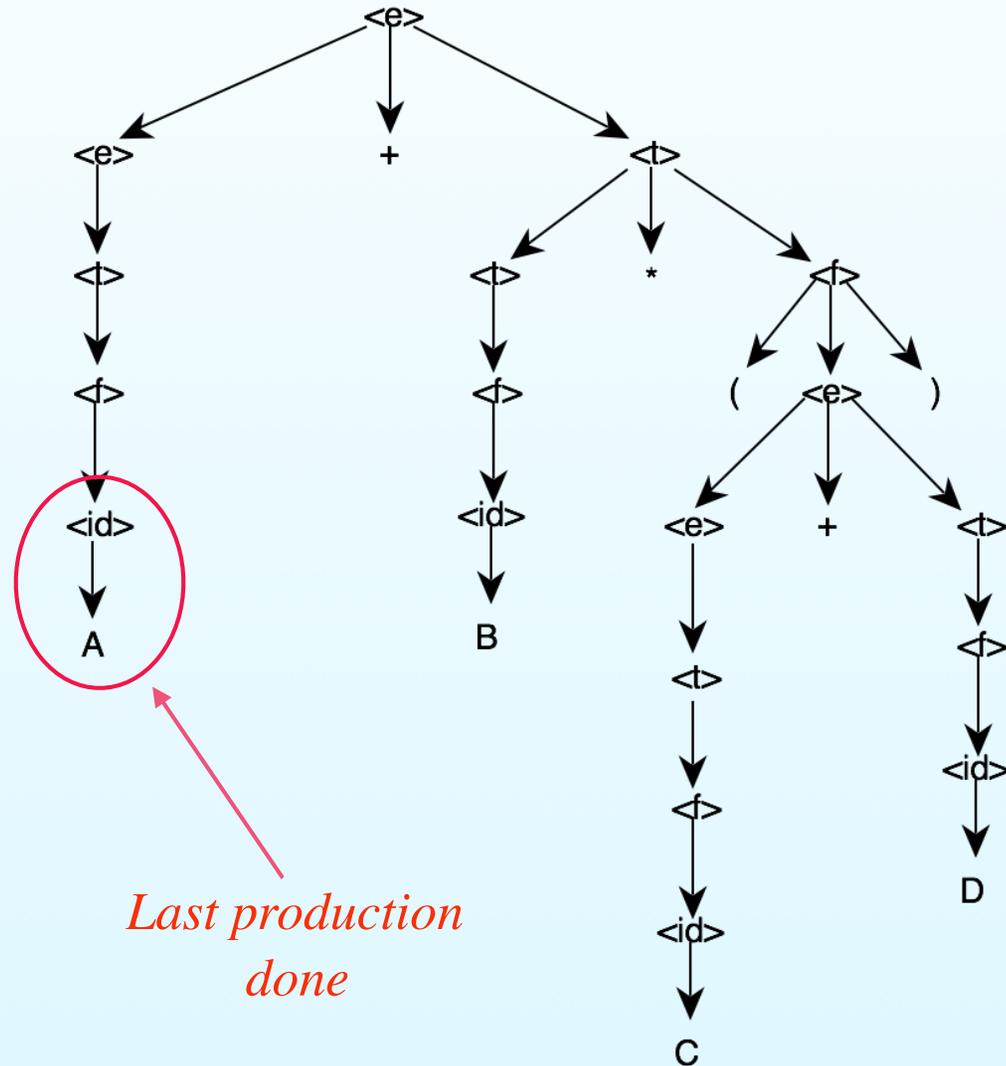


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B * (C + D)$
 $\langle t \rangle + B * (C + D)$
 $\langle f \rangle + B * (C + D)$
 $\langle id \rangle + B * (C + D)$
 $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



Last production done

Finding the handle

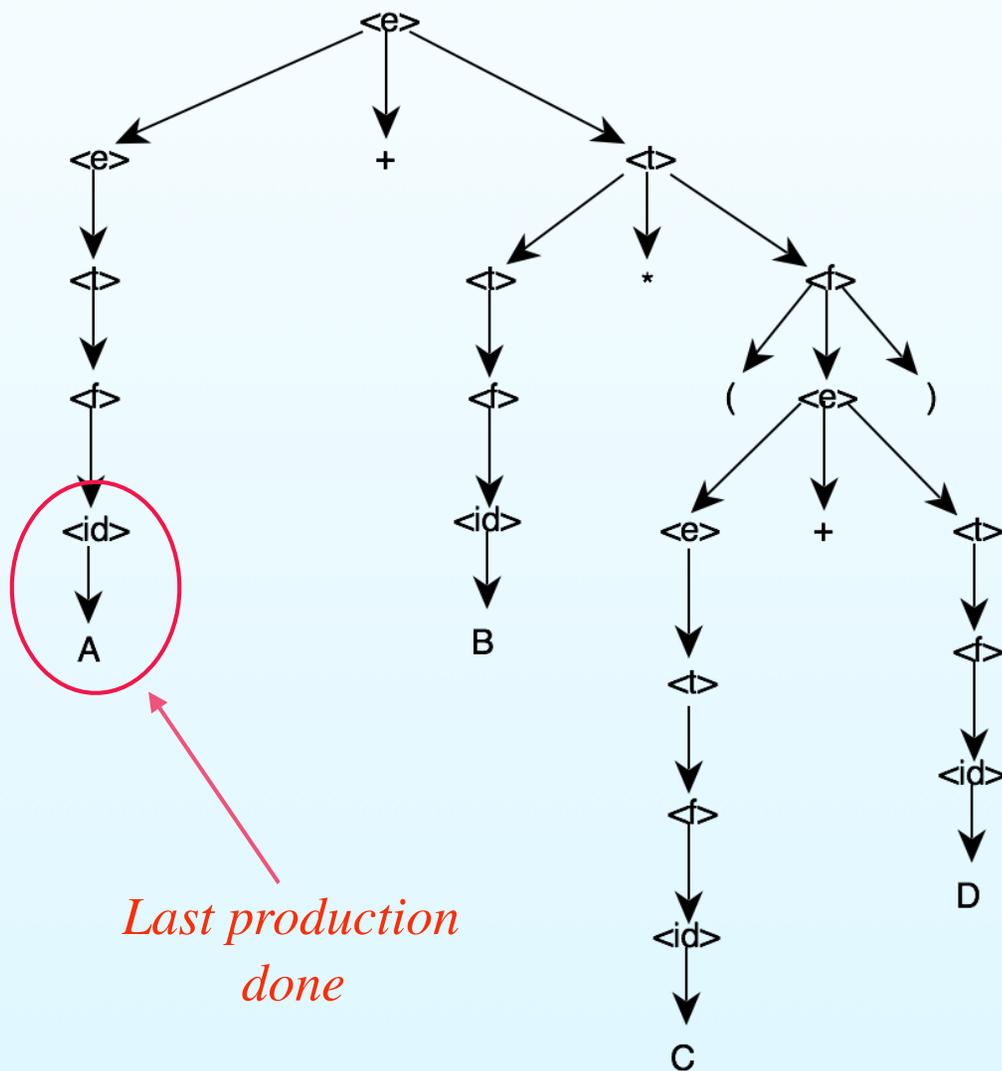
Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- $A + B * (C + D)$

Handle

Last production done

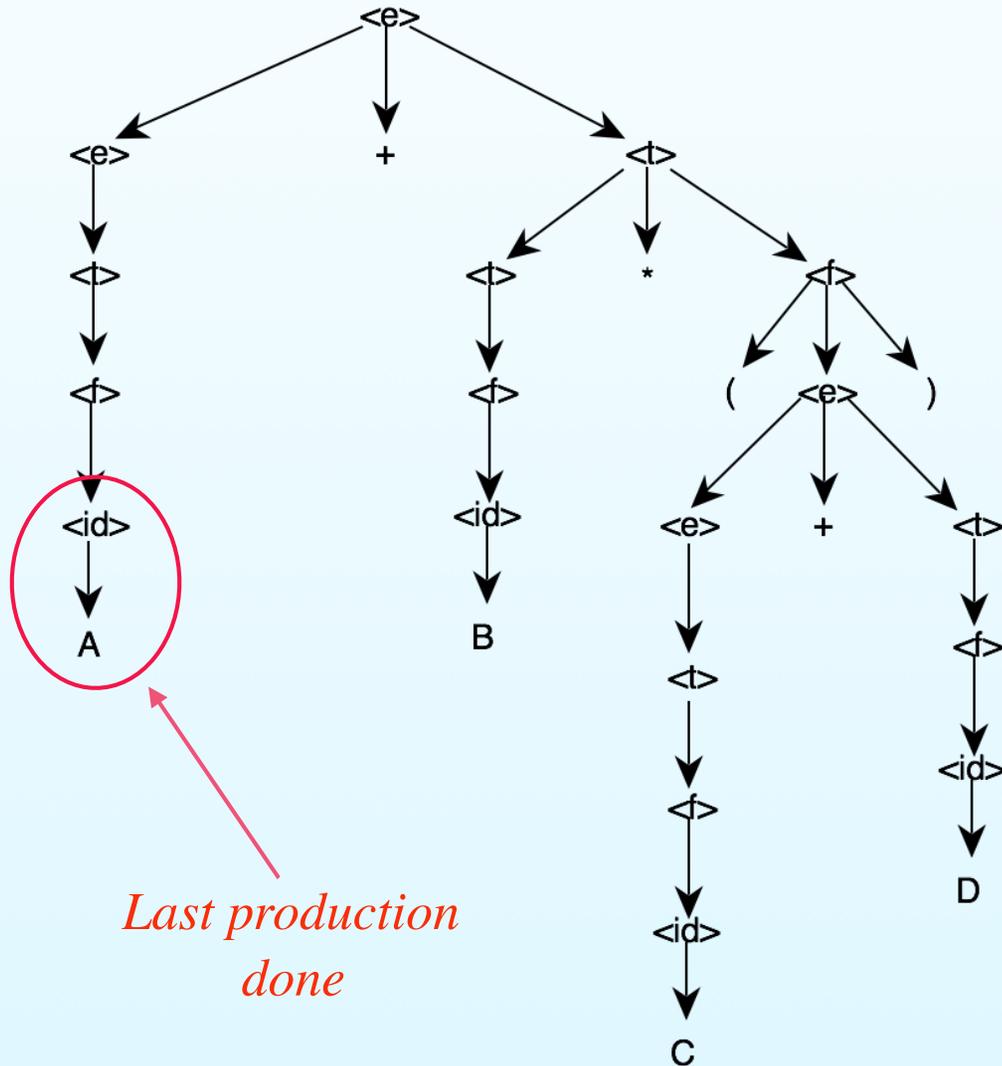


Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



Handle

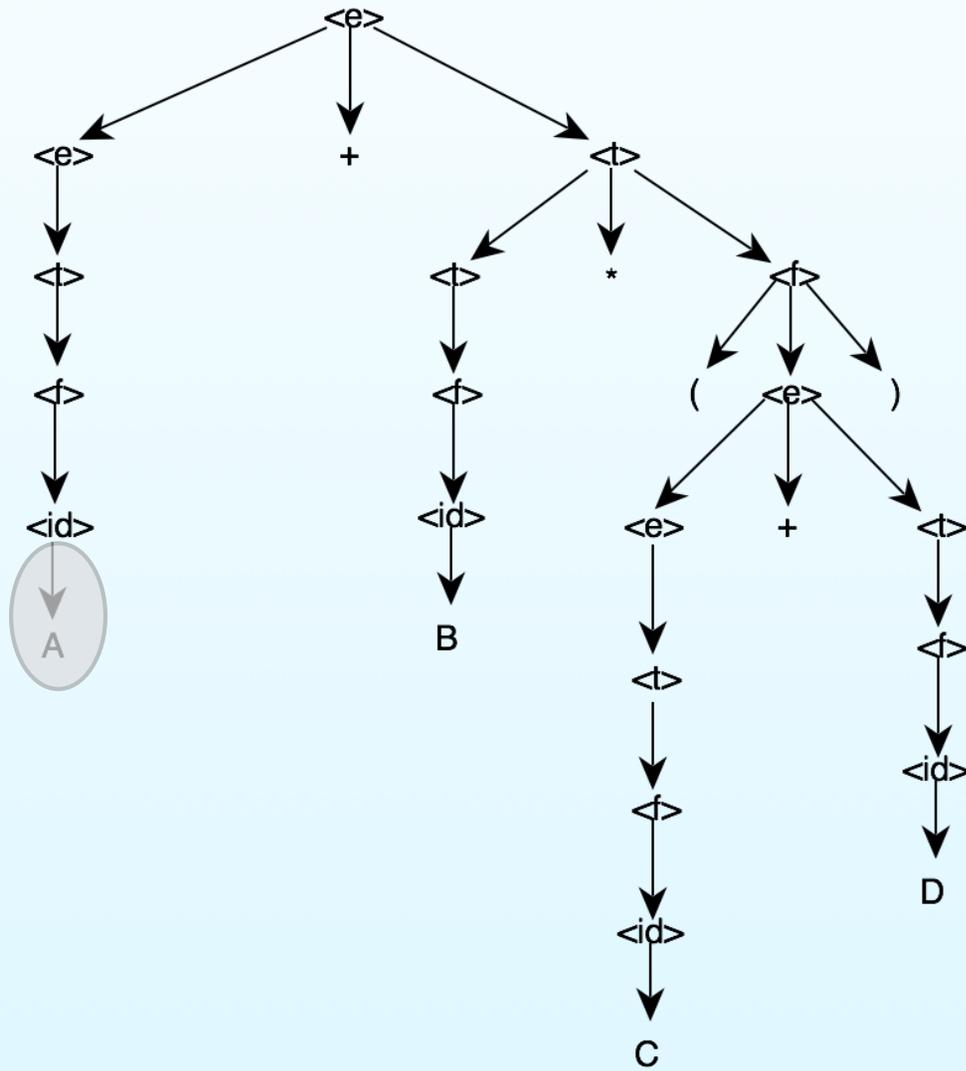
Last production done

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- ~~$A + B * (C + D)$~~

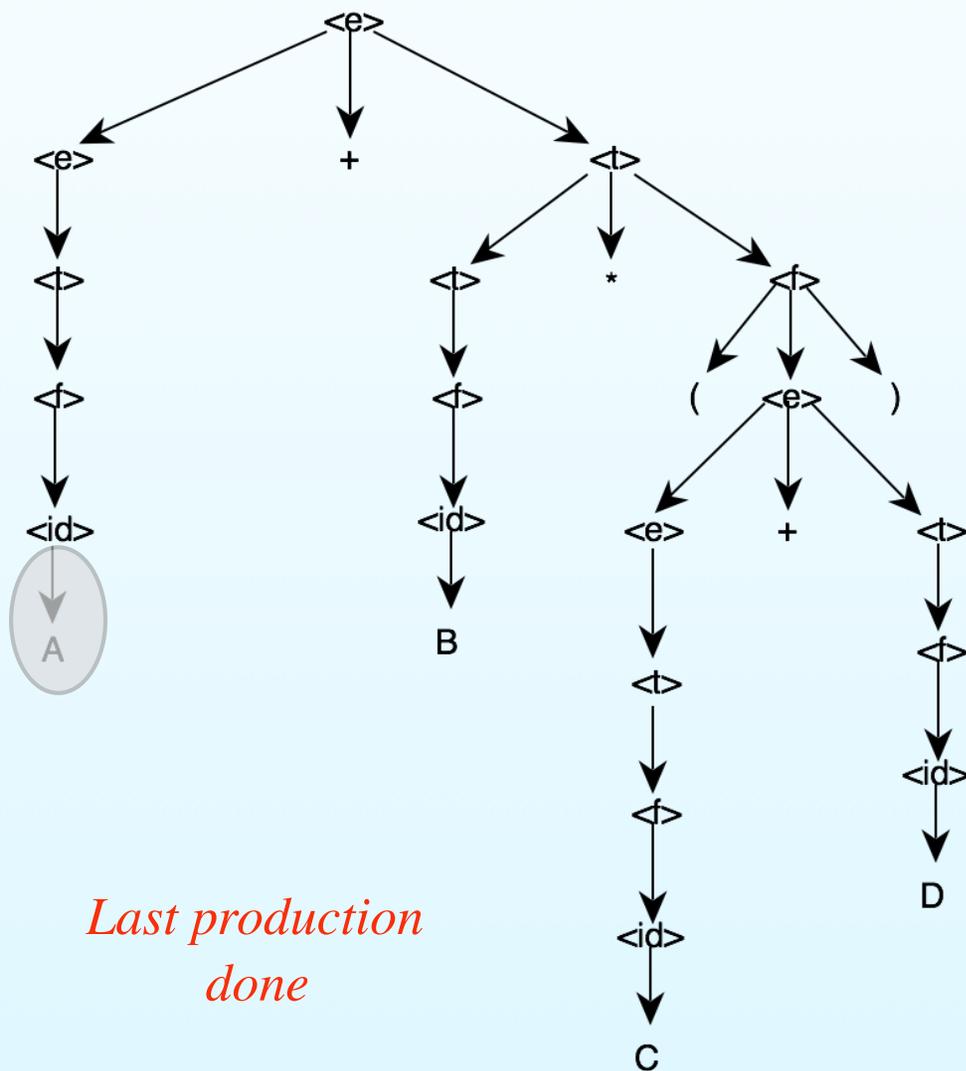


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- ~~$A + B * (C + D)$~~



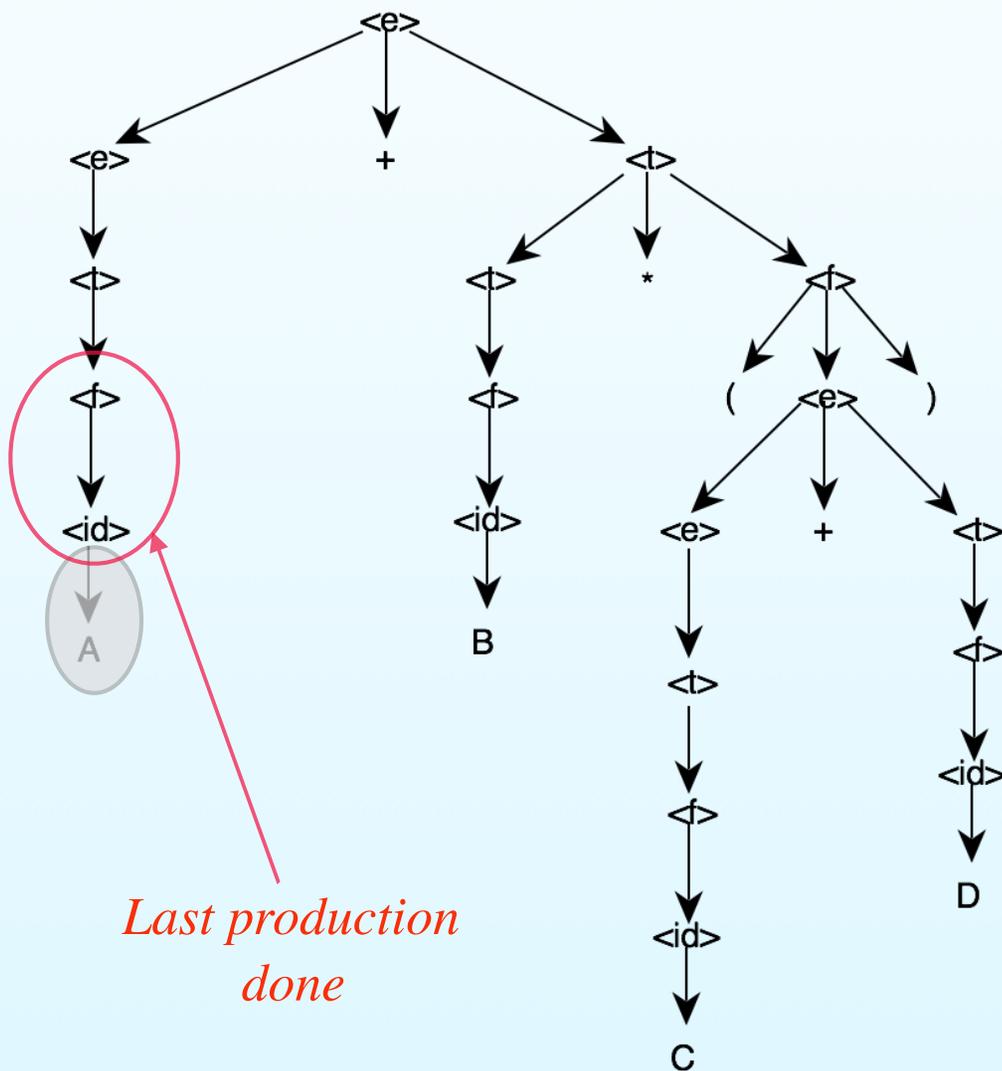
Last production done

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- ~~$A + B * (C + D)$~~



Last production done

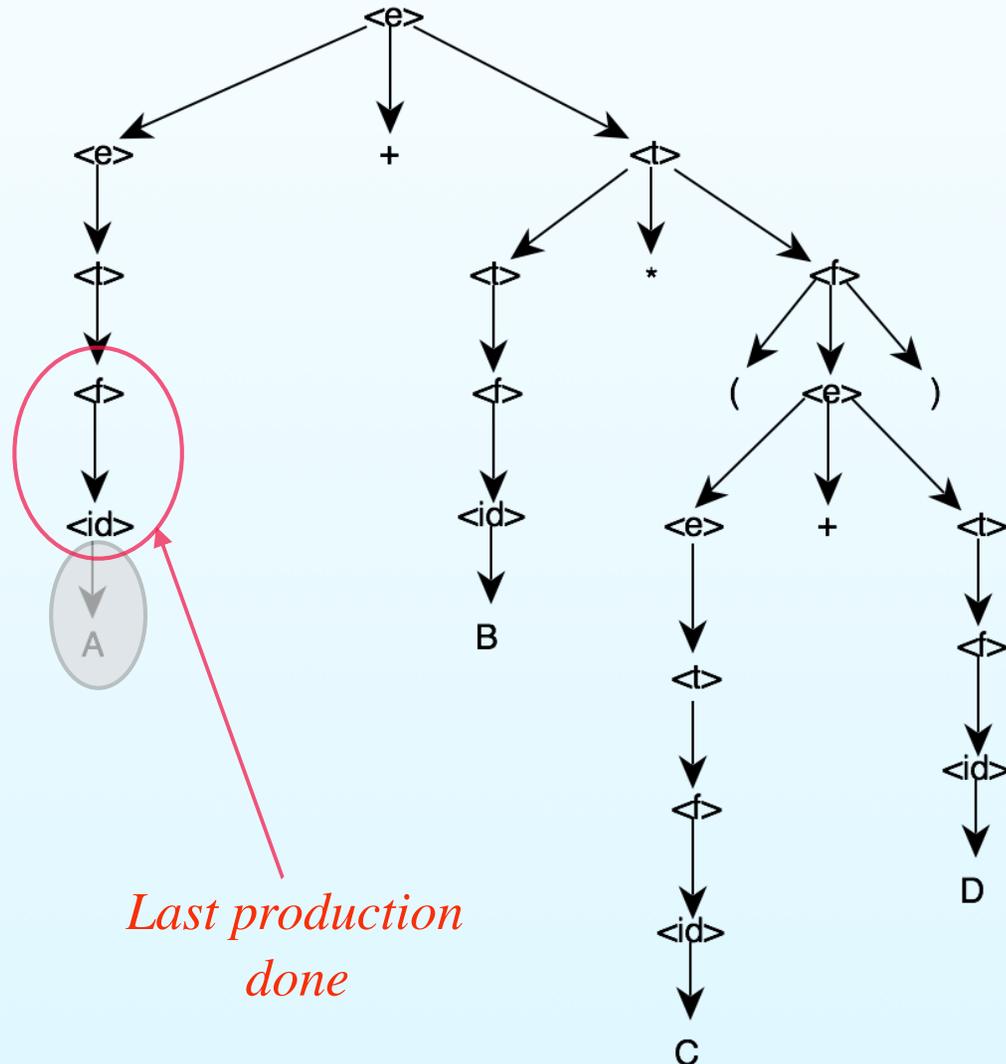
Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- ~~$A + B * (C + D)$~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



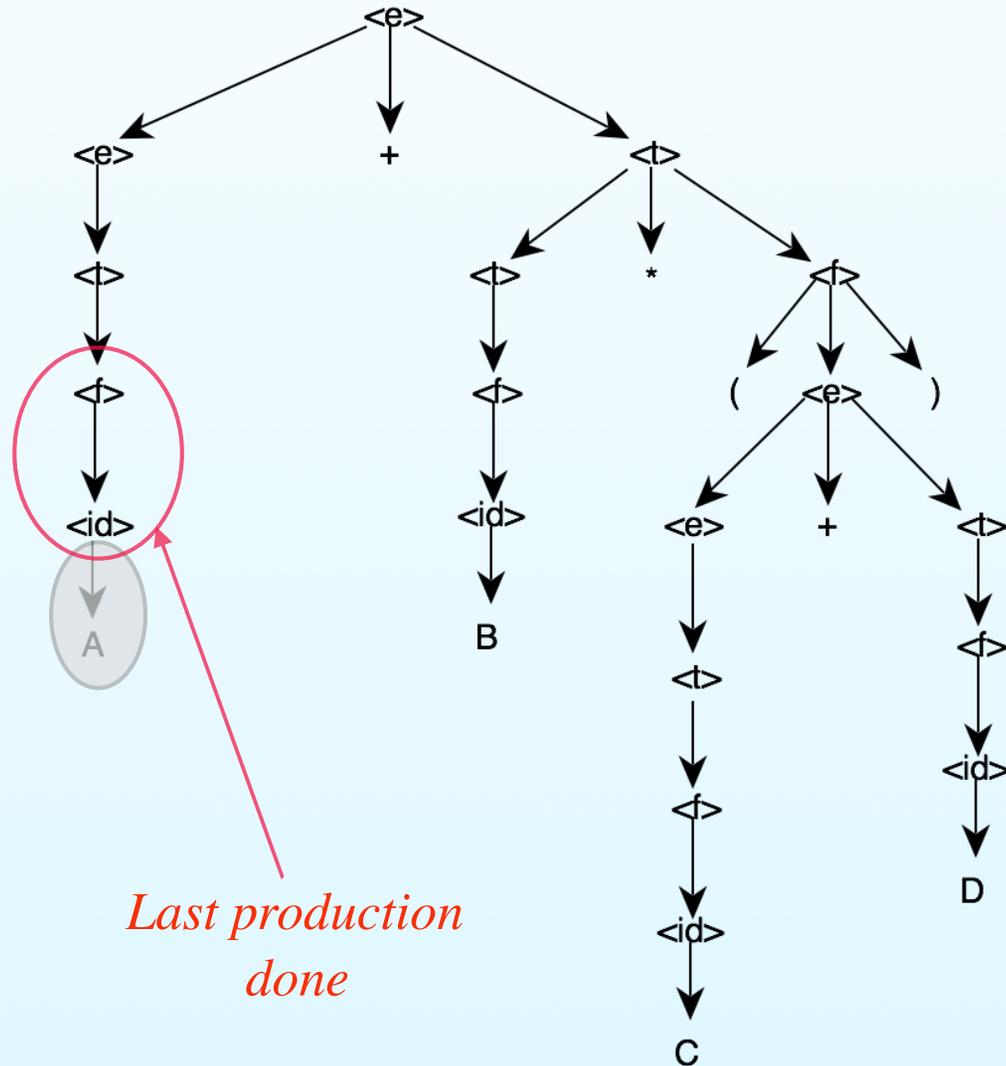
*Last production
done*

Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



Last production done

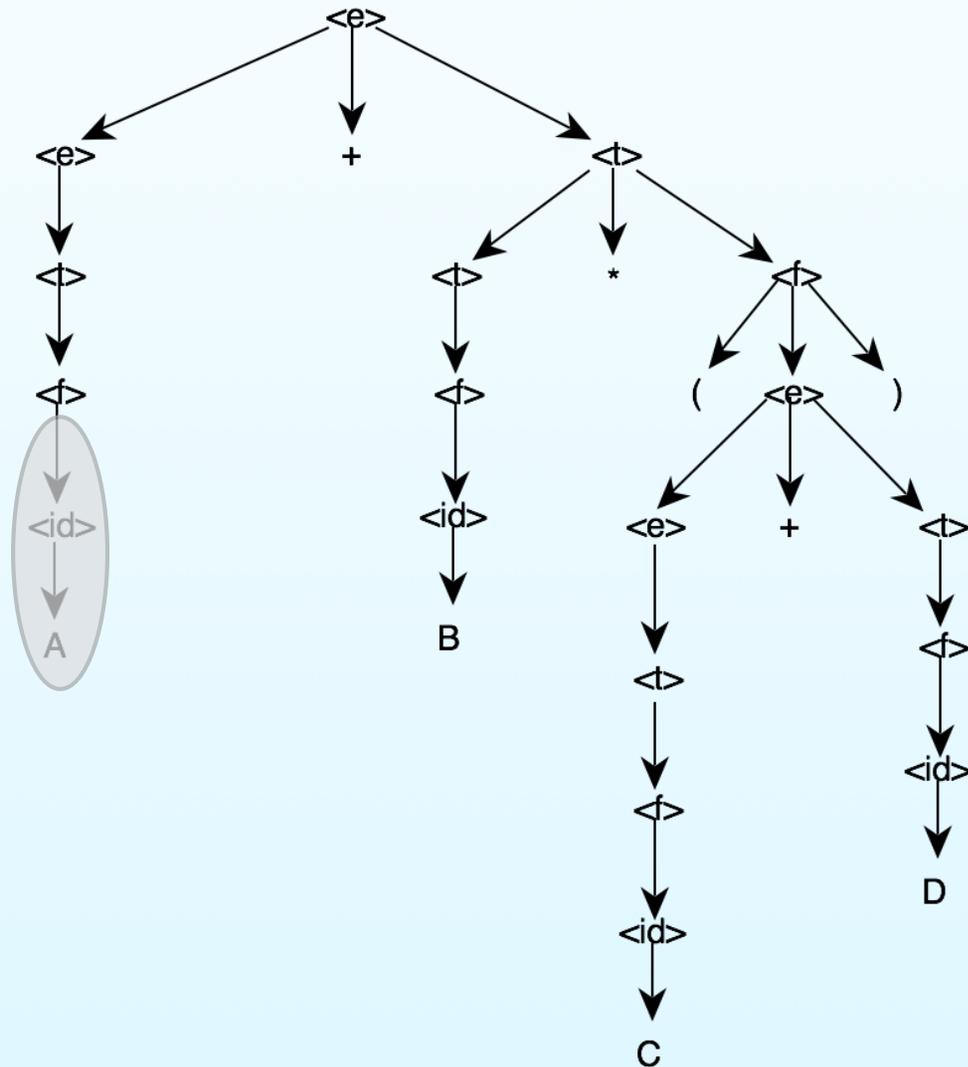
Handle

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B * (C + D)$
 $\langle t \rangle + B * (C + D)$
 $\langle f \rangle + B * (C + D)$
 ~~$\langle id \rangle + B * (C + D)$~~
 ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

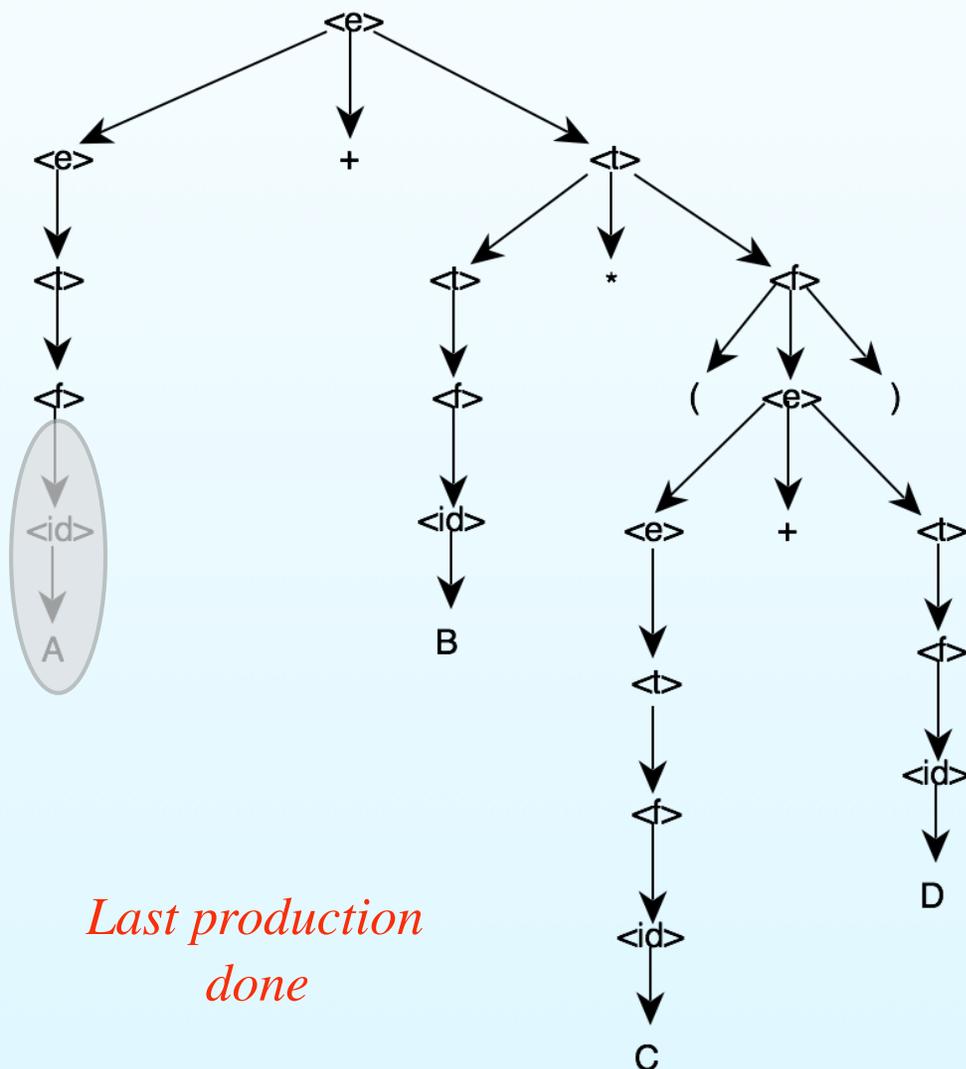


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

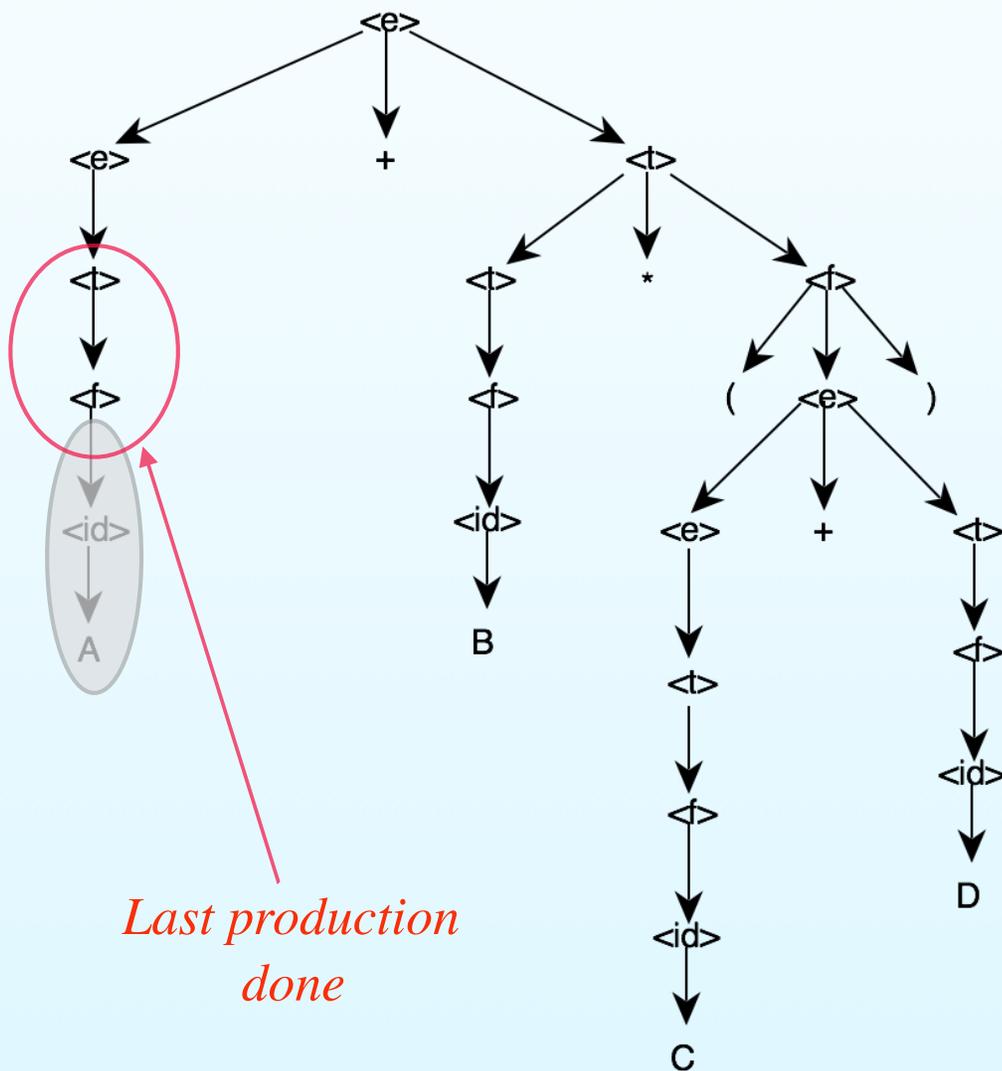


Finding the handle

Parse: A + B * (C + D)

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~



Last production done

Finding the handle

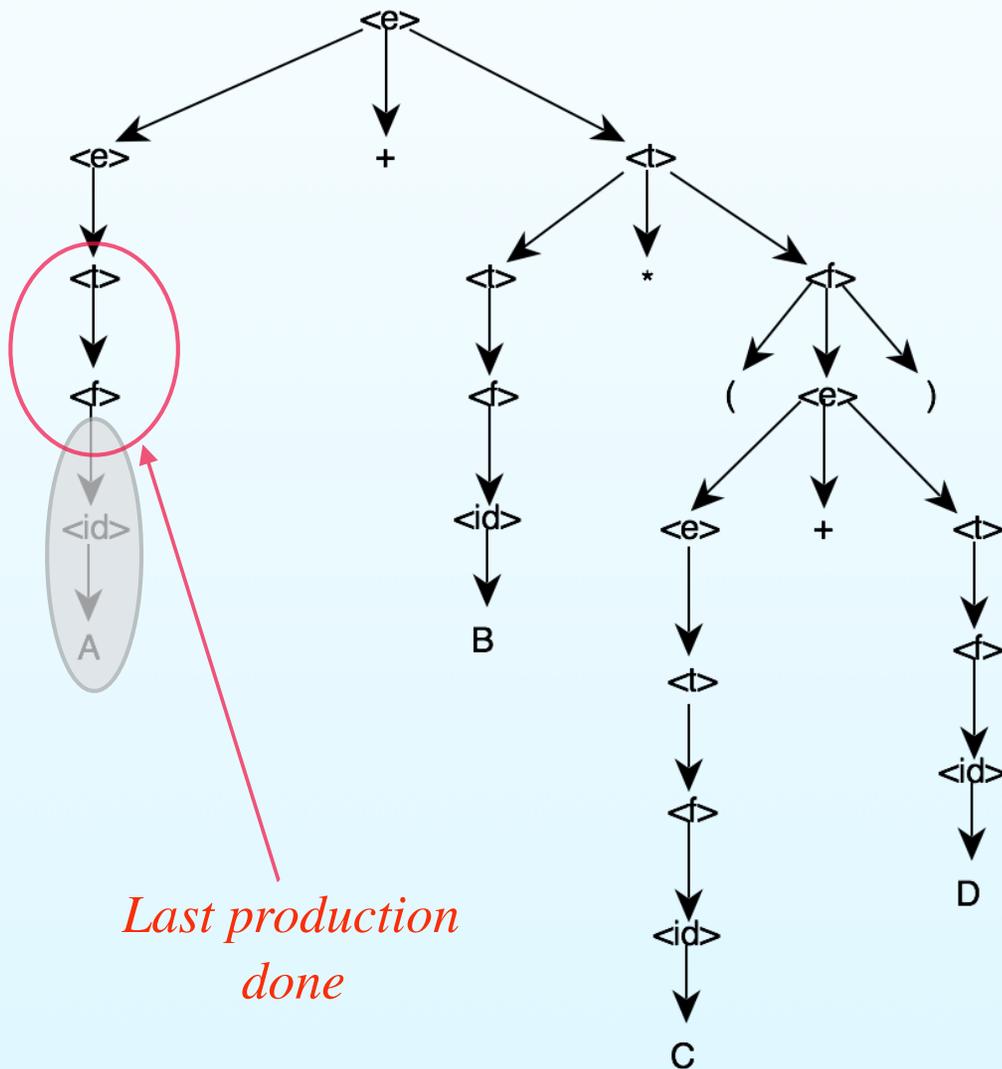
Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

Last production done

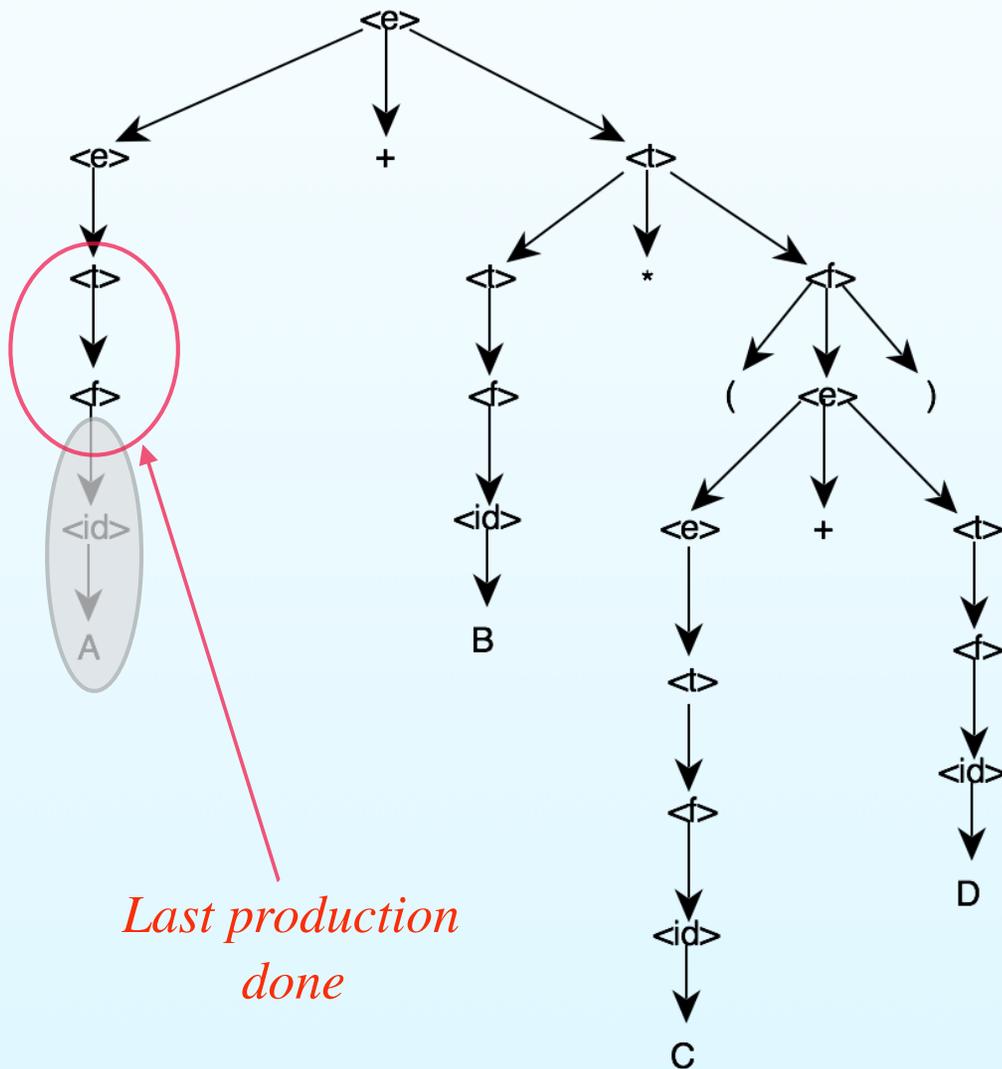


Finding the handle

Parse: A + B * (C + D)

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~



Handle

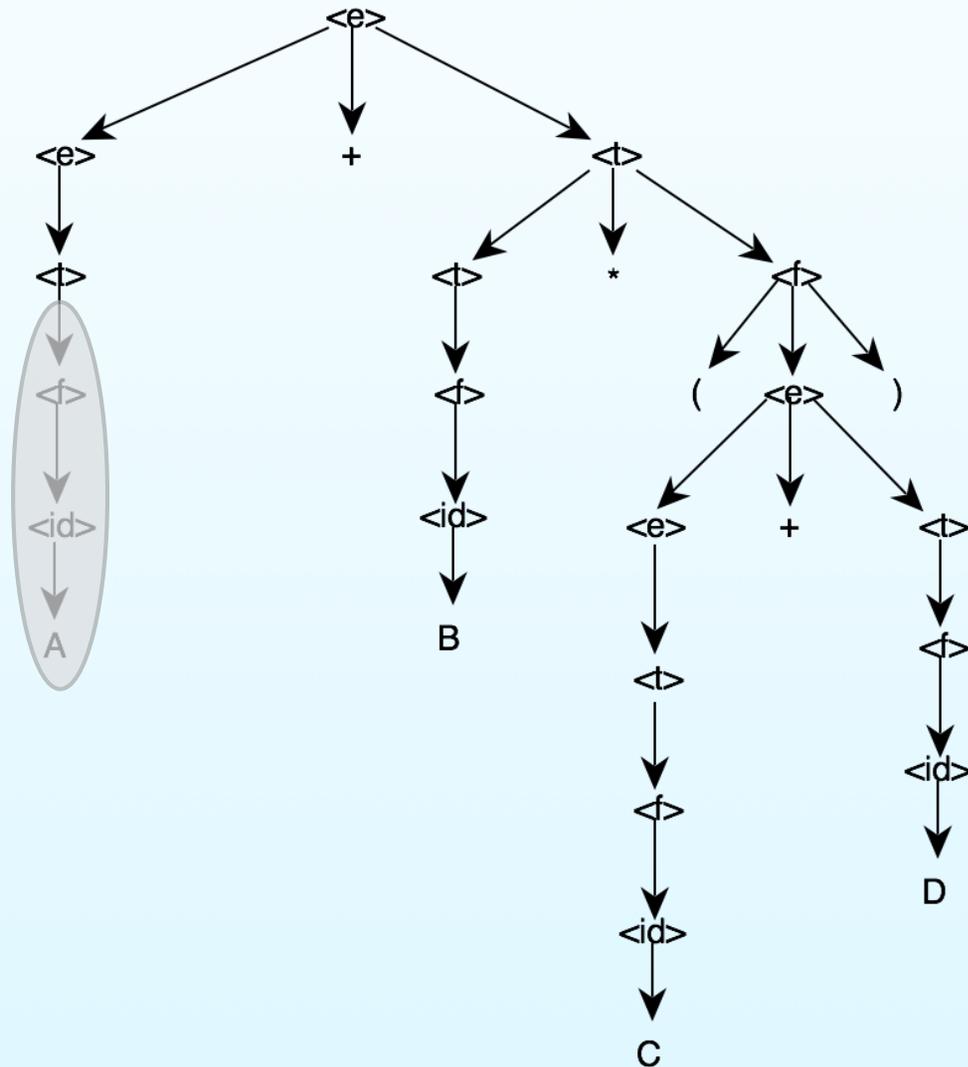
Last production done

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B * (C + D)$
 $\langle t \rangle + B * (C + D)$
 ~~$\langle f \rangle + B * (C + D)$~~
 ~~$\langle id \rangle + B * (C + D)$~~
 ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

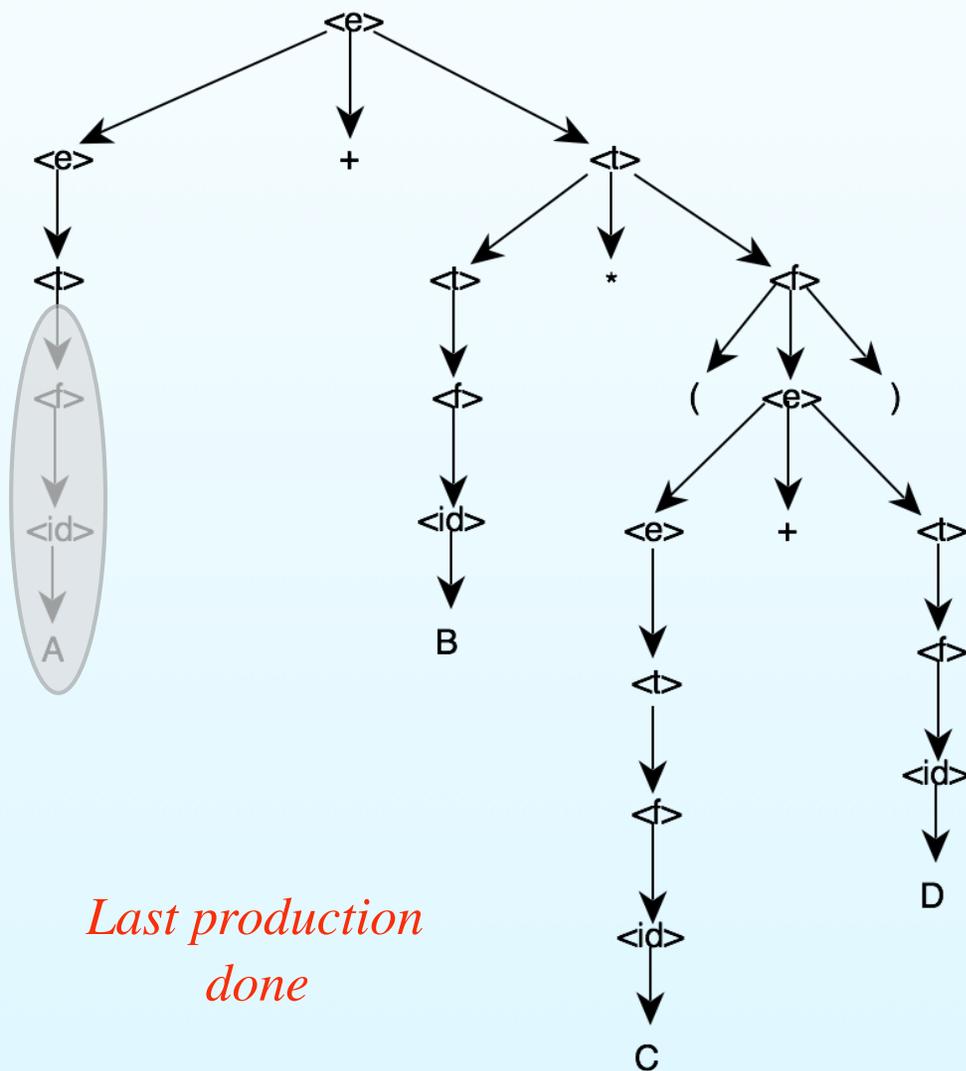


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~



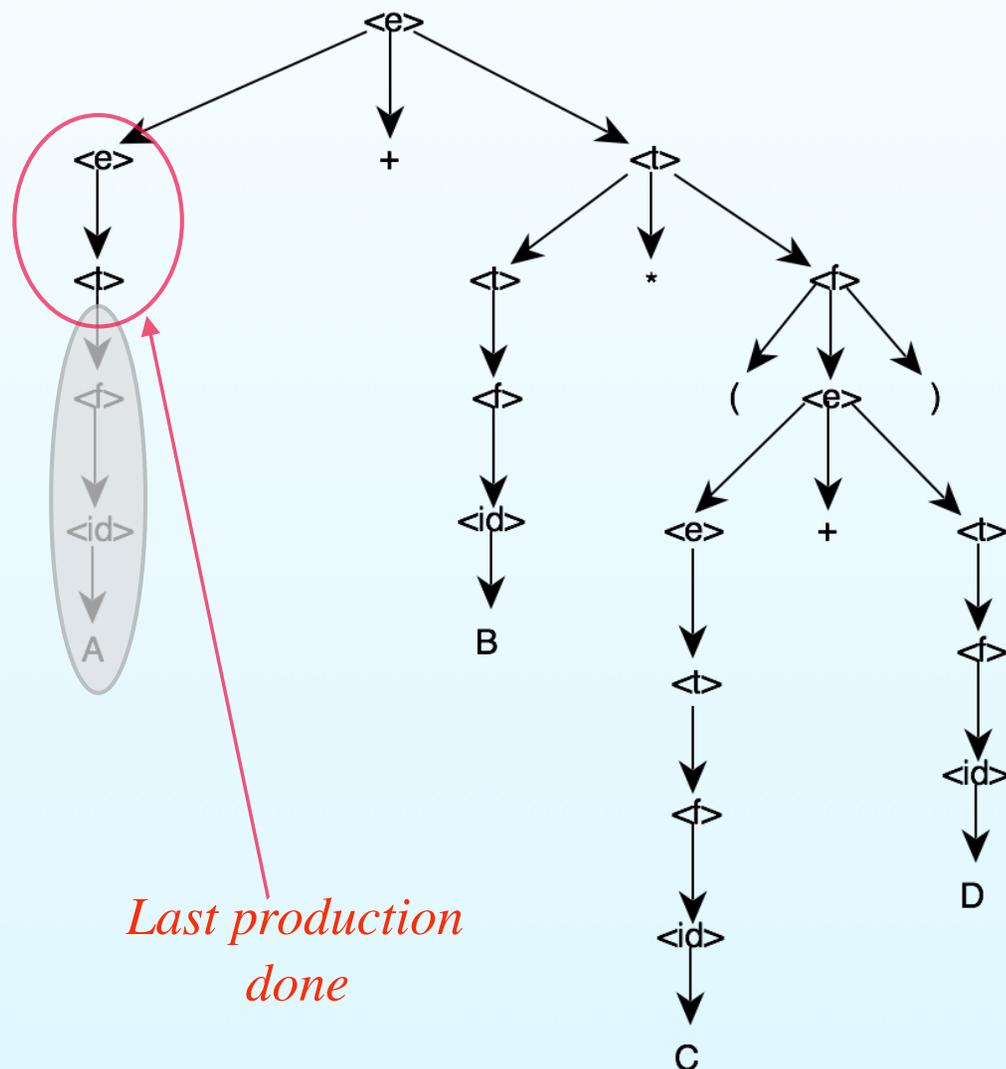
Last production done

Finding the handle

Parse: A + B * (C + D)

- <e>
- <e> + <t>
- <e> + <t> * <f>
- <e> + <t> * (<e>)
- <e> + <t> * (<e> + <t>)
- <e> + <t> * (<e> + <f>)
- <e> + <t> * (<e> + <id>)
- <e> + <t> * (<e> + D)
- <e> + <t> * (<t> + D)
- <e> + <t> * (<f> + D)
- <e> + <t> * (<id> + D)
- <e> + <t> * (C + D)
- <e> + <f> * (C + D)
- <e> + <id> * (C + D)
- <e> + B * (C + D)
- <t> + B * (C + D)
- ~~<f> + B * (C + D)~~
- ~~<id> + B * (C + D)~~
- ~~A + B * (C + D)~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



Last production done

Finding the handle

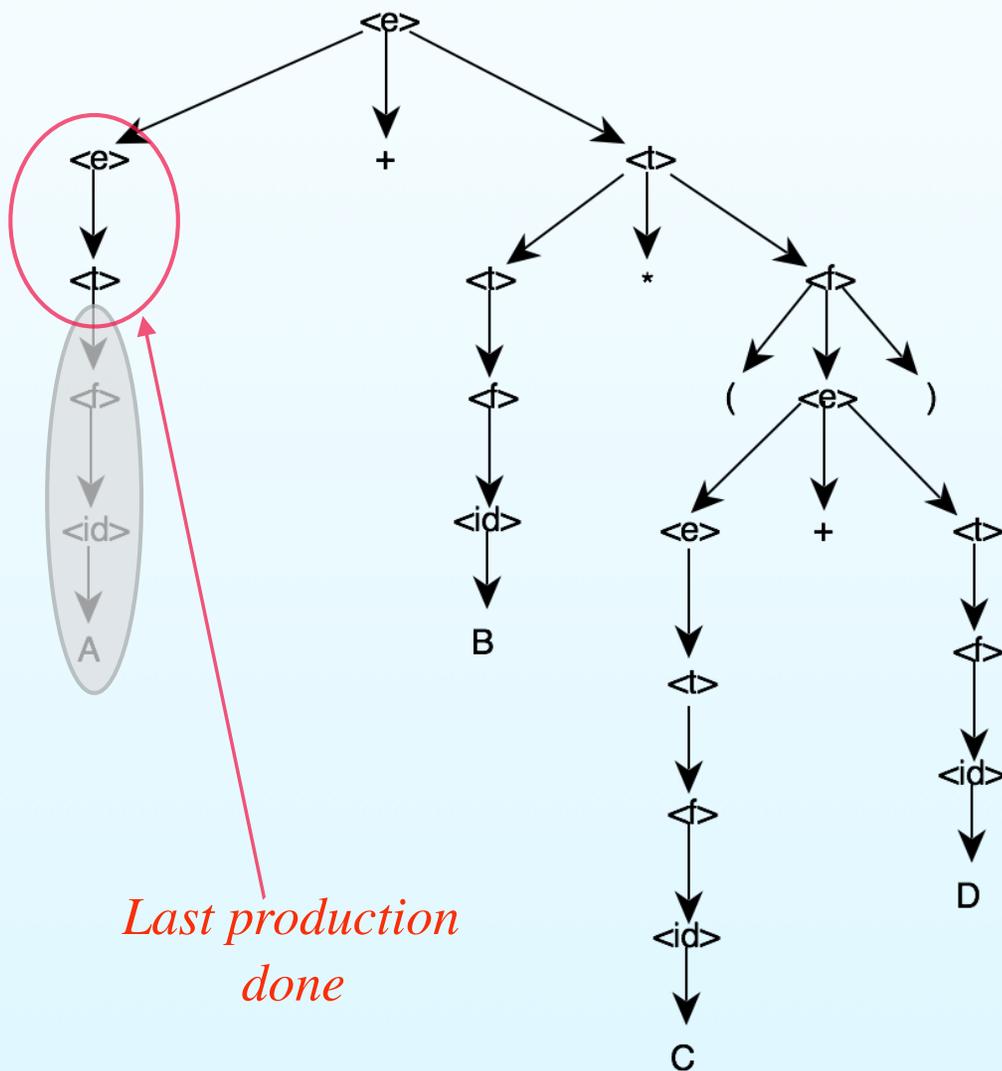
Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

Last production done



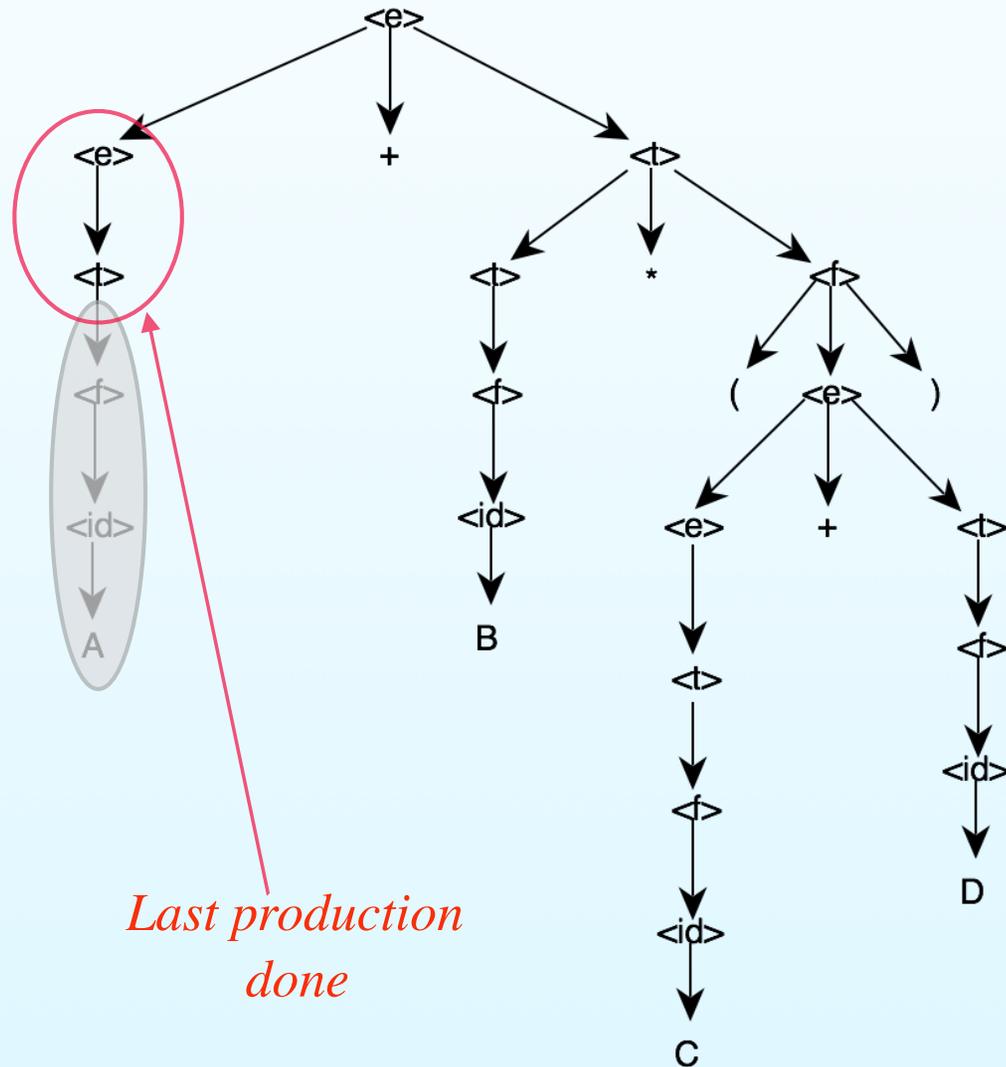
Finding the handle

Parse: A + B * (C + D)

- <e>
- <e> + <t>
- <e> + <t> * <f>
- <e> + <t> * (<e>)
- <e> + <t> * (<e> + <t>)
- <e> + <t> * (<e> + <f>)
- <e> + <t> * (<e> + <id>)
- <e> + <t> * (<e> + D)
- <e> + <t> * (<t> + D)
- <e> + <t> * (<f> + D)
- <e> + <t> * (<id> + D)
- <e> + <t> * (C + D)
- <e> + <f> * (C + D)
- <e> + <id> * (C + D)
- <e> + B * (C + D)
- <t> + B * (C + D)**
- ~~<f> + B * (C + D)~~
- ~~<id> + B * (C + D)~~
- ~~A + B * (C + D)~~

Handle

<e> → <e> + <t> | <t>
 <t> → <t> * <f> | <f>
 <f> → (<e>) | <id>



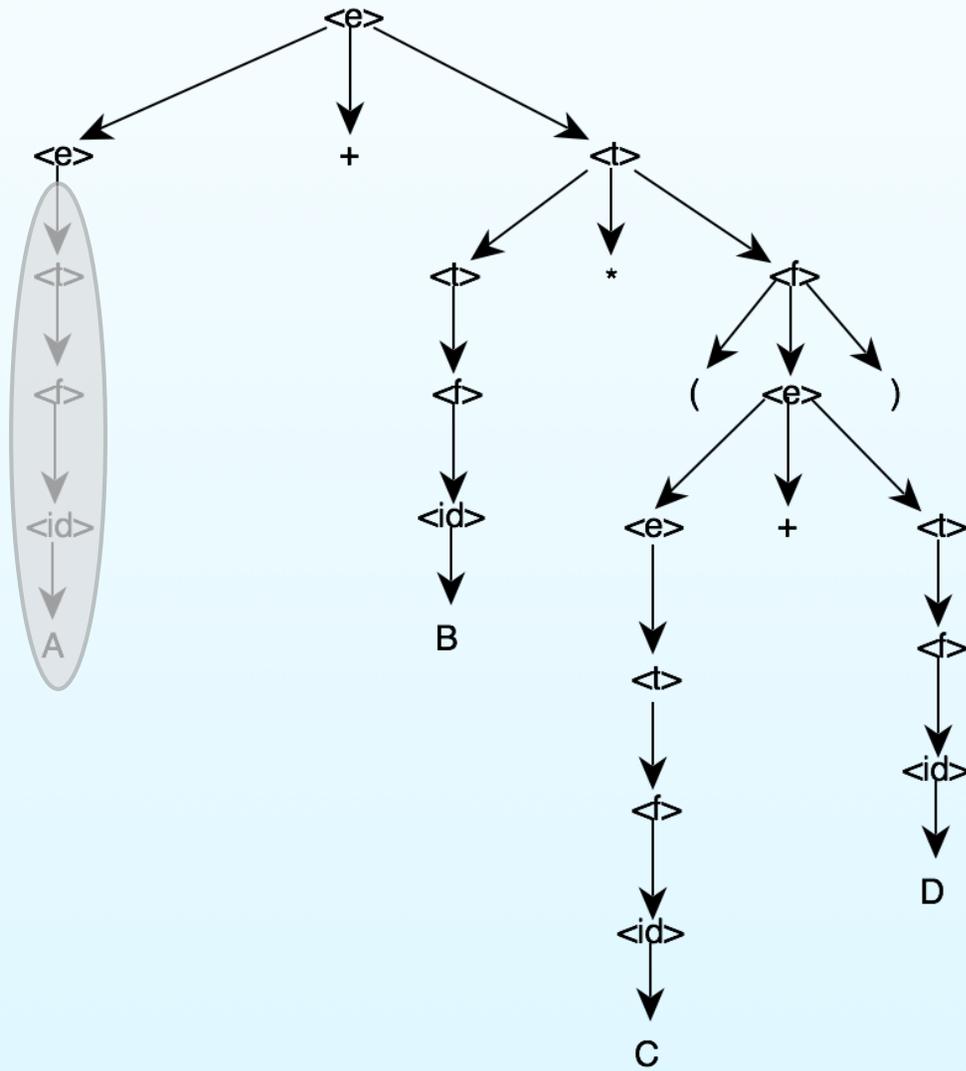
Last production done

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 $\langle e \rangle + \langle f \rangle * (C + D)$
 $\langle e \rangle + \langle id \rangle * (C + D)$
 $\langle e \rangle + B * (C + D)$
 ~~$\langle t \rangle + B * (C + D)$~~
 ~~$\langle f \rangle + B * (C + D)$~~
 ~~$\langle id \rangle + B * (C + D)$~~
 ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

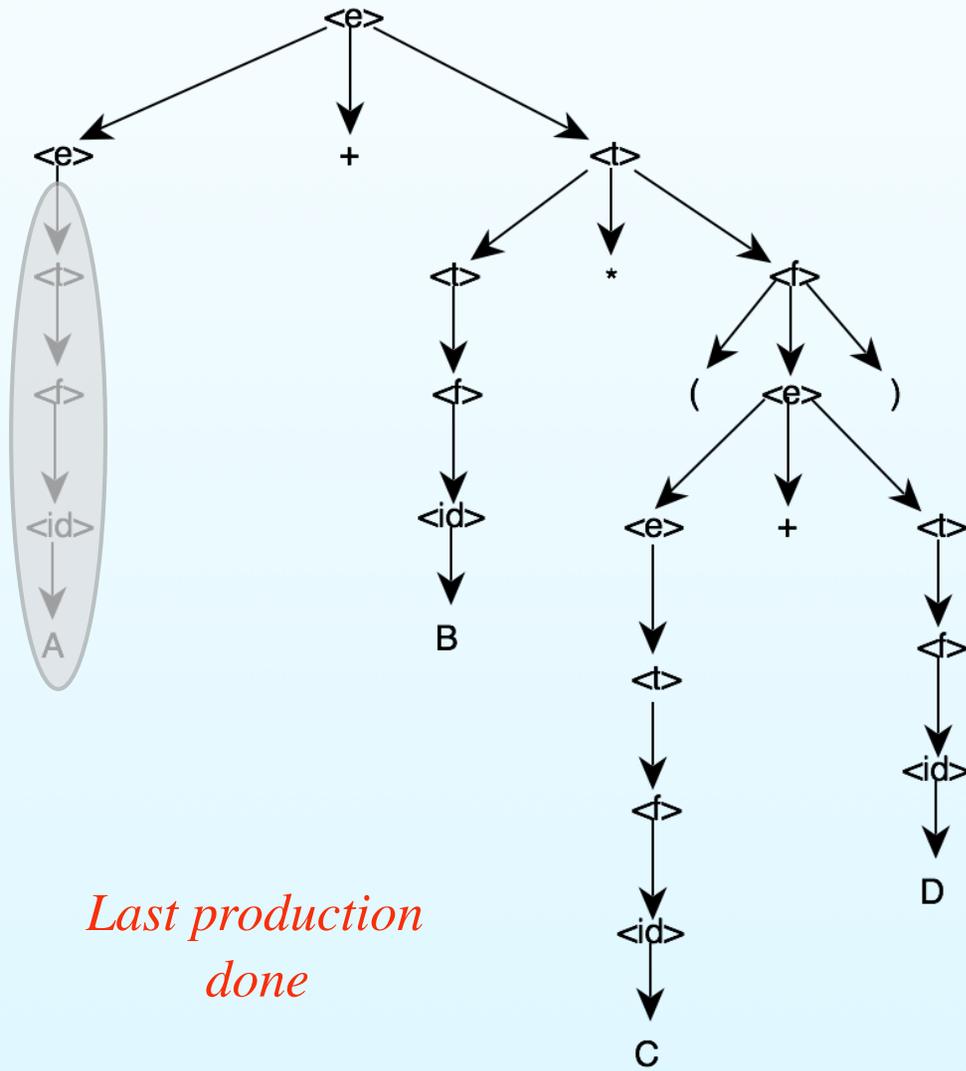


Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



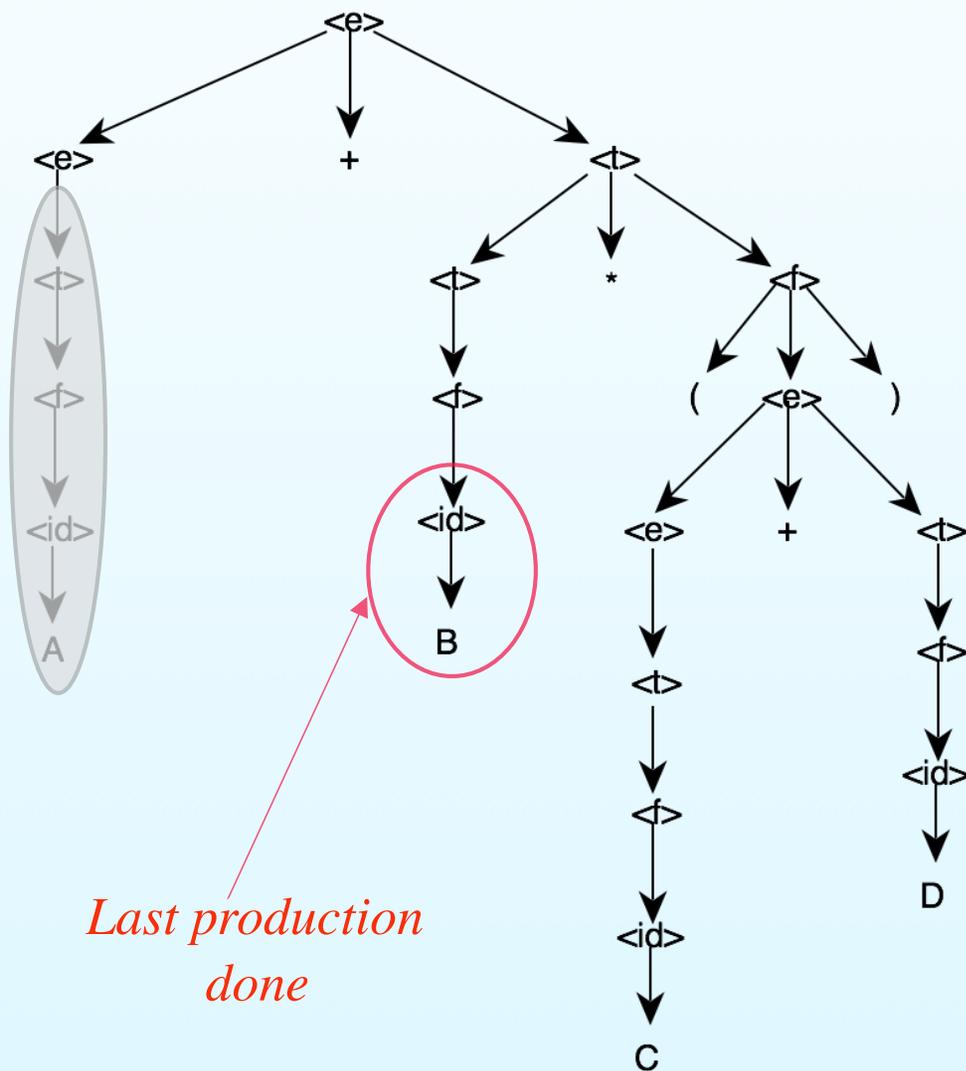
Last production done

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~



Last production done

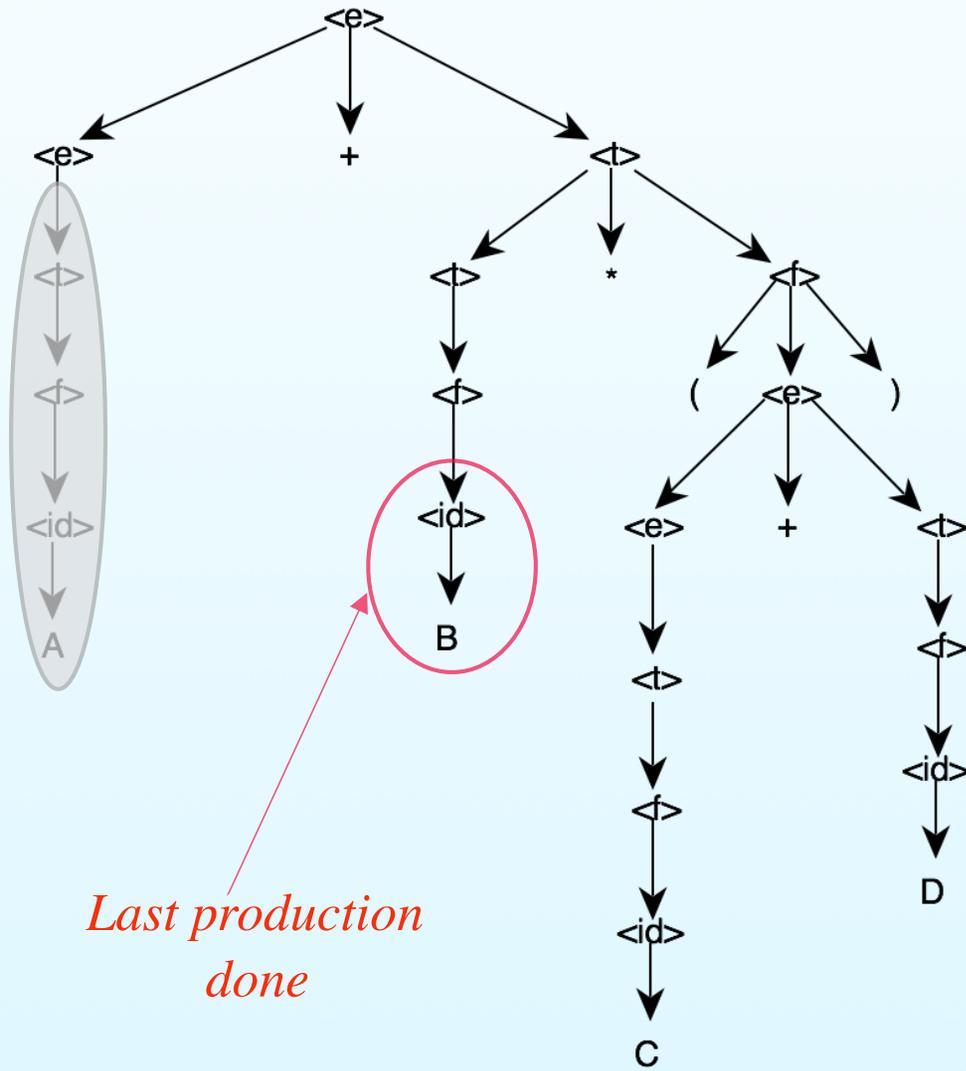
Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



Last production done

Finding the handle

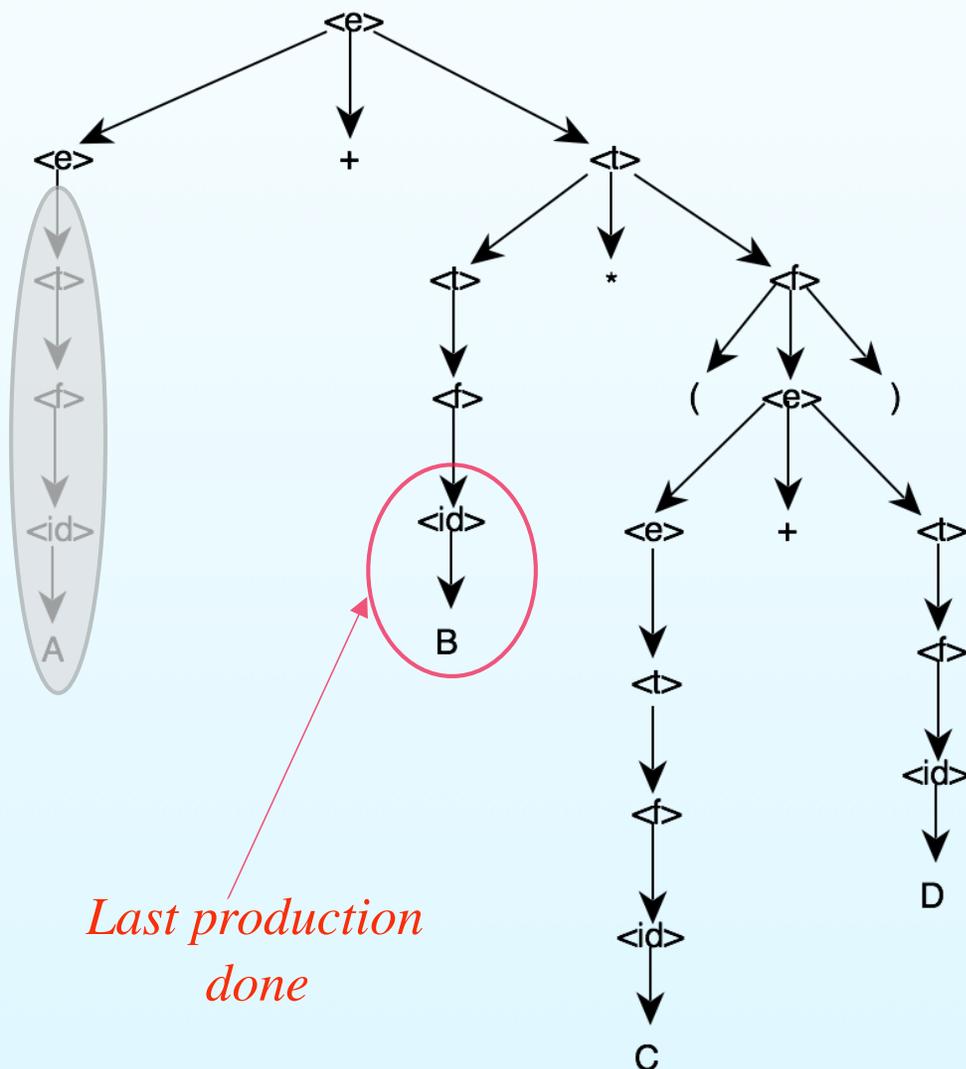
Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

Last production done

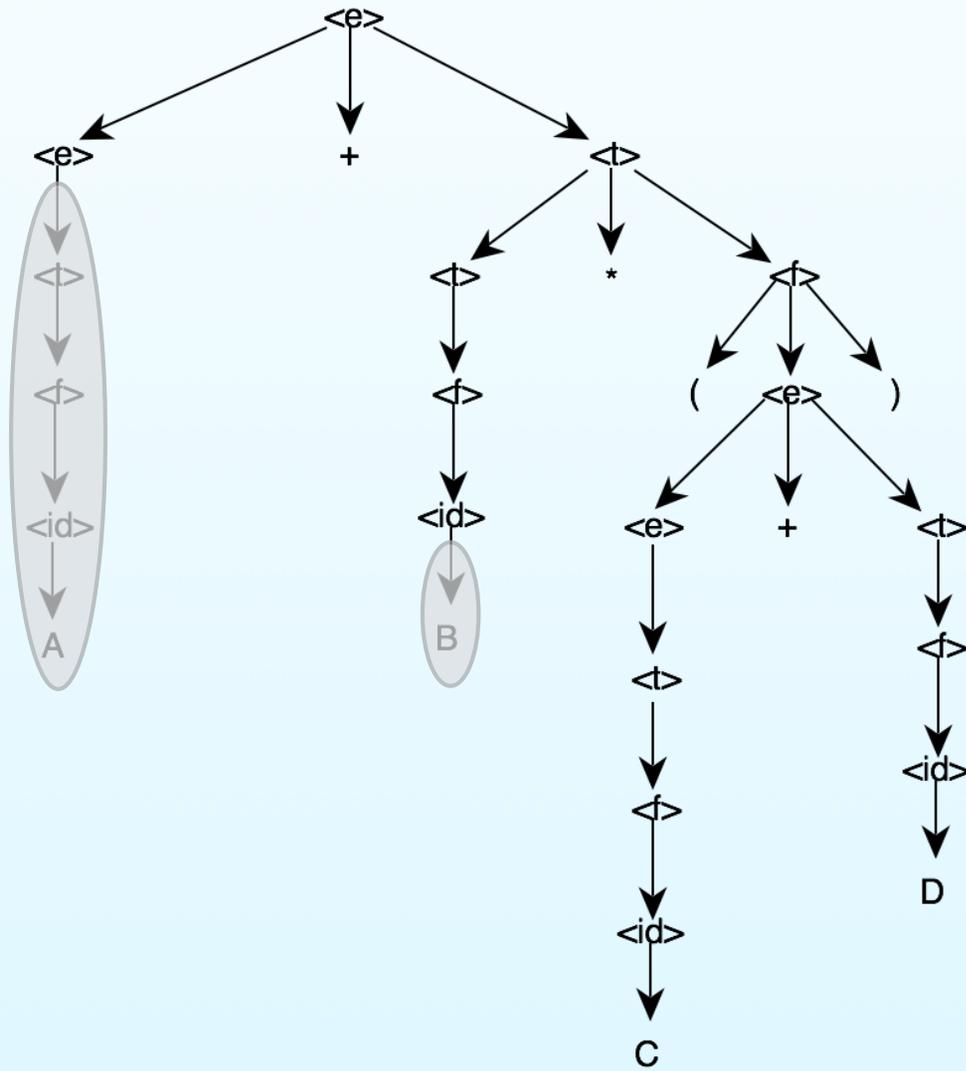


Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

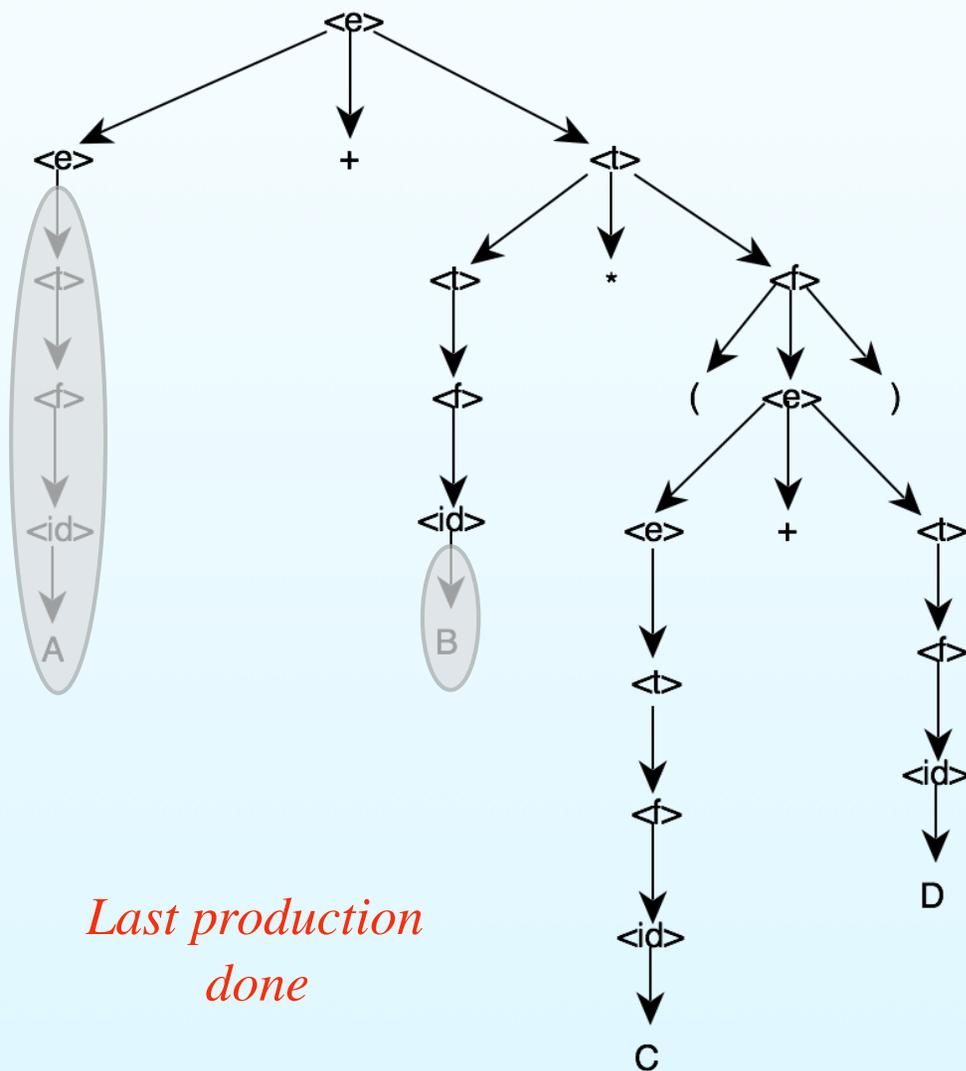


Finding the handle

Parse: A + B * (C + D)

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~



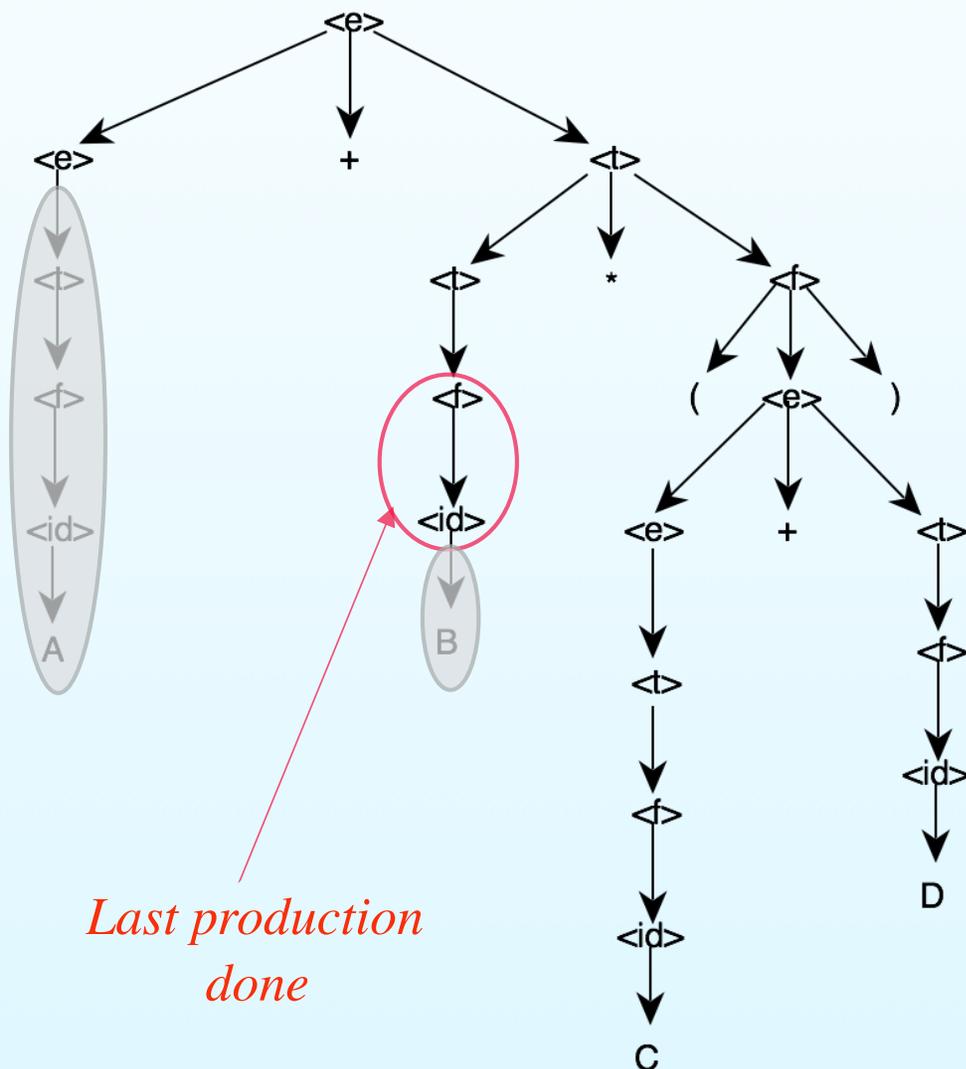
Last production done

Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~



Last production done

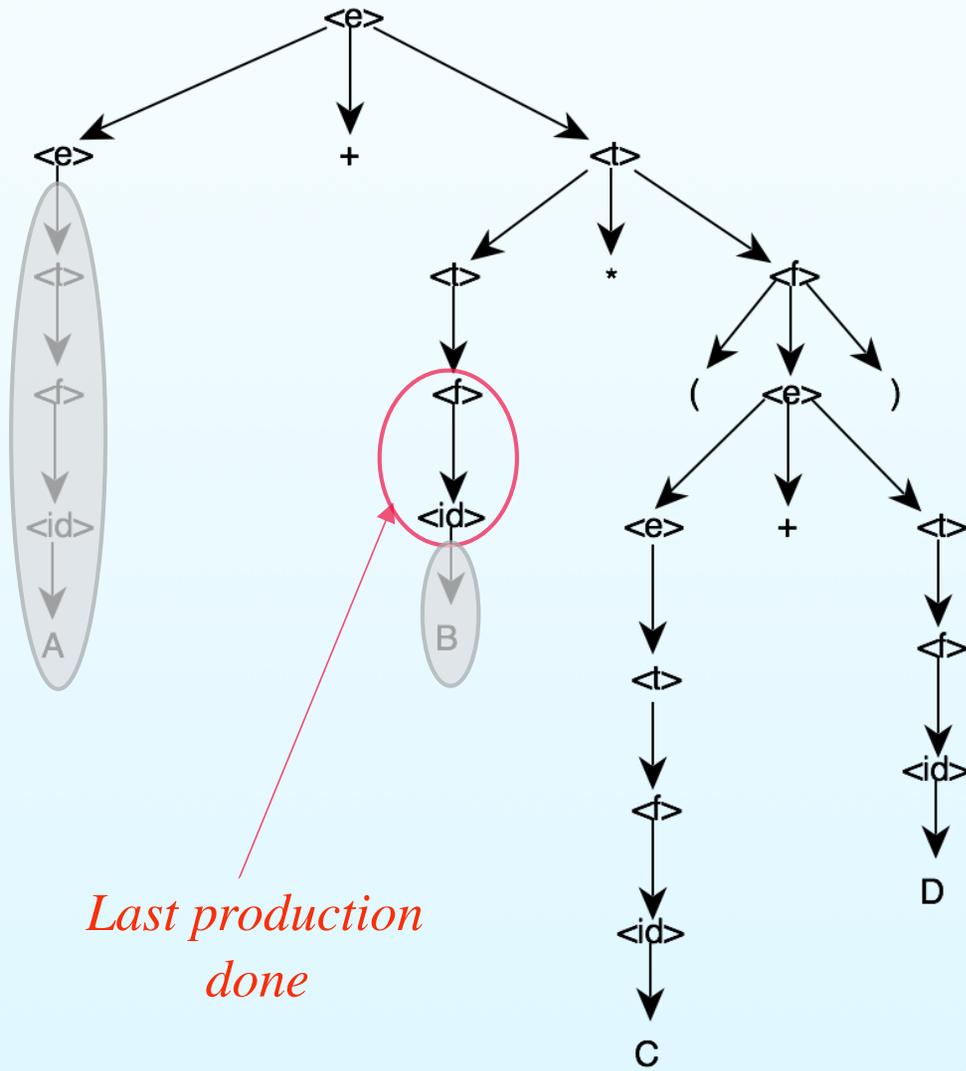
Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



Last production done

Finding the handle

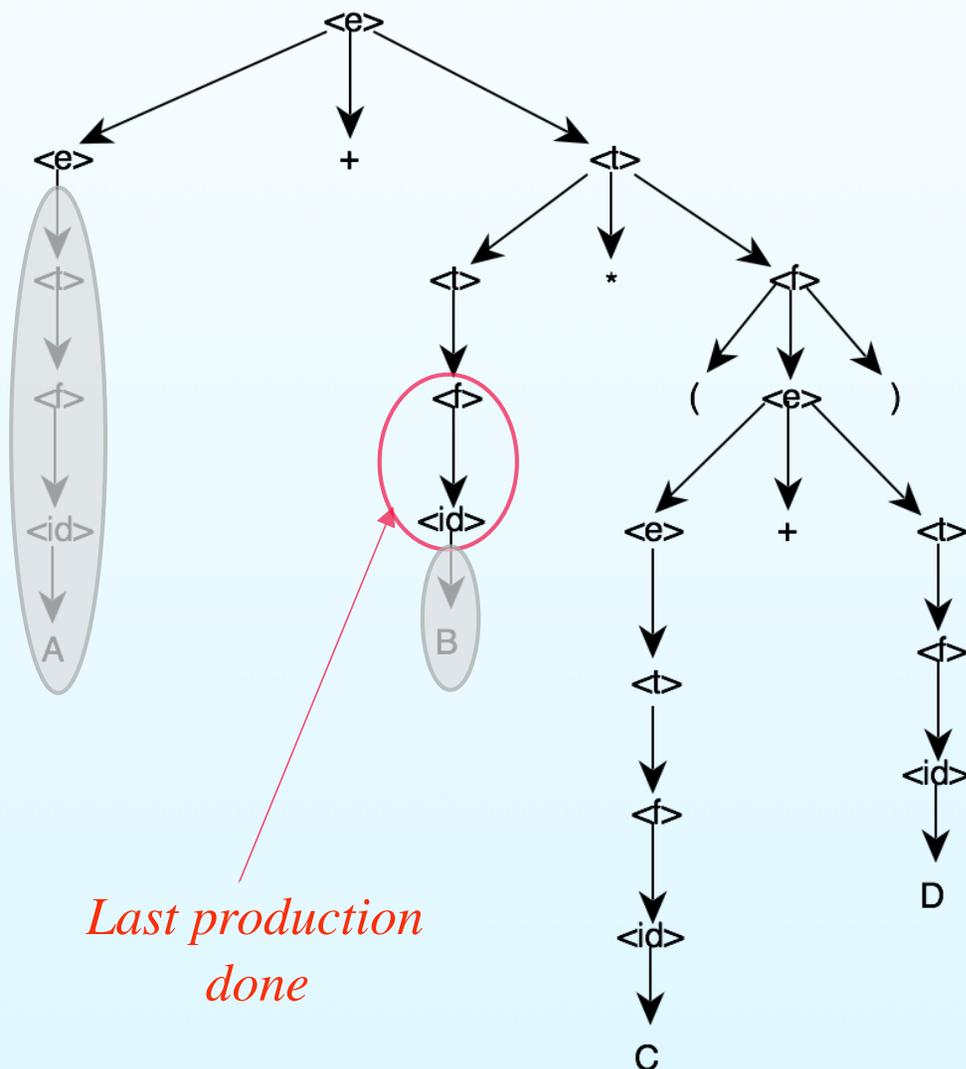
Parse: A + B * (C + D)

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

Last production done

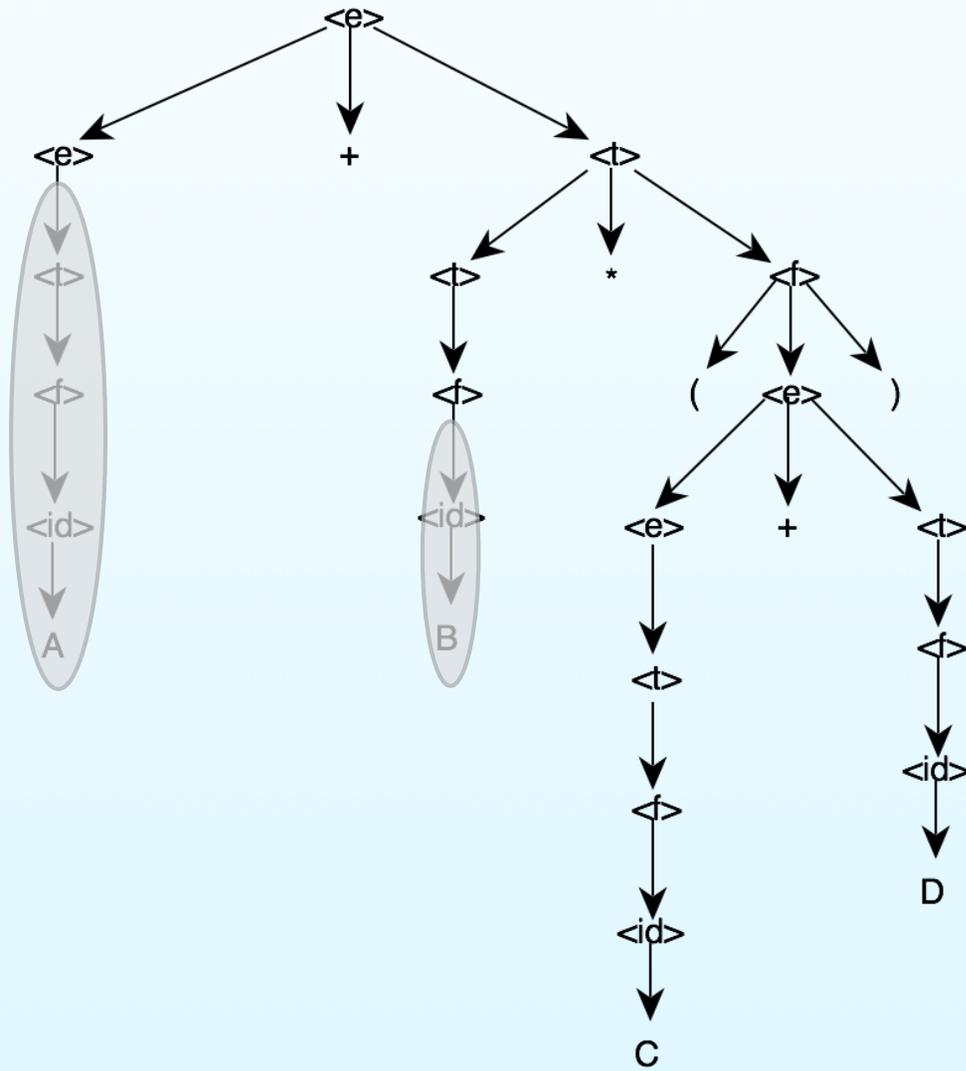


Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

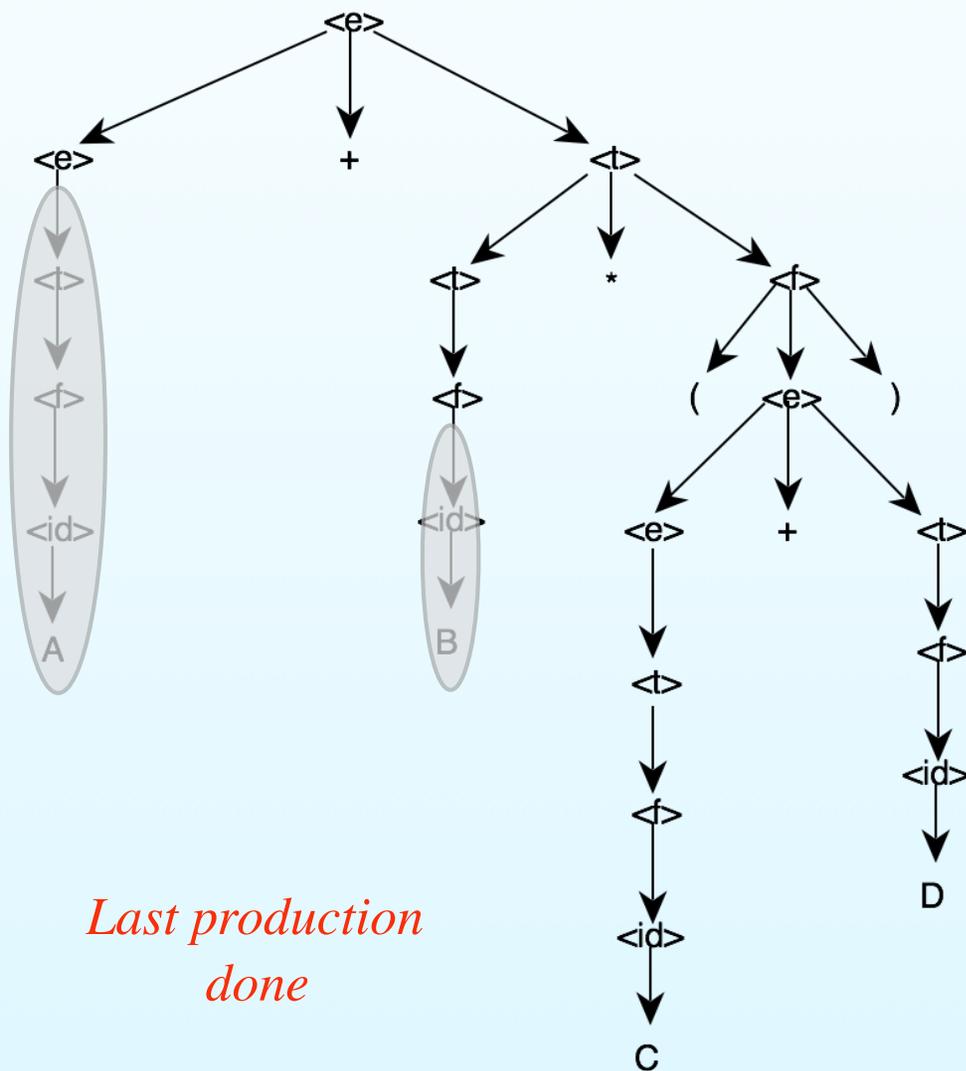


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

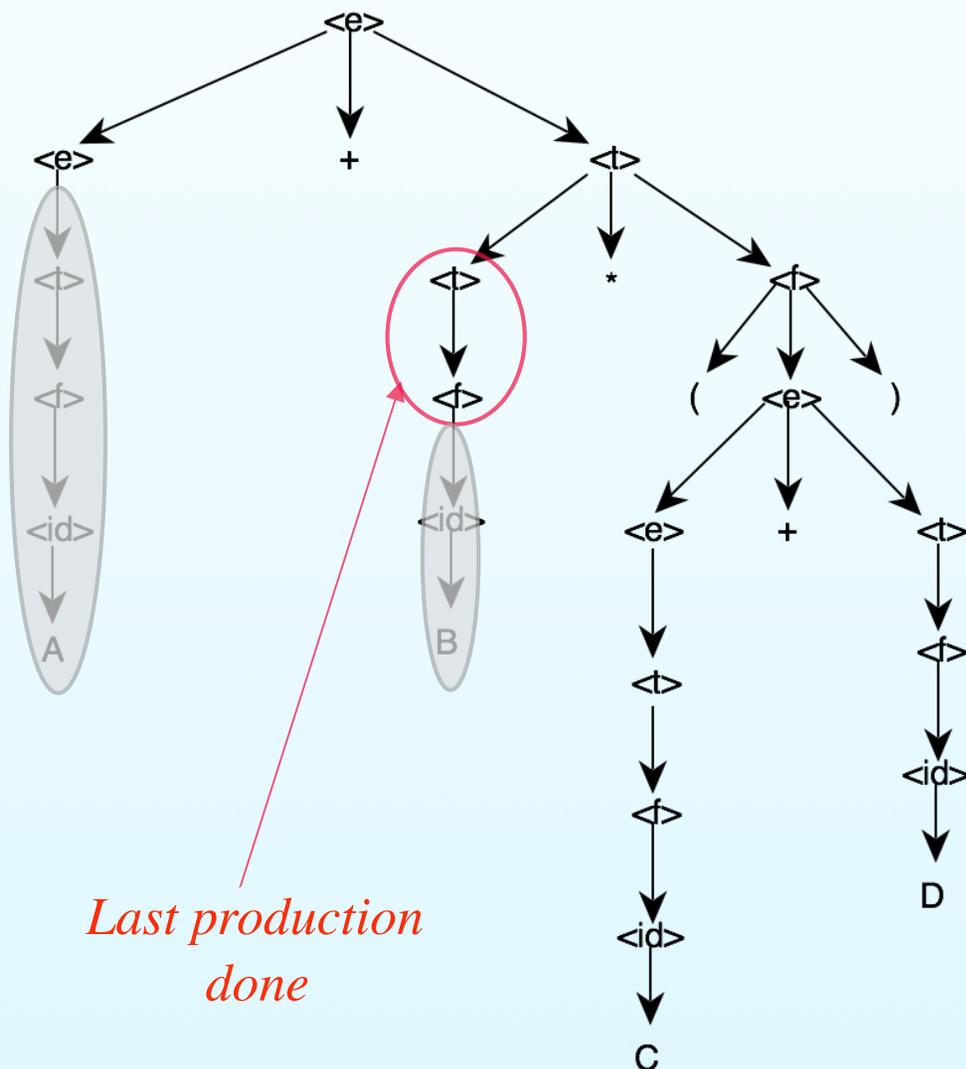


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~



Last production done

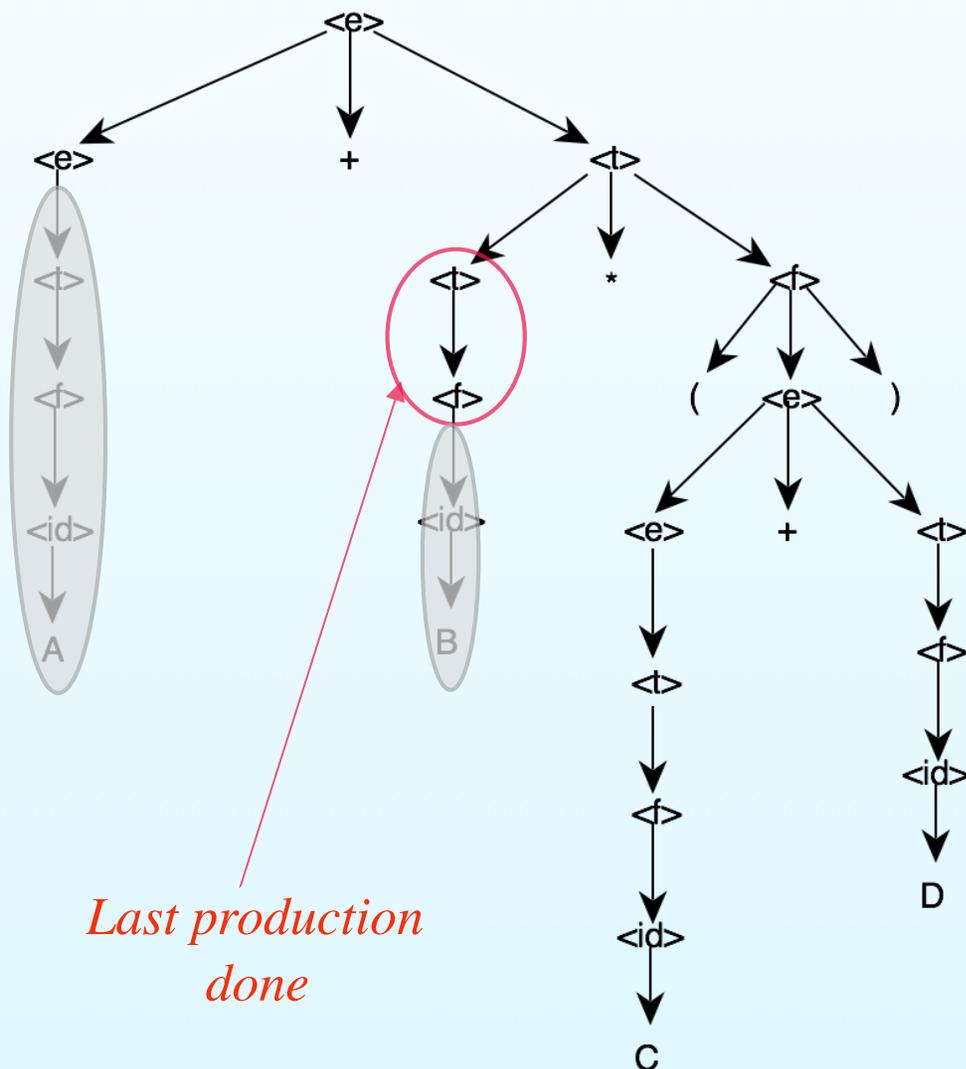
Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



Last production done

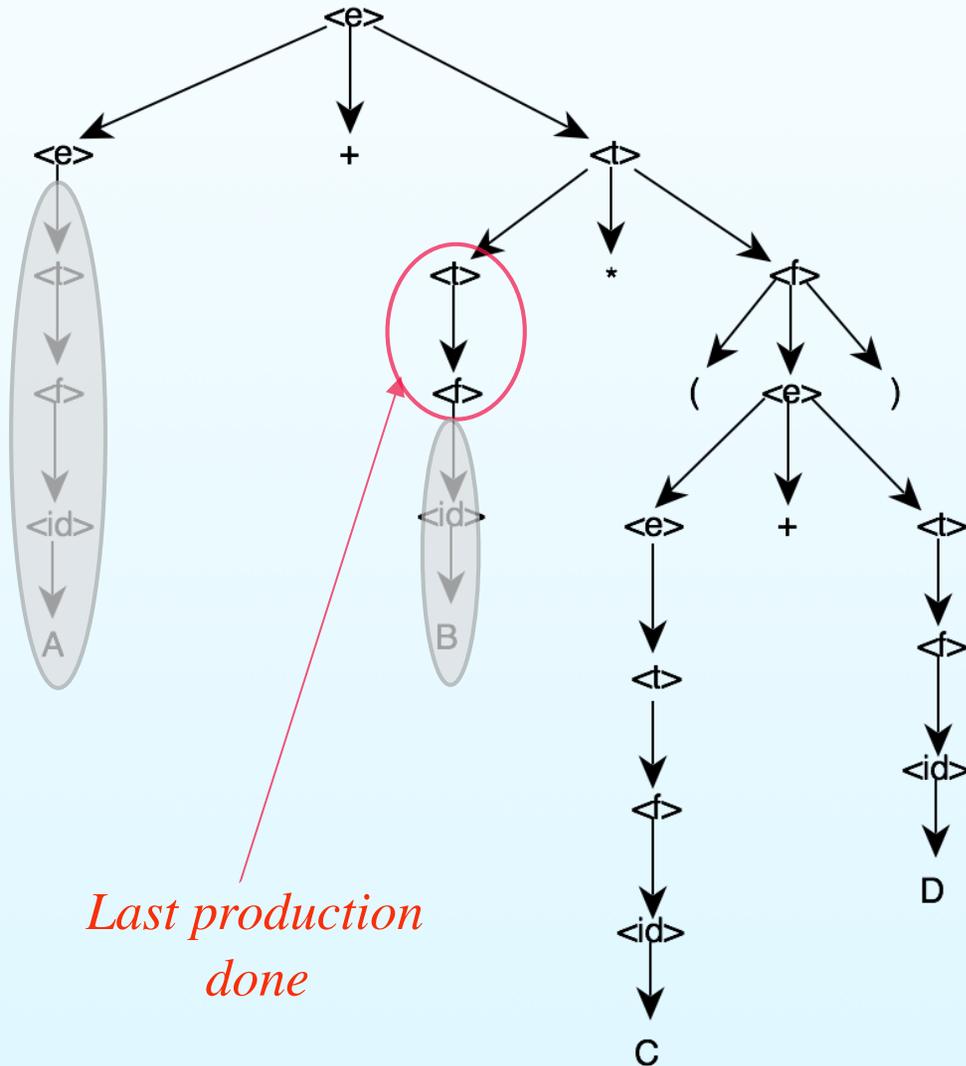
Finding the handle

Parse: A + B * (C + D)

- <e>
- <e> + <t>
- <e> + <t> * <f>
- <e> + <t> * (<e>)
- <e> + <t> * (<e> + <t>)
- <e> + <t> * (<e> + <f>)
- <e> + <t> * (<e> + <id>)
- <e> + <t> * (<e> + D)
- <e> + <t> * (<t> + D)
- <e> + <t> * (<f> + D)
- <e> + <t> * (<id> + D)
- <e> + <t> * (C + D)
- <e> + **<f>** * (C + D)
- ~~<e> + <id> * (C + D)~~
- ~~<e> + B * (C + D)~~
- ~~<t> + B * (C + D)~~
- ~~<f> + B * (C + D)~~
- ~~<id> + B * (C + D)~~
- ~~A + B * (C + D)~~

Handle

<e> → <e> + <t> | <t>
 <t> → <t> * <f> | <f>
 <f> → (<e>) | <id>

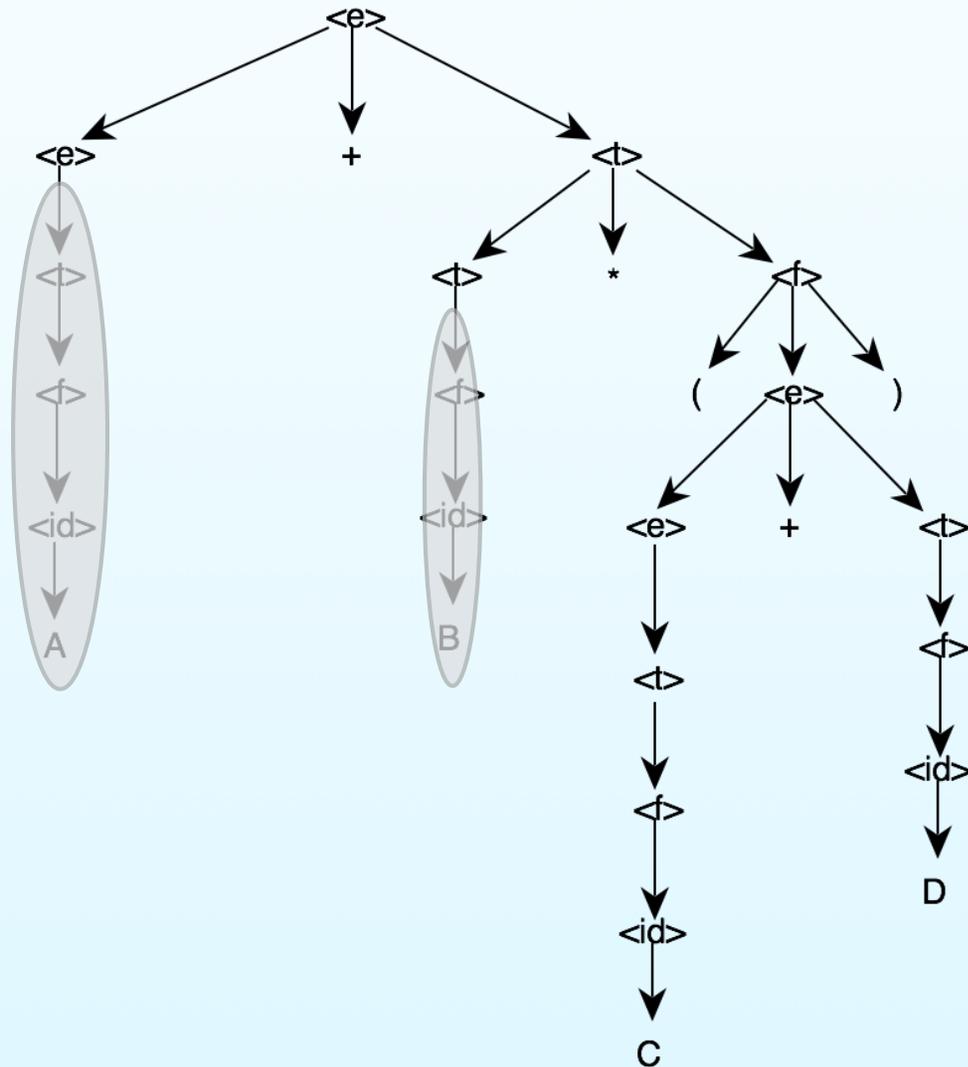


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
 ~~$\langle e \rangle + B * (C + D)$~~
 ~~$\langle t \rangle + B * (C + D)$~~
 ~~$\langle f \rangle + B * (C + D)$~~
 ~~$\langle id \rangle + B * (C + D)$~~
 ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

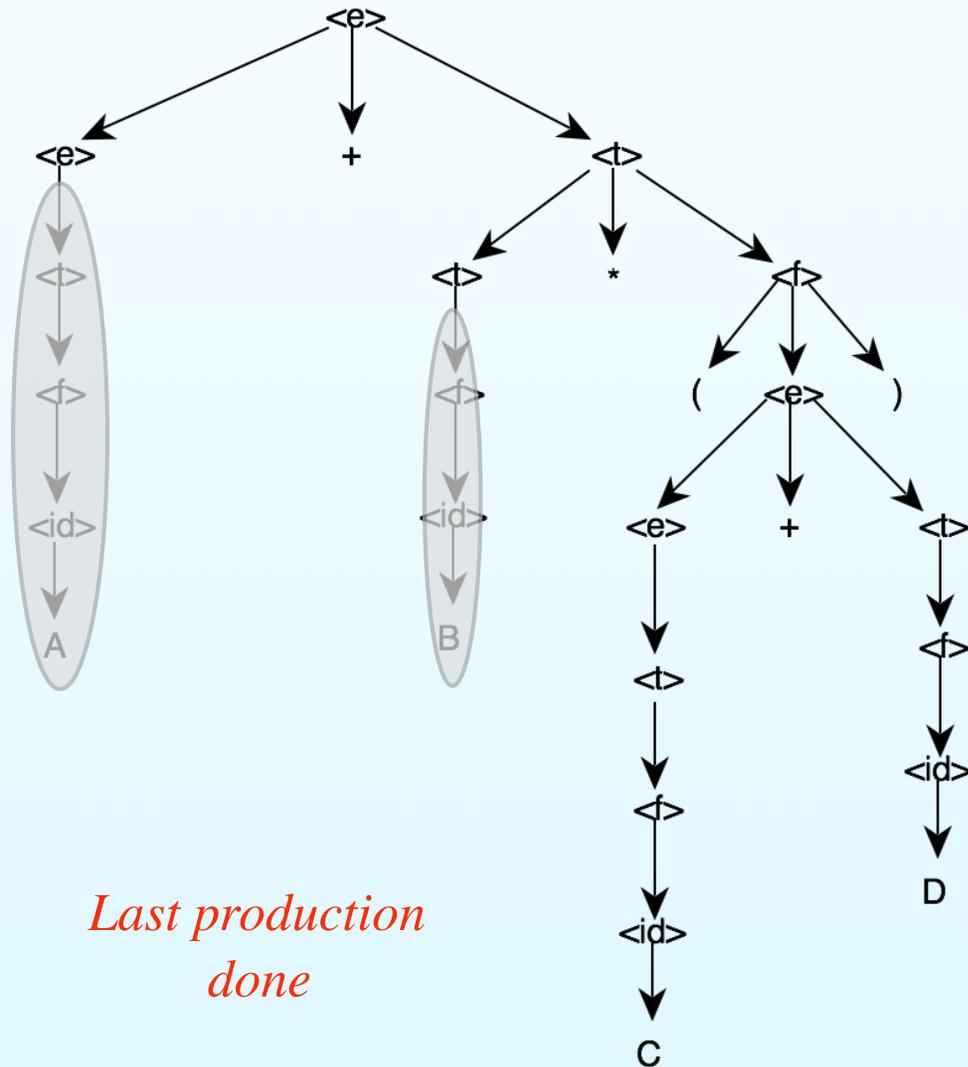


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
 $\langle e \rangle + \langle t \rangle * (C + D)$
 ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
 ~~$\langle e \rangle + B * (C + D)$~~
 ~~$\langle t \rangle + B * (C + D)$~~
 ~~$\langle f \rangle + B * (C + D)$~~
 ~~$\langle id \rangle + B * (C + D)$~~
 ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

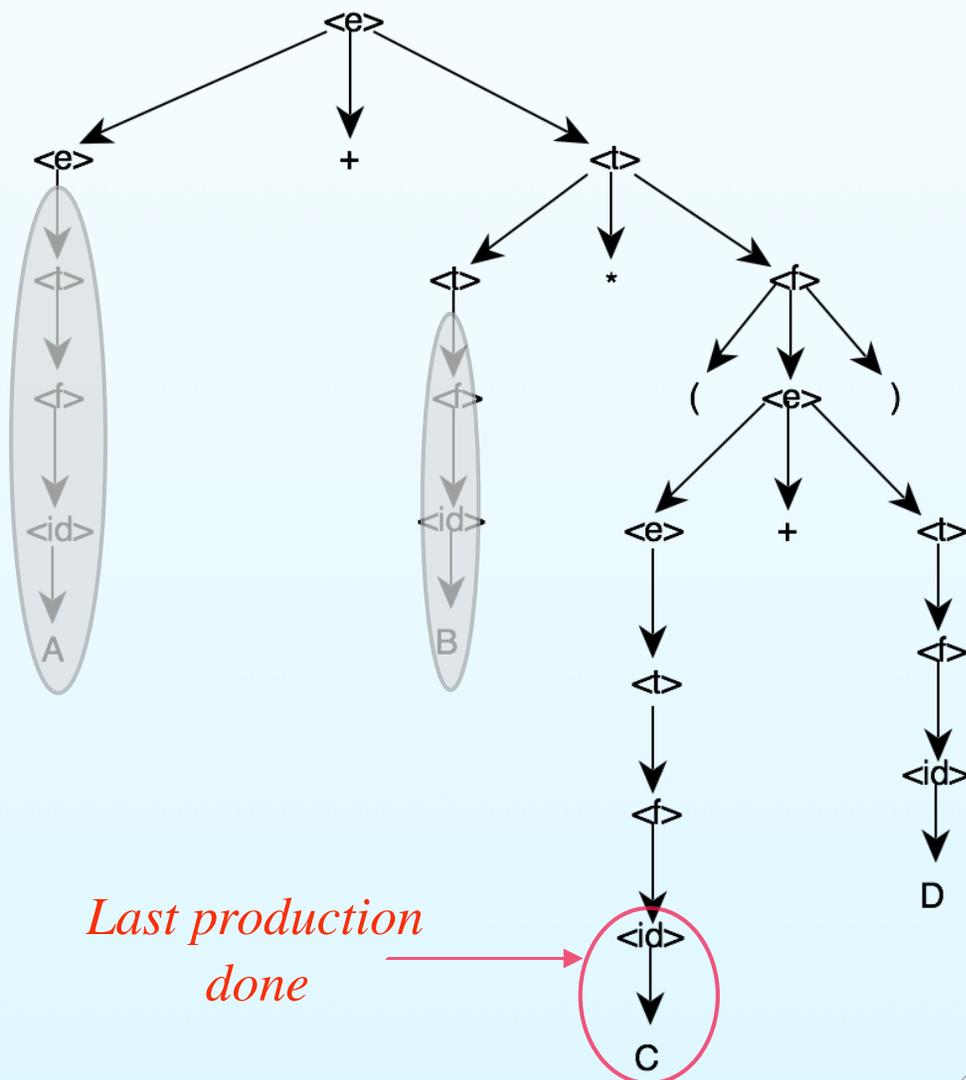


Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~



Last production done

Finding the handle

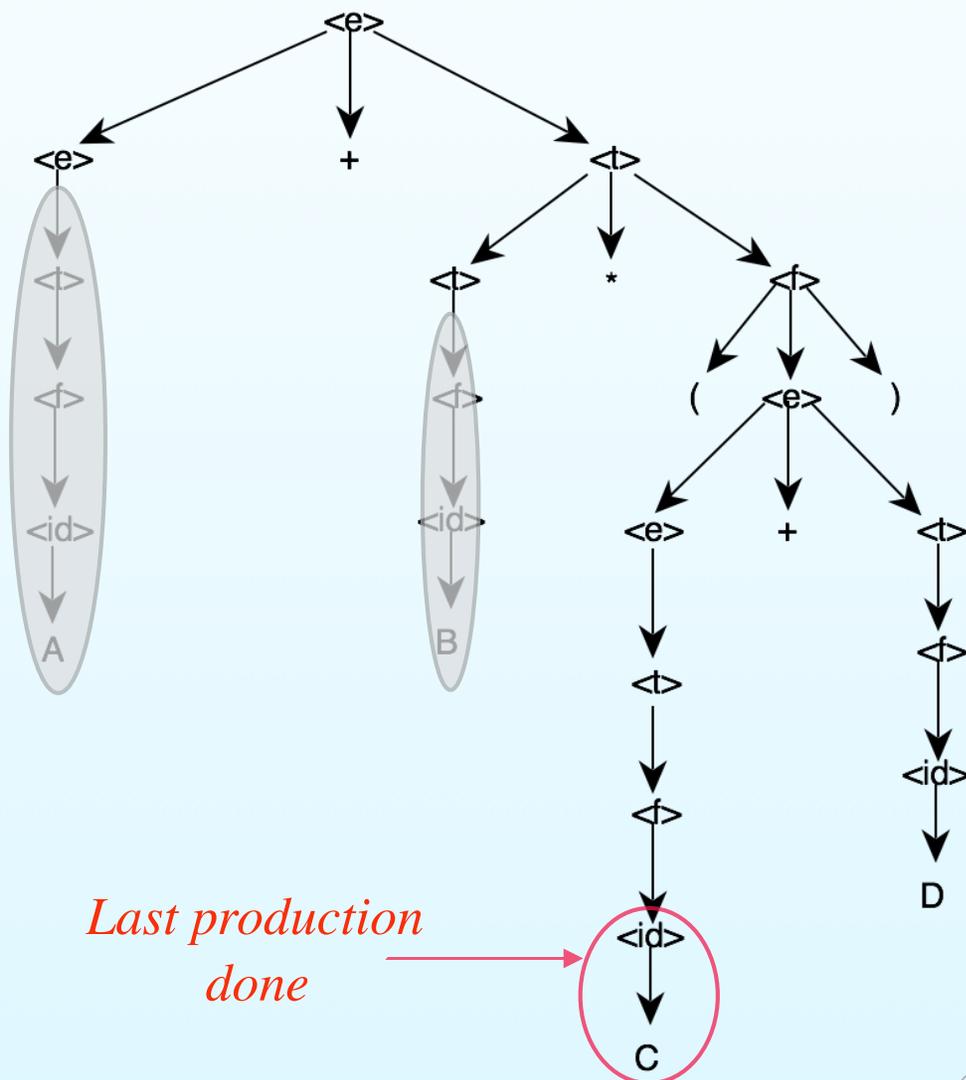
Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

Last production done



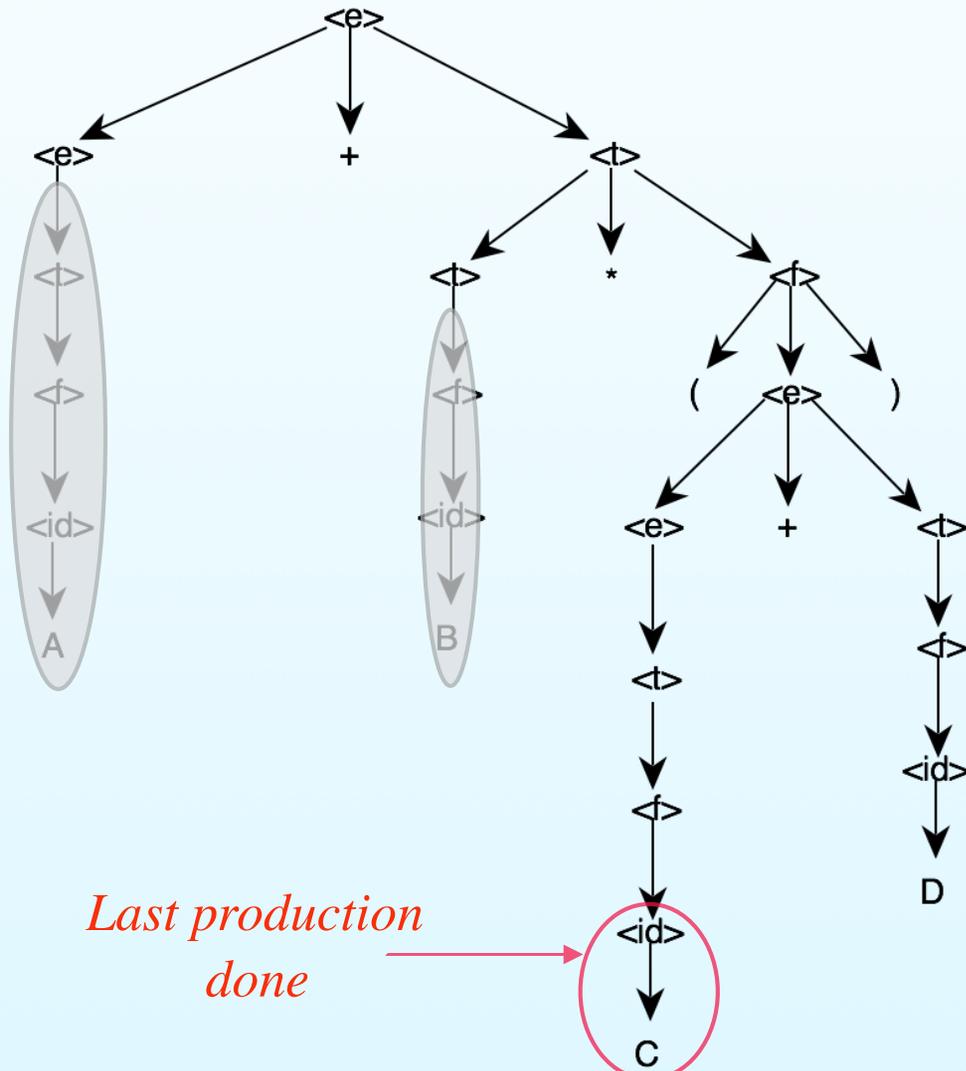
Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\mathbf{C} + D)$
- ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle



Last production done

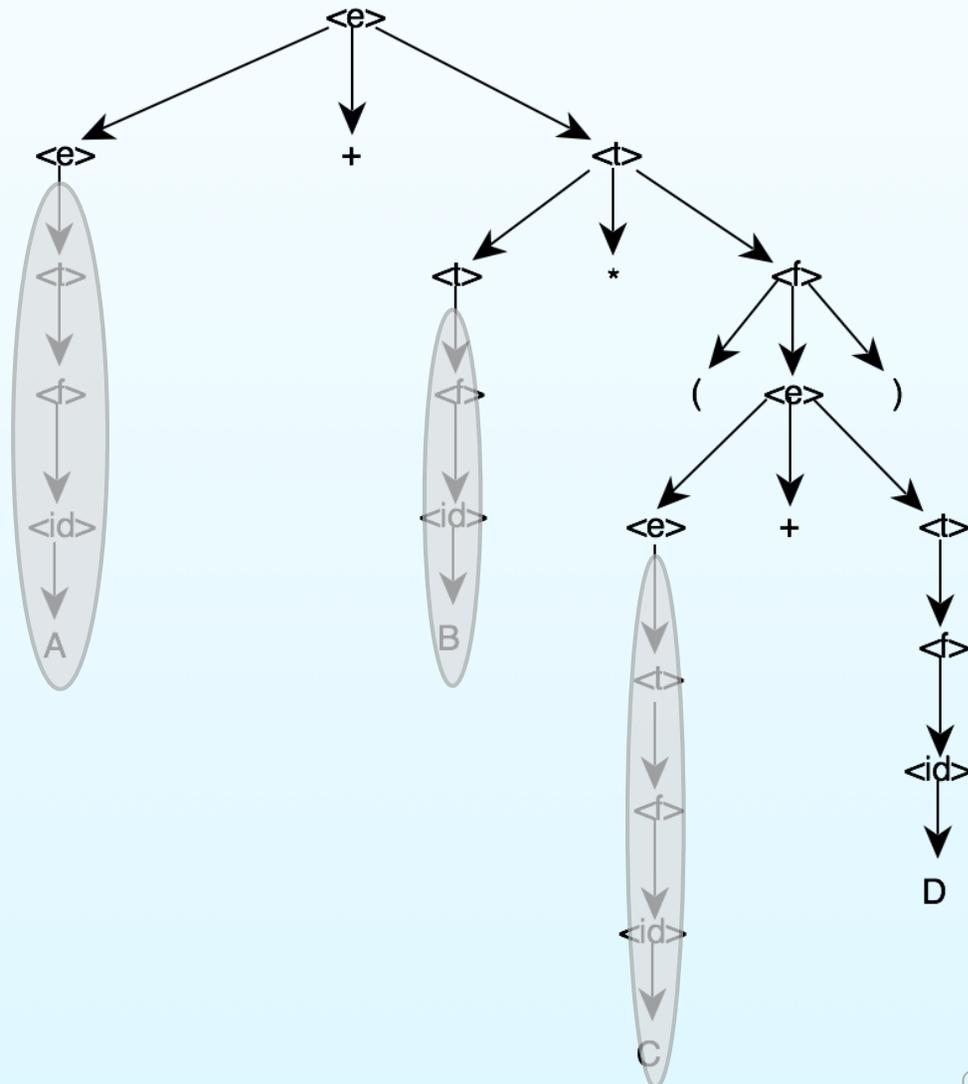
Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
 $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
 ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
 ~~$\langle e \rangle + B * (C + D)$~~
 ~~$\langle t \rangle + B * (C + D)$~~
 ~~$\langle f \rangle + B * (C + D)$~~
 ~~$\langle id \rangle + B * (C + D)$~~
 ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



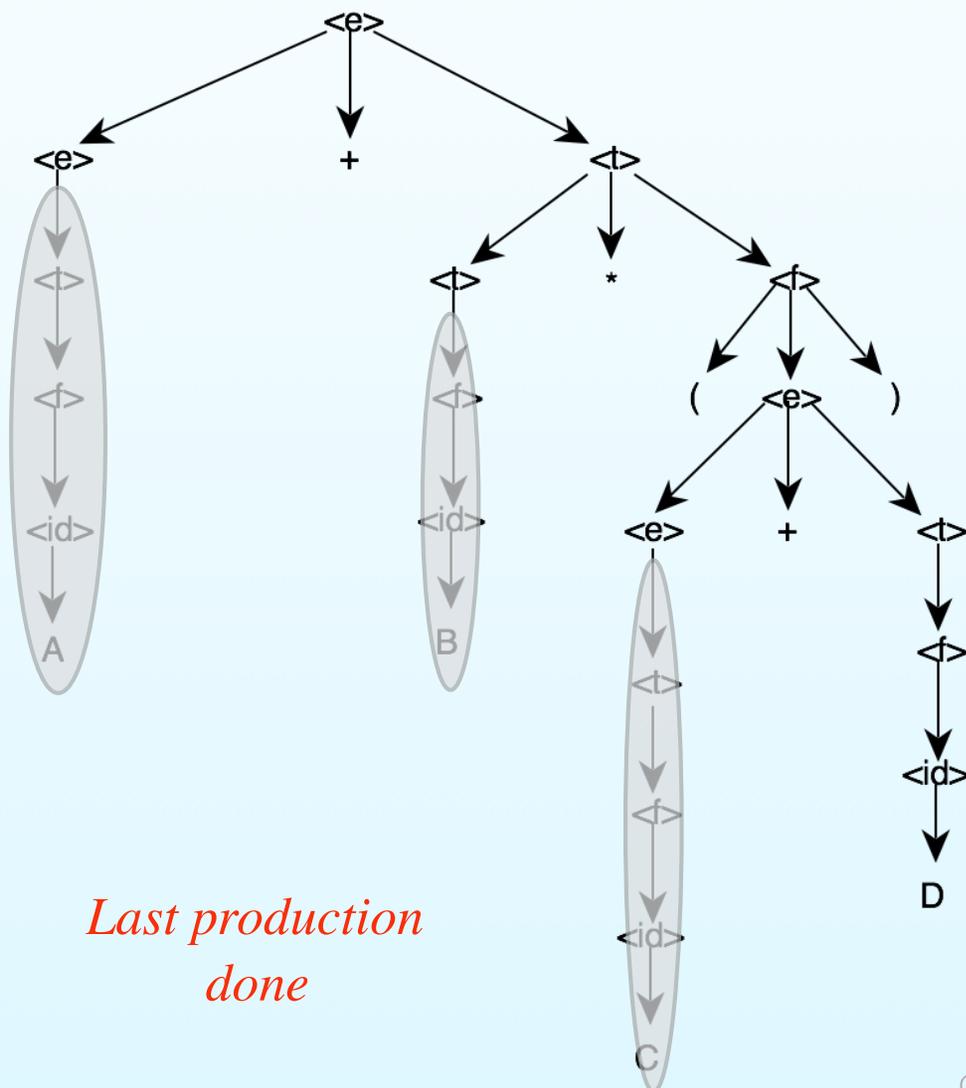
Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



Last production done

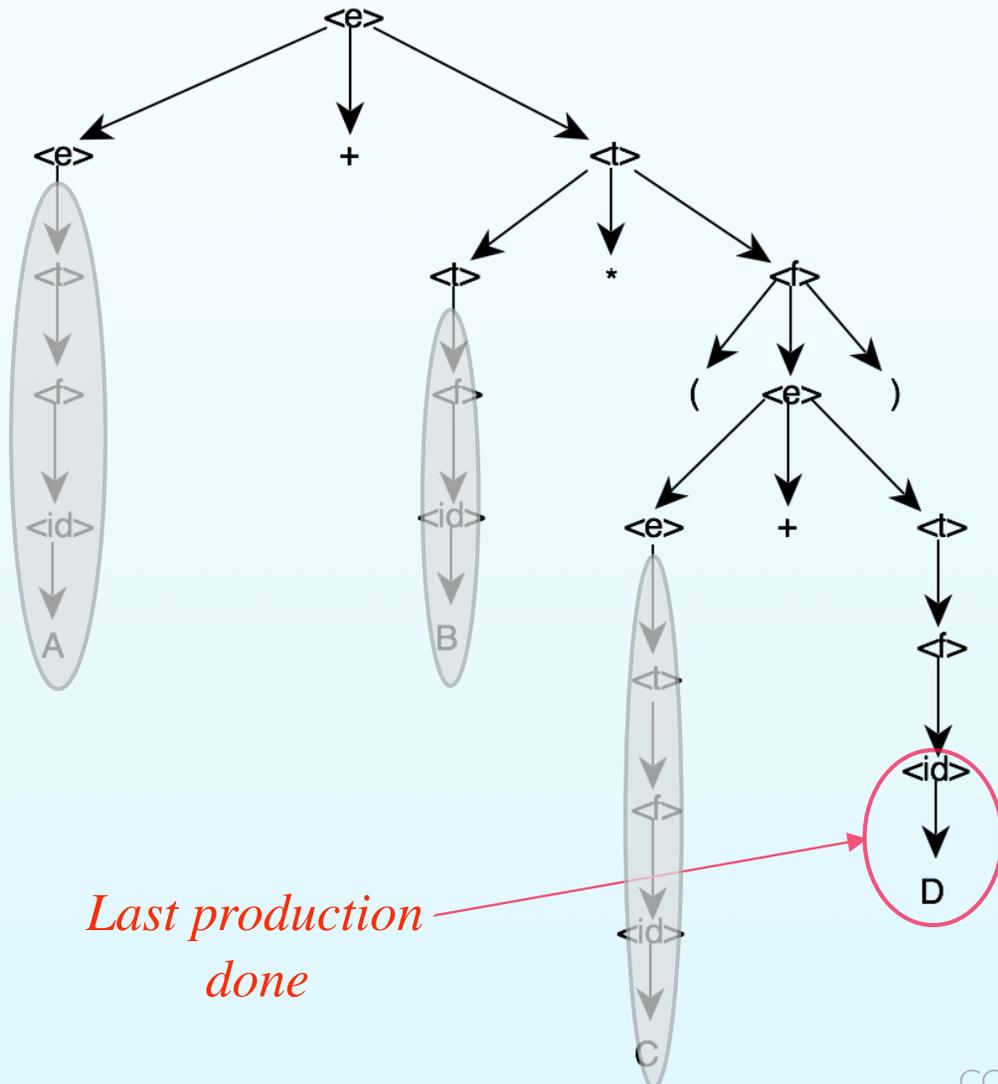
Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



Last production done

Finding the handle

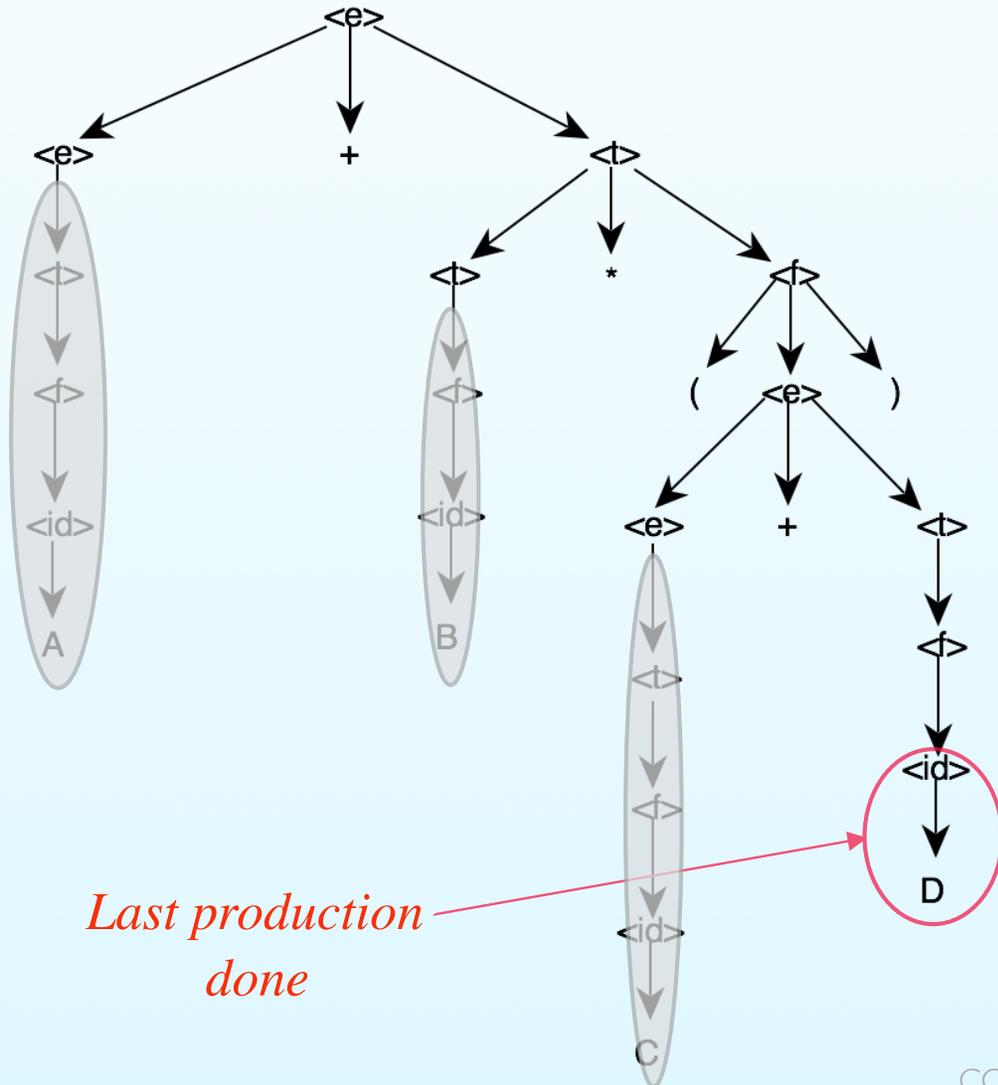
Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



Last production done

Finding the handle

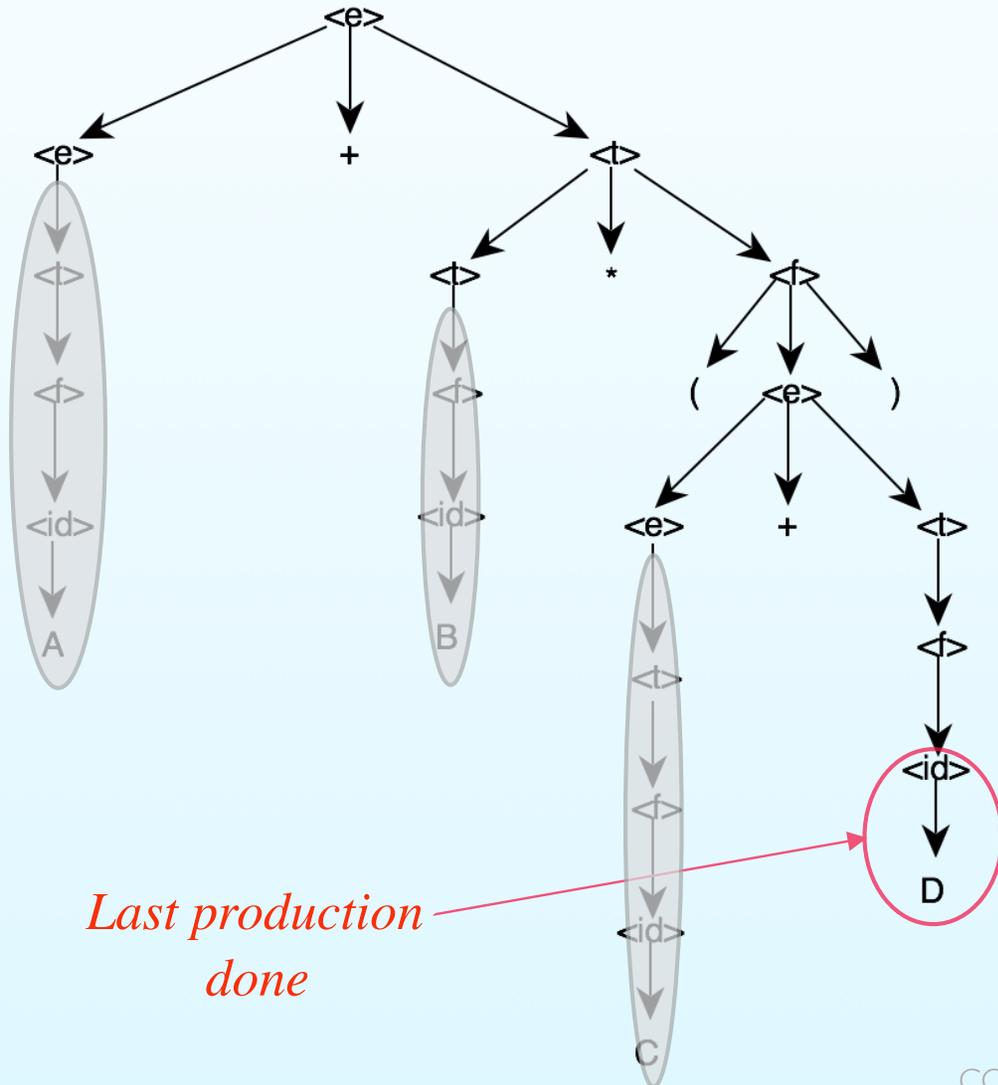
Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



Last production done

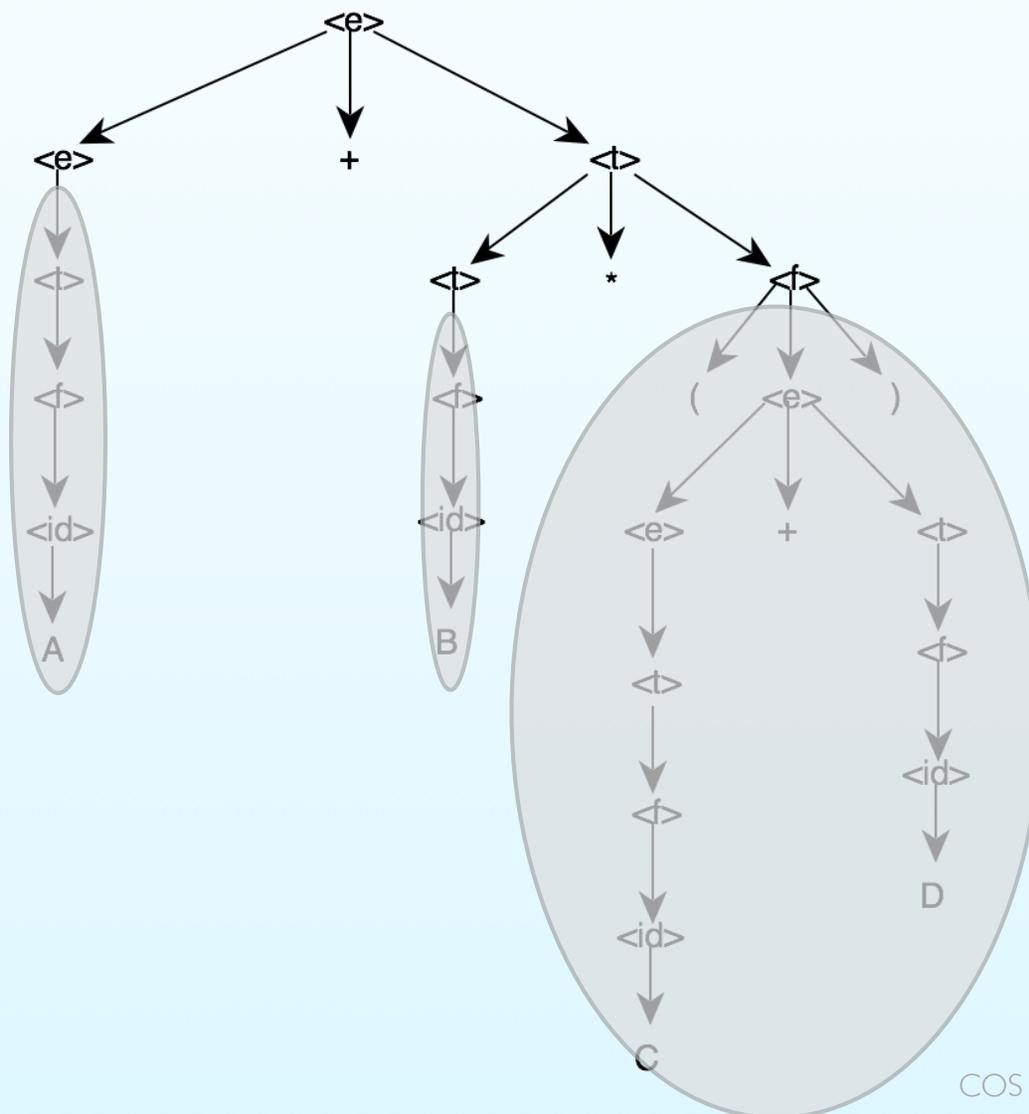
Finding the handle

Parse: A + B * (C + D)

- <e>
- <e> + <t>
- <e> + <t> * <f>
- ~~<e> + <t> * (<e>)~~
- ~~<e> + <t> * (<e> + <t>)~~
- ~~<e> + <t> * (<e> + <f>)~~
- ~~<e> + <t> * (<e> + <id>)~~
- ~~<e> + <t> * (<e> + D)~~
- ~~<e> + <t> * (<t> + D)~~
- ~~<e> + <t> * (<f> + D)~~
- ~~<e> + <t> * (<id> + D)~~
- ~~<e> + <t> * (C + D)~~
- ~~<e> + <f> * (C + D)~~
- ~~<e> + <id> * (C + D)~~
- ~~<e> + B * (C + D)~~
- ~~<t> + B * (C + D)~~
- ~~<f> + B * (C + D)~~
- ~~<id> + B * (C + D)~~
- A + B * (C + D)

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



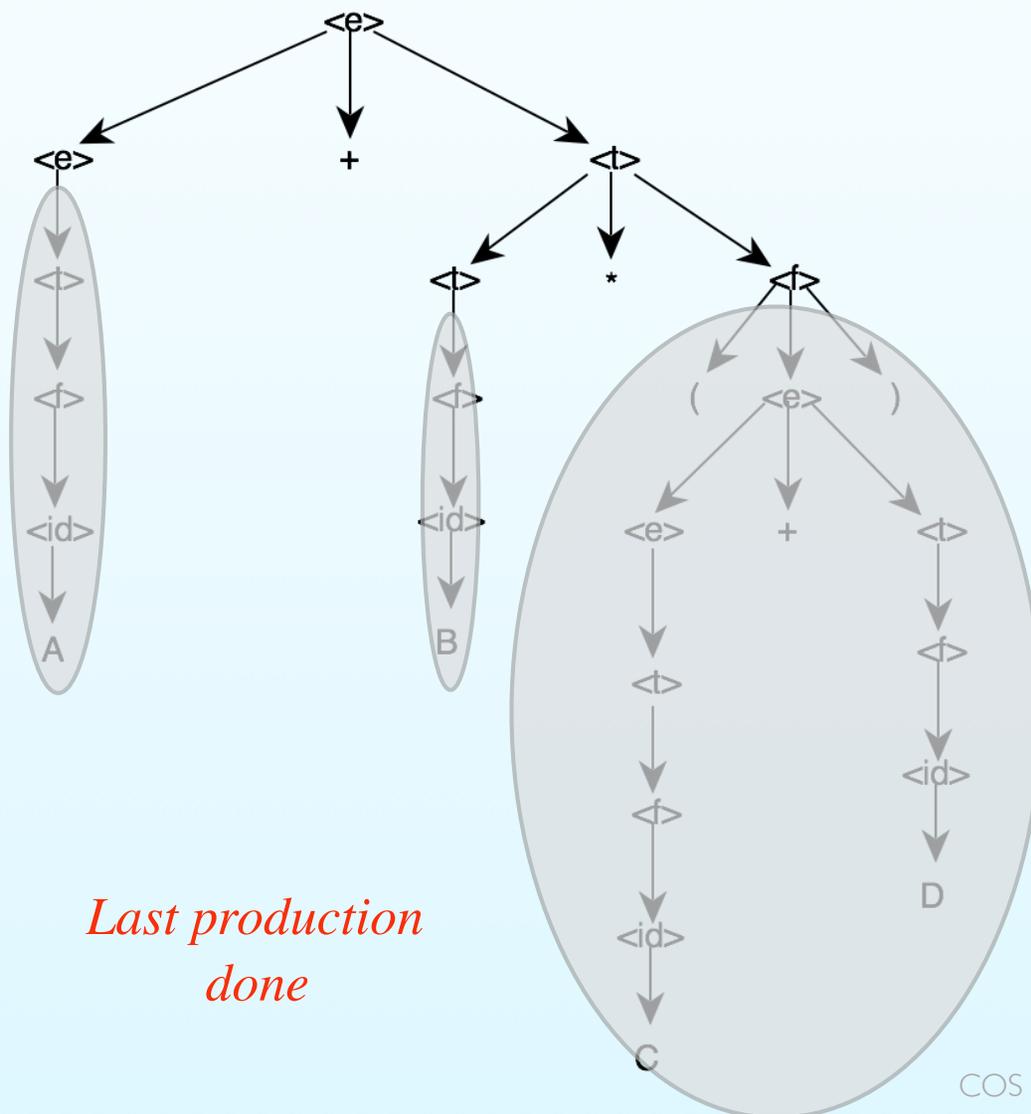
Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



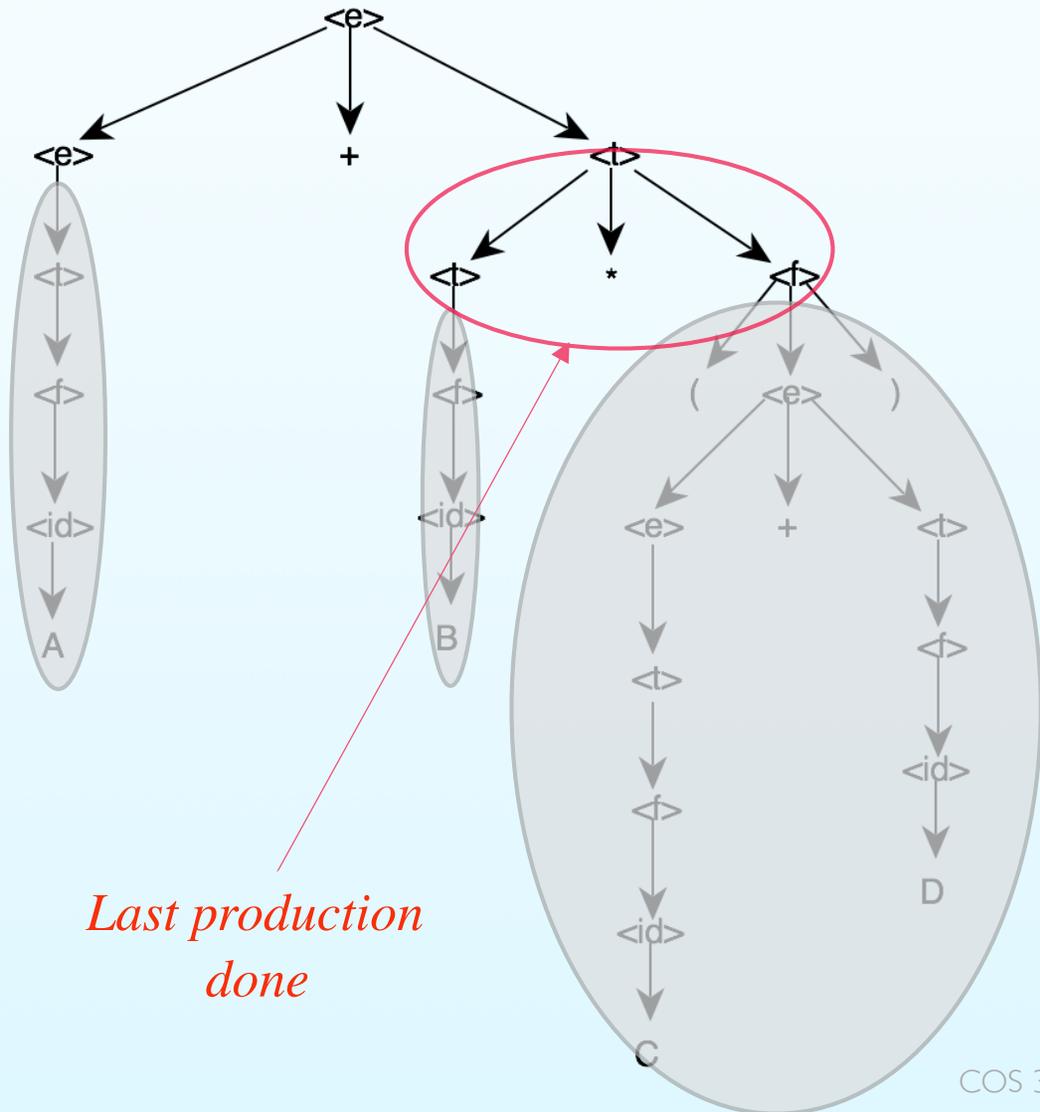
Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
 ~~$\langle e \rangle + B * (C + D)$~~
 ~~$\langle t \rangle + B * (C + D)$~~
 ~~$\langle f \rangle + B * (C + D)$~~
 ~~$\langle id \rangle + B * (C + D)$~~
 ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



Last production done

Finding the handle

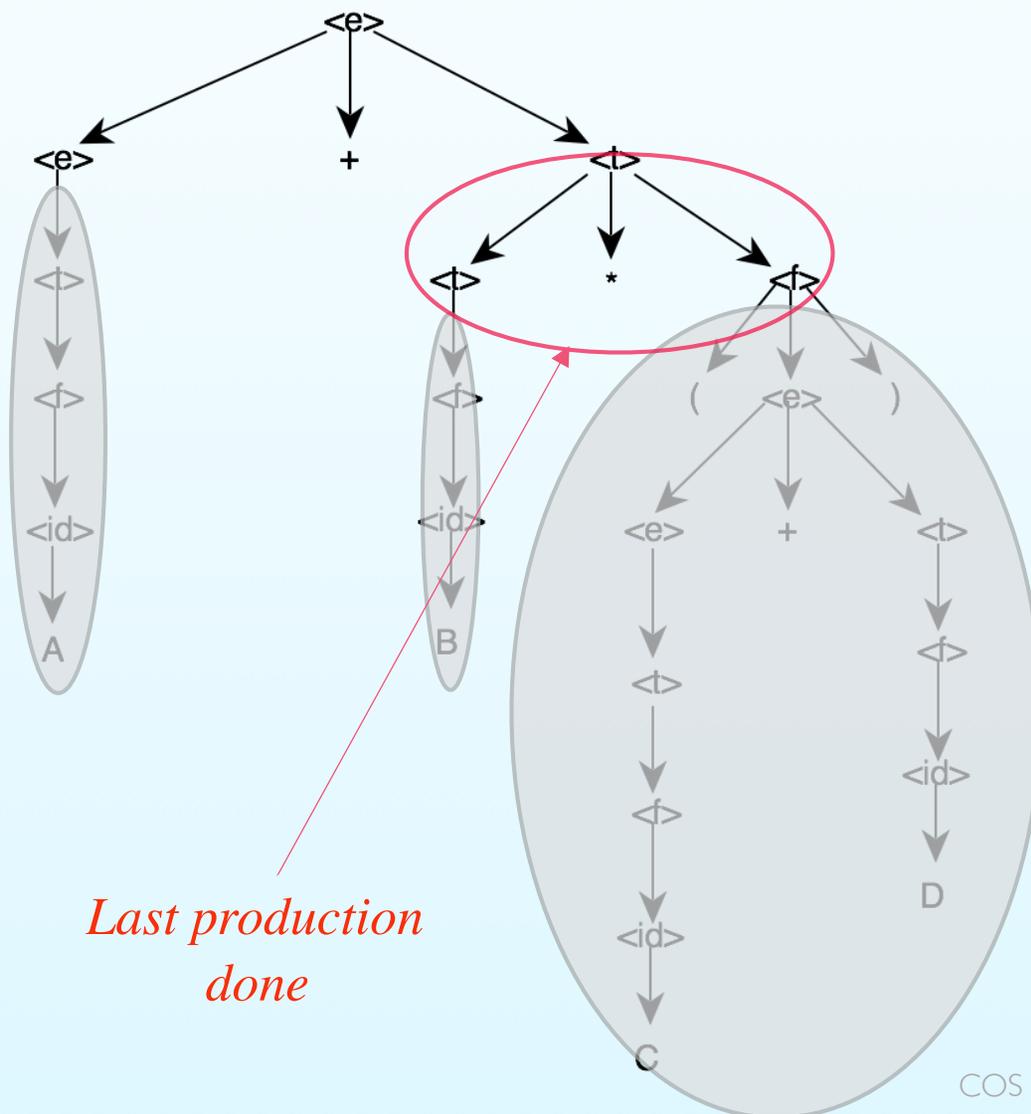
Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



Last production done

Finding the handle

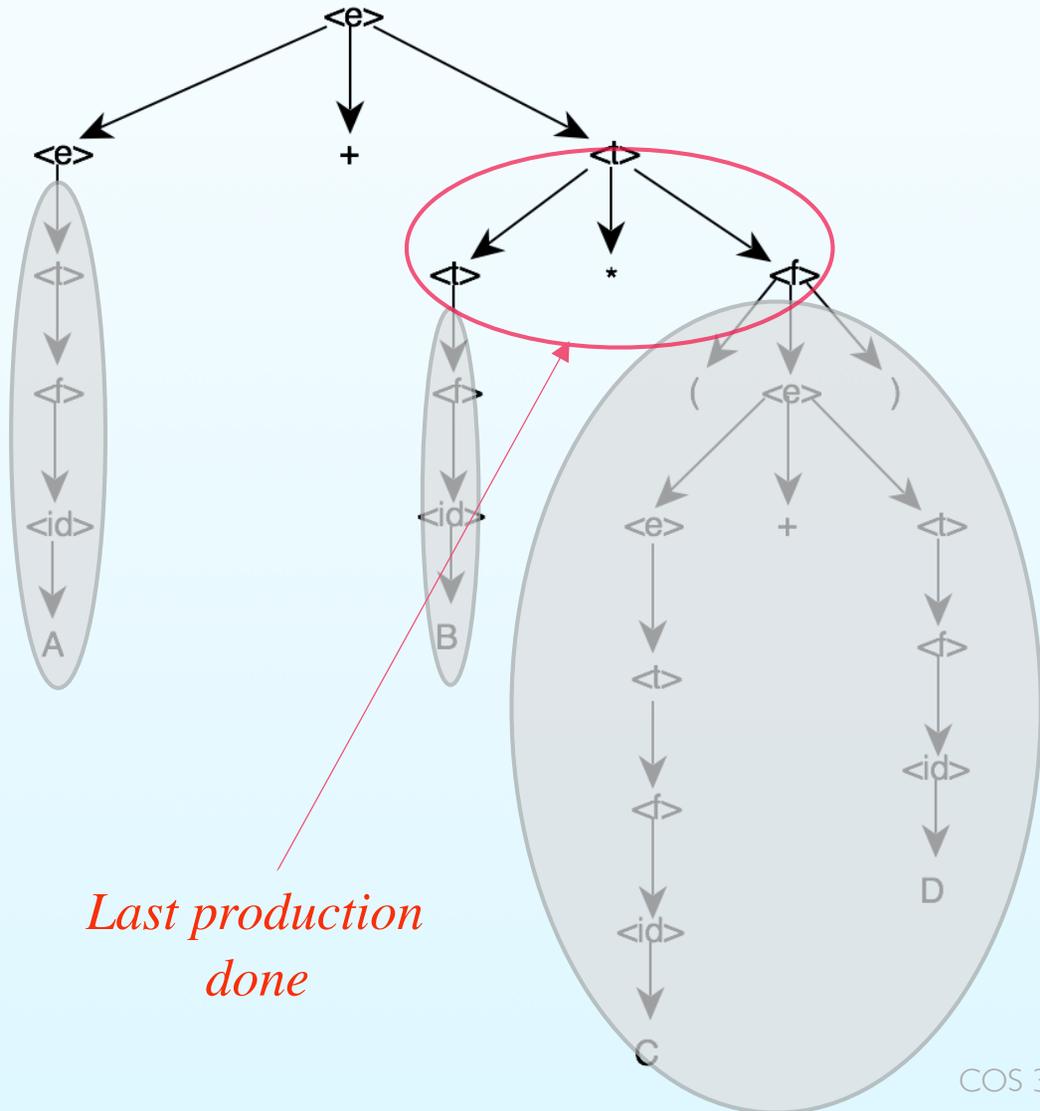
Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
 ~~$\langle e \rangle + B * (C + D)$~~
 ~~$\langle t \rangle + B * (C + D)$~~
 ~~$\langle f \rangle + B * (C + D)$~~
 ~~$\langle id \rangle + B * (C + D)$~~
 ~~$A + B * (C + D)$~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



Last production done

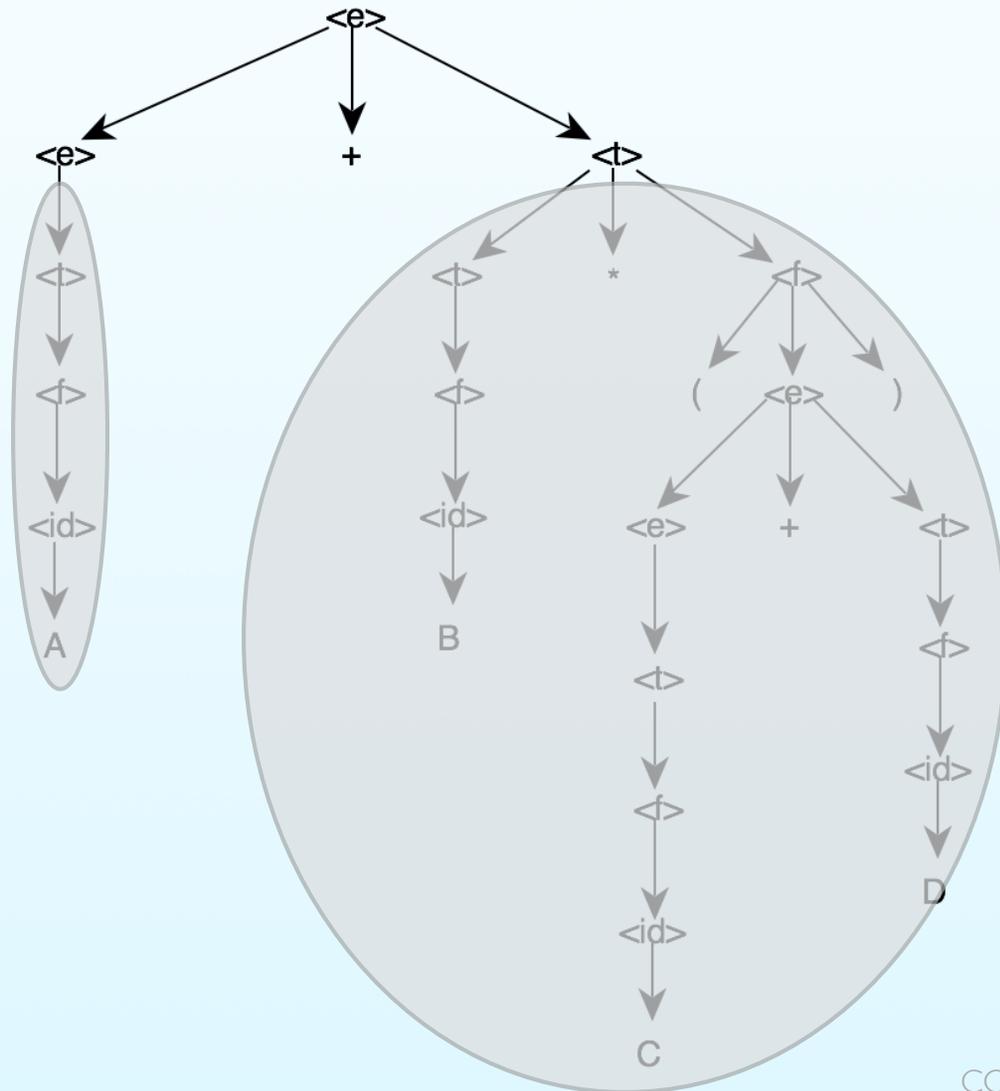
Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 ~~$\langle e \rangle + \langle t \rangle * \langle f \rangle$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
 ~~$\langle e \rangle + B * (C + D)$~~
 ~~$\langle t \rangle + B * (C + D)$~~
 ~~$\langle f \rangle + B * (C + D)$~~
 ~~$\langle id \rangle + B * (C + D)$~~
 ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



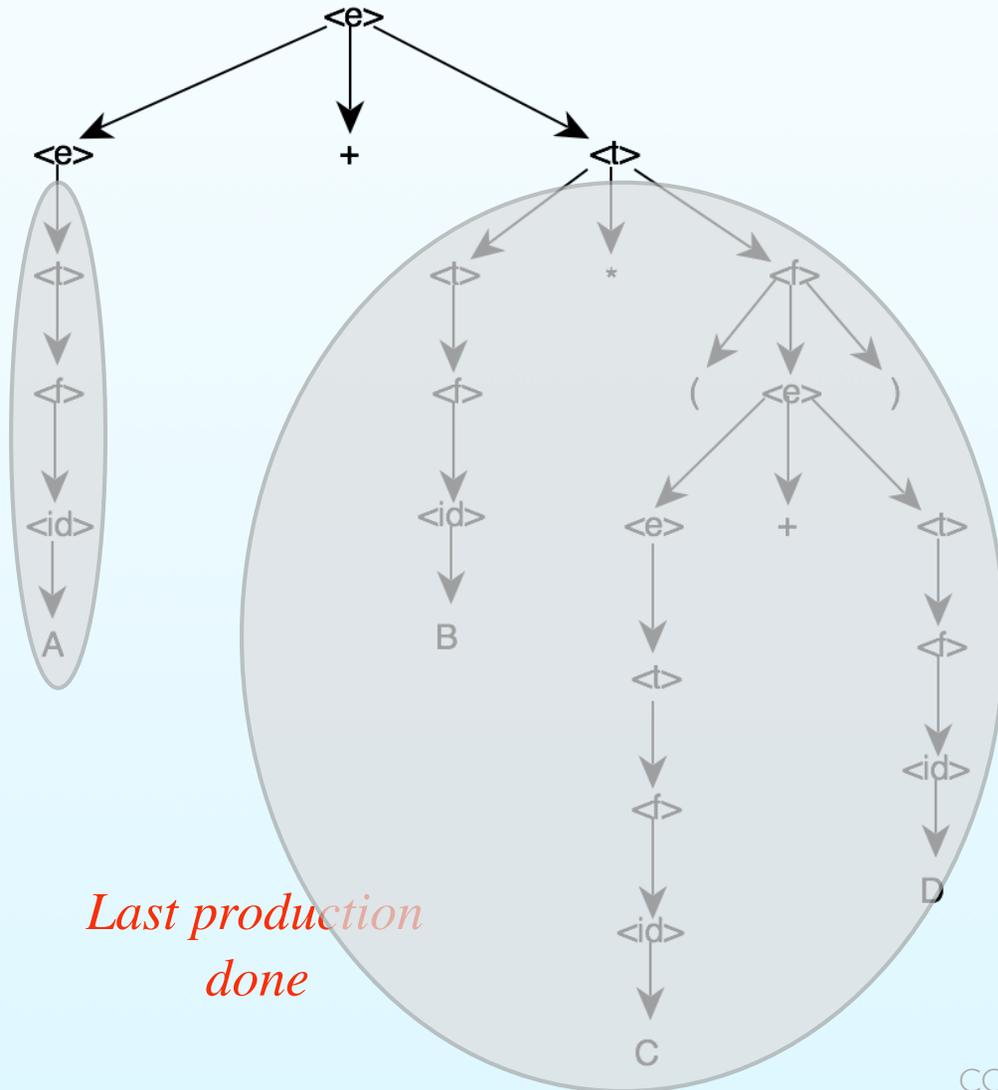
Finding the handle

Parse: $A + B * (C + D)$

$\langle e \rangle$
 $\langle e \rangle + \langle t \rangle$
 ~~$\langle e \rangle + \langle t \rangle * \langle f \rangle$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
 ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
 ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
 ~~$\langle e \rangle + B * (C + D)$~~
 ~~$\langle t \rangle + B * (C + D)$~~
 ~~$\langle f \rangle + B * (C + D)$~~
 ~~$\langle id \rangle + B * (C + D)$~~
 ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



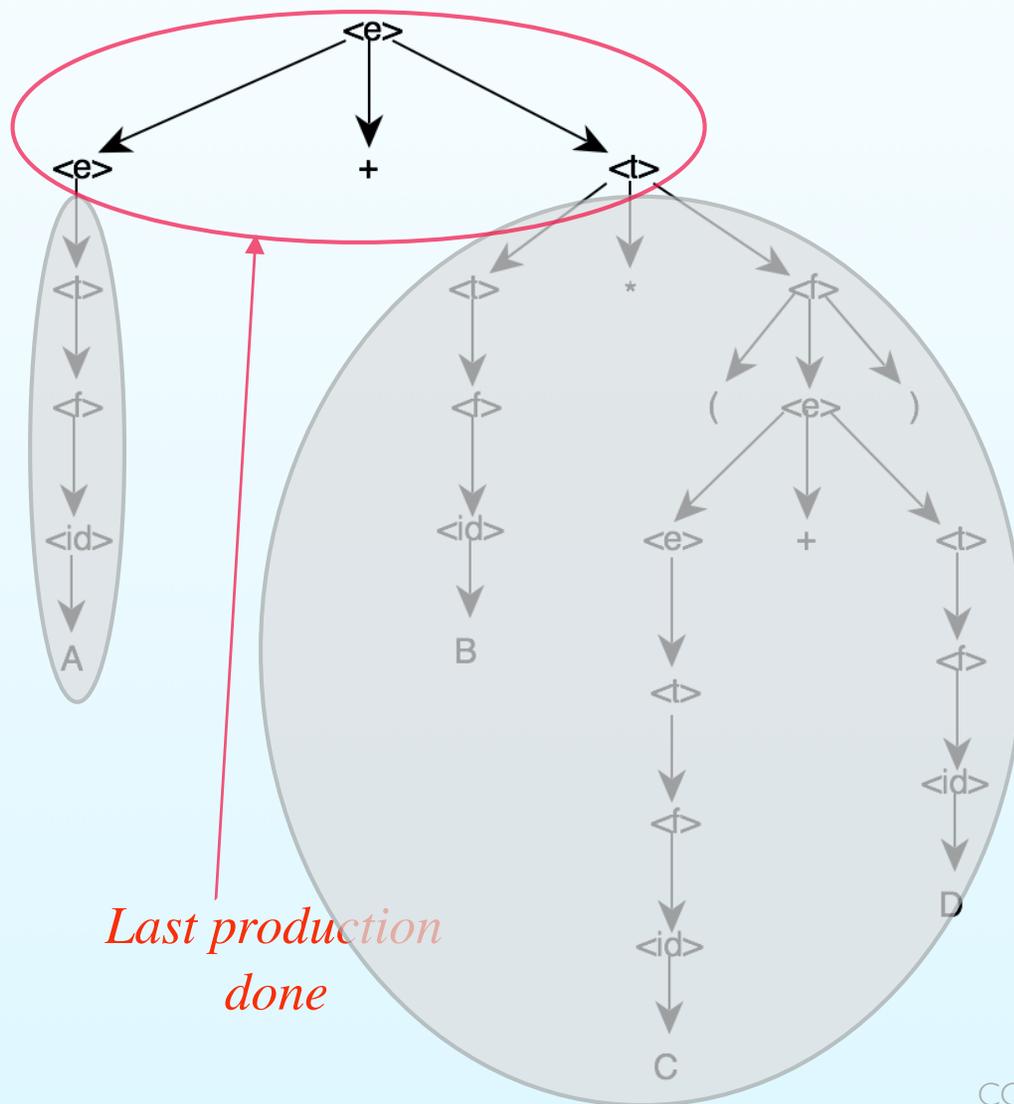
Finding the handle

Parse: A + B * (C + D)

~~<e>~~
~~<e> + <t>~~
~~<e> + <t> * <f>~~
~~<e> + <t> * (<e>)~~
~~<e> + <t> * (<e> + <t>)~~
~~<e> + <t> * (<e> + <f>)~~
~~<e> + <t> * (<e> + <id>)~~
~~<e> + <t> * (<e> + D)~~
~~<e> + <t> * (<t> + D)~~
~~<e> + <t> * (<f> + D)~~
~~<e> + <t> * (<id> + D)~~
~~<e> + <t> * (C + D)~~
~~<e> + <f> * (C + D)~~
~~<e> + <id> * (C + D)~~
~~<e> + B * (C + D)~~
~~<t> + B * (C + D)~~
~~<f> + B * (C + D)~~
~~<id> + B * (C + D)~~
 A + B * (C + D)

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



Last production done

Finding the handle

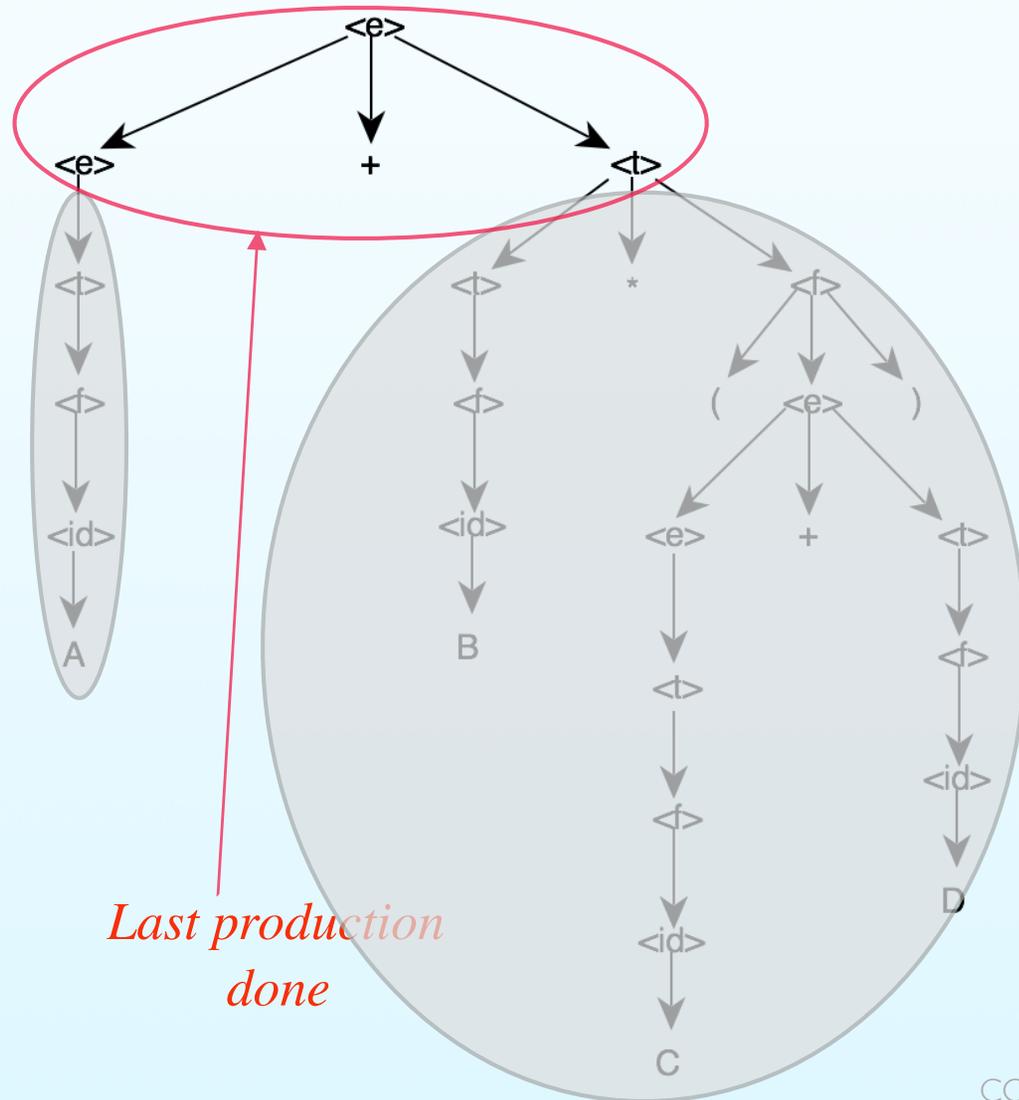
Parse: A + B * (C + D)

~~<e>~~
~~<e> + <t>~~
~~<e> + <t> * <f>~~
~~<e> + <t> * (<e>)~~
~~<e> + <t> * (<e> + <t>)~~
~~<e> + <t> * (<e> + <f>)~~
~~<e> + <t> * (<e> + <id>)~~
~~<e> + <t> * (<e> + D)~~
~~<e> + <t> * (<t> + D)~~
~~<e> + <t> * (<f> + D)~~
~~<e> + <t> * (<id> + D)~~
~~<e> + <t> * (C + D)~~
~~<e> + <f> * (C + D)~~
~~<e> + <id> * (C + D)~~
~~<e> + B * (C + D)~~
~~<t> + B * (C + D)~~
~~<f> + B * (C + D)~~
~~<id> + B * (C + D)~~
A + B * (C + D)

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Skip a few steps...



Last production done

Finding the handle

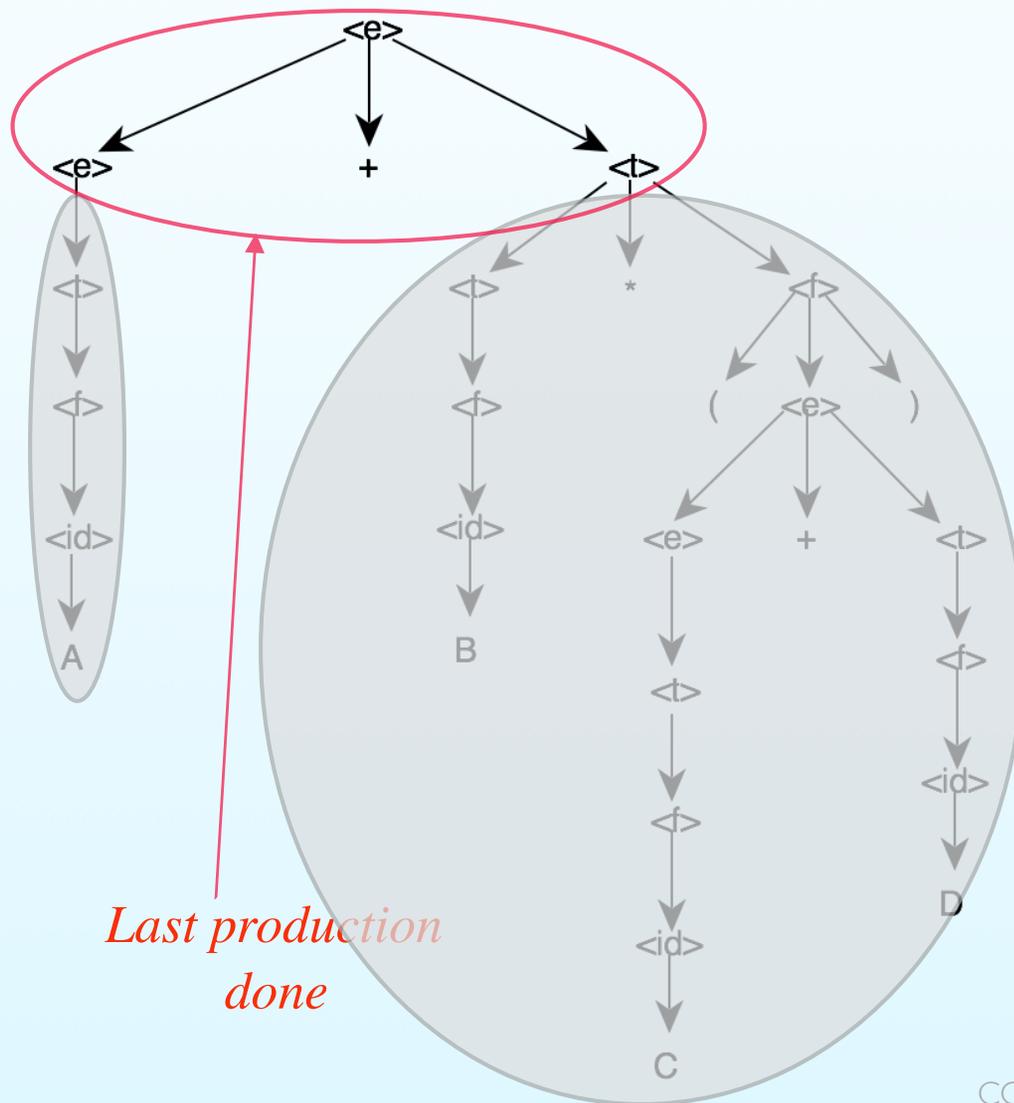
Parse: A + B * (C + D)

- <e>
- <e> + <t>**
- ~~<e> + <t> * <f>~~
- ~~<e> + <t> * (<e>)~~
- ~~<e> + <t> * (<e> + <t>)~~
- ~~<e> + <t> * (<e> + <f>)~~
- ~~<e> + <t> * (<e> + <id>)~~
- ~~<e> + <t> * (<e> + D)~~
- ~~<e> + <t> * (<t> + D)~~
- ~~<e> + <t> * (<f> + D)~~
- ~~<e> + <t> * (<id> + D)~~
- ~~<e> + <t> * (C + D)~~
- ~~<e> + <f> * (C + D)~~
- ~~<e> + <id> * (C + D)~~
- ~~<e> + B * (C + D)~~
- ~~<t> + B * (C + D)~~
- ~~<f> + B * (C + D)~~
- ~~<id> + B * (C + D)~~
- A + B * (C + D)

Handle

<e> → <e> + <t> | <t>
 <t> → <t> * <f> | <f>
 <f> → (<e>) | <id>

Skip a few steps...



Last production done

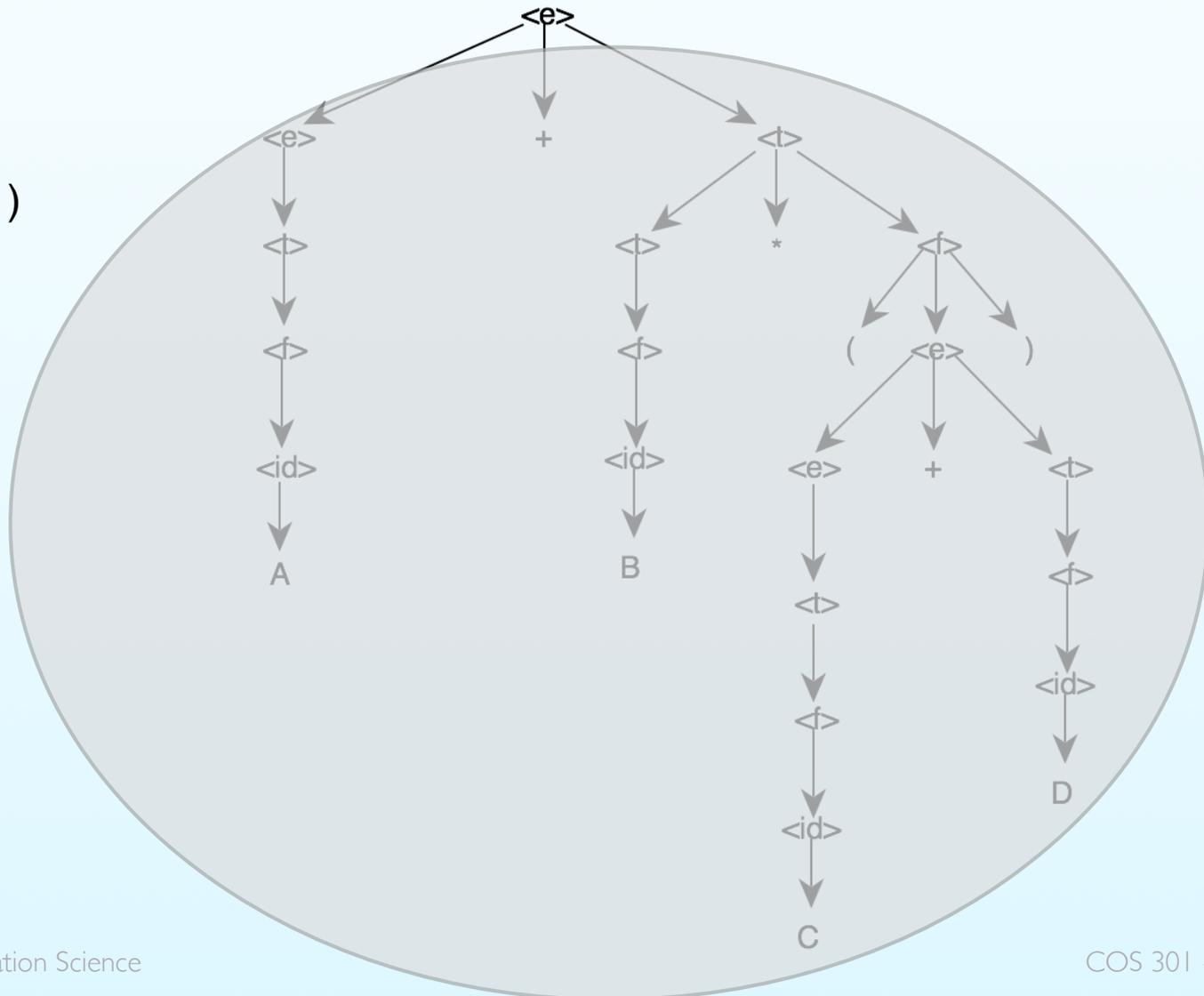
Finding the handle

Parse: A + B * (C + D)

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Done

- ~~$\langle e \rangle$~~
- ~~$\langle e \rangle + \langle t \rangle$~~
- ~~$\langle e \rangle + \langle t \rangle * \langle f \rangle$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~



Finding the handle

- What was common to all handles?

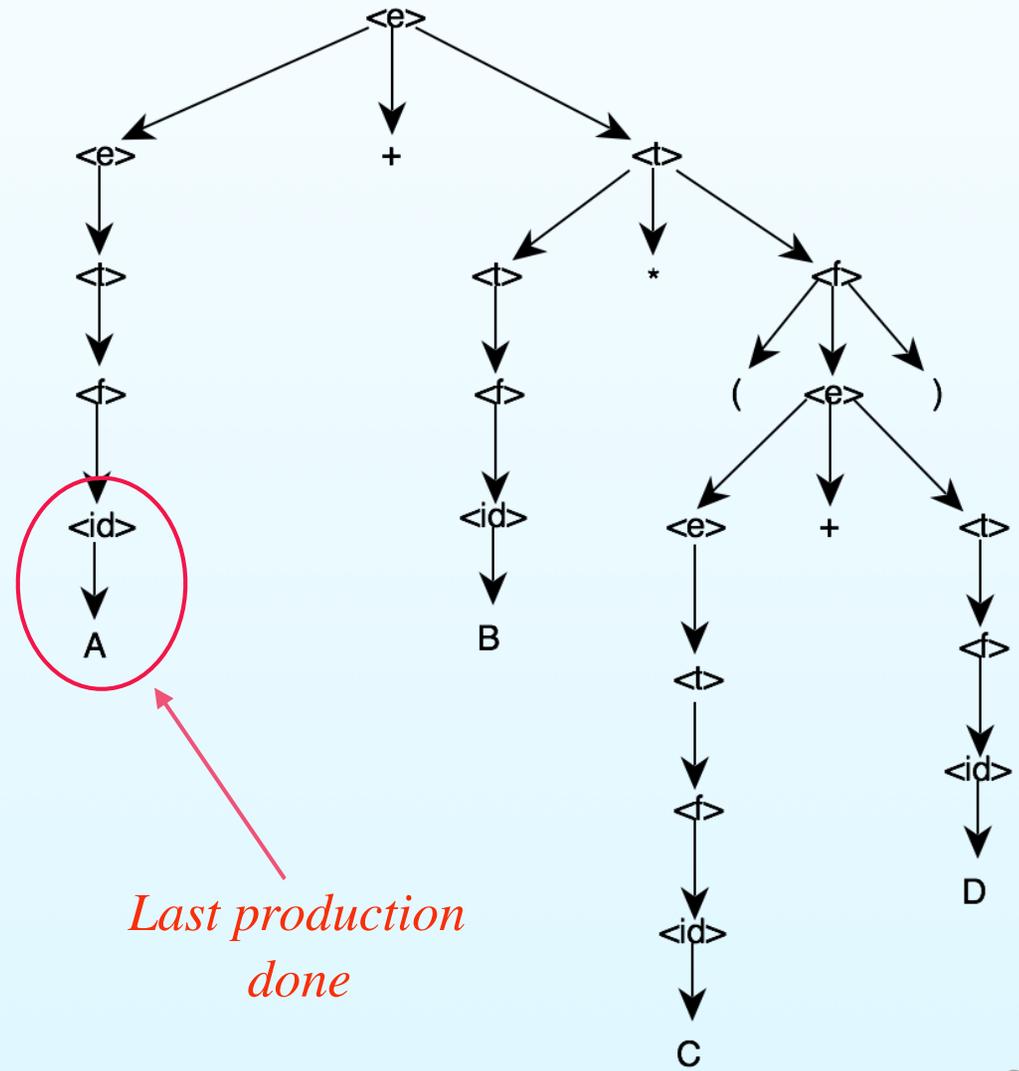
Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- $A + B * (C + D)$

Handle → $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



Last production done

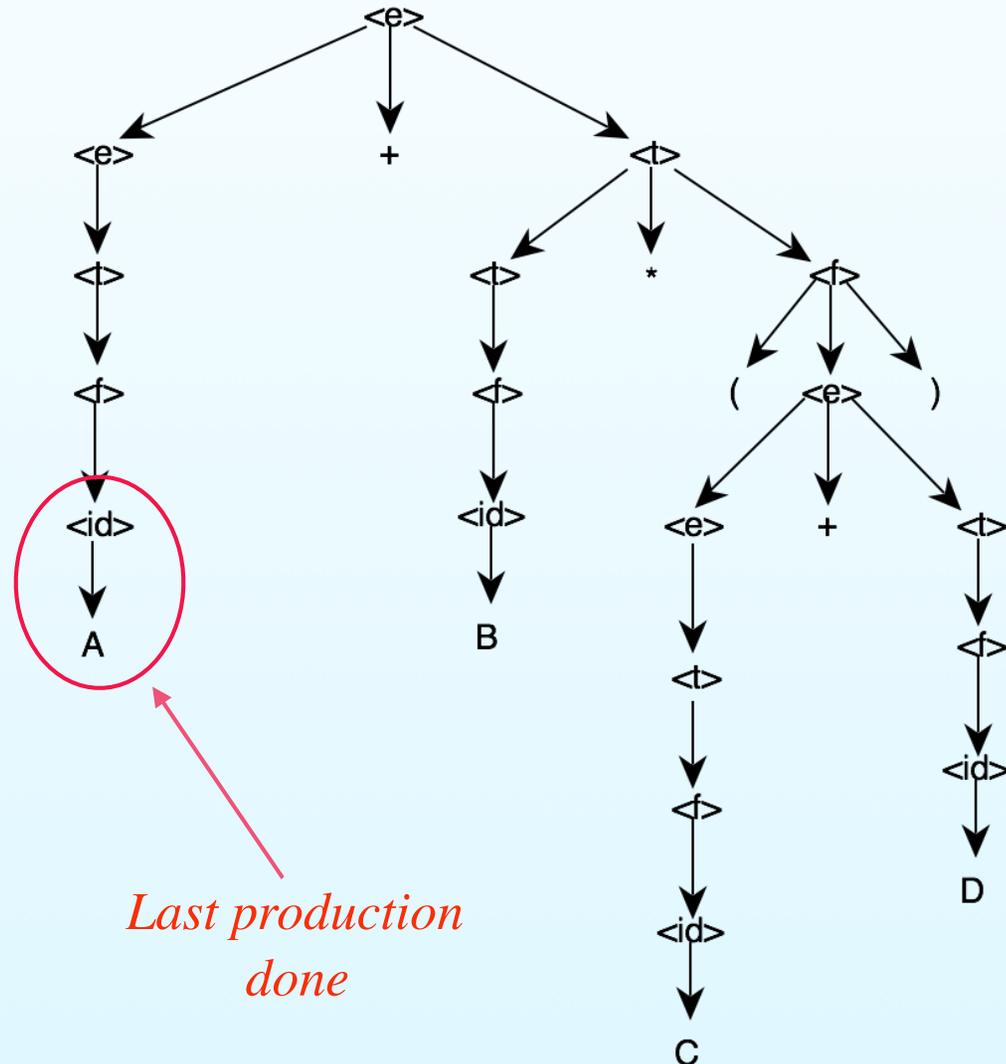
Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Hypothesis: Handle is leftmost RHS?



Handle

Last production done

Finding the handle

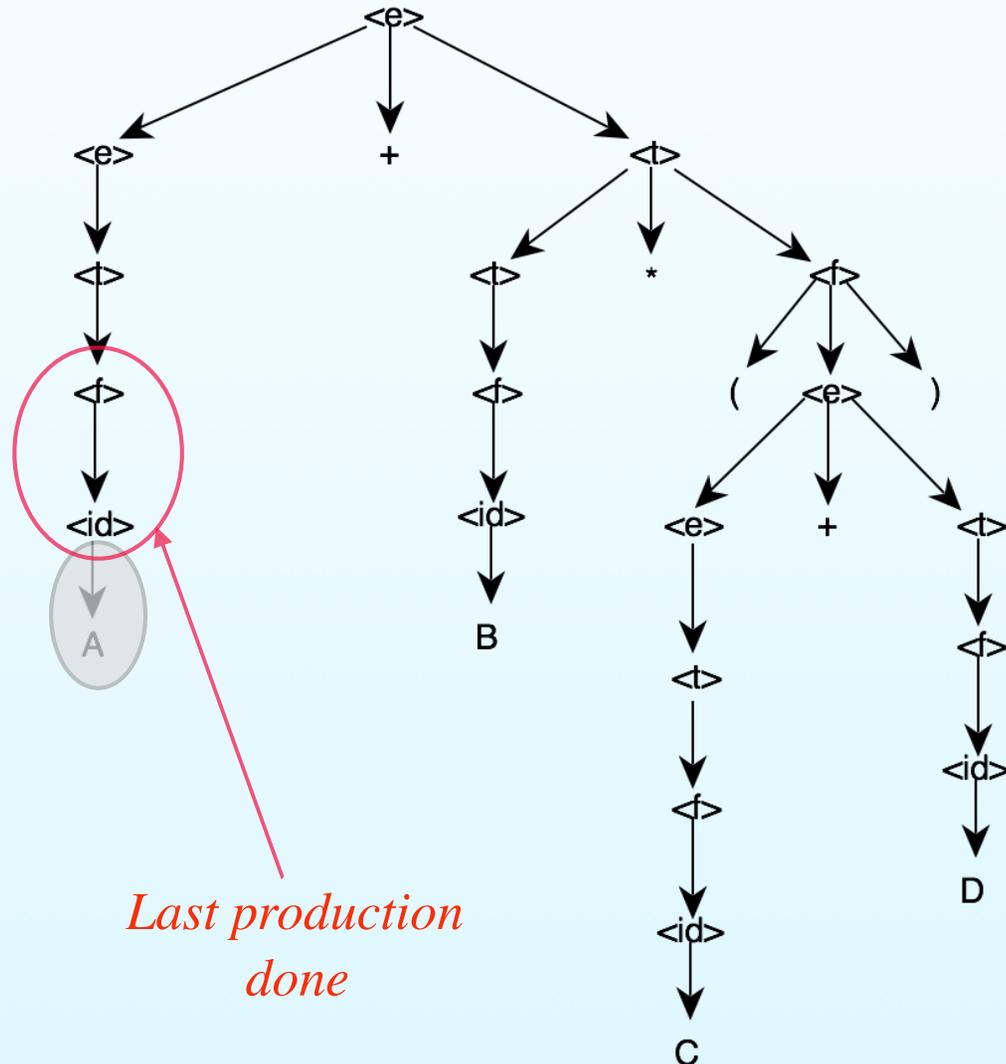
Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- ~~$A + B * (C + D)$~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Hypothesis: Handle is leftmost RHS?



Last production done

Finding the handle

Parse: $A + B * (C + D)$

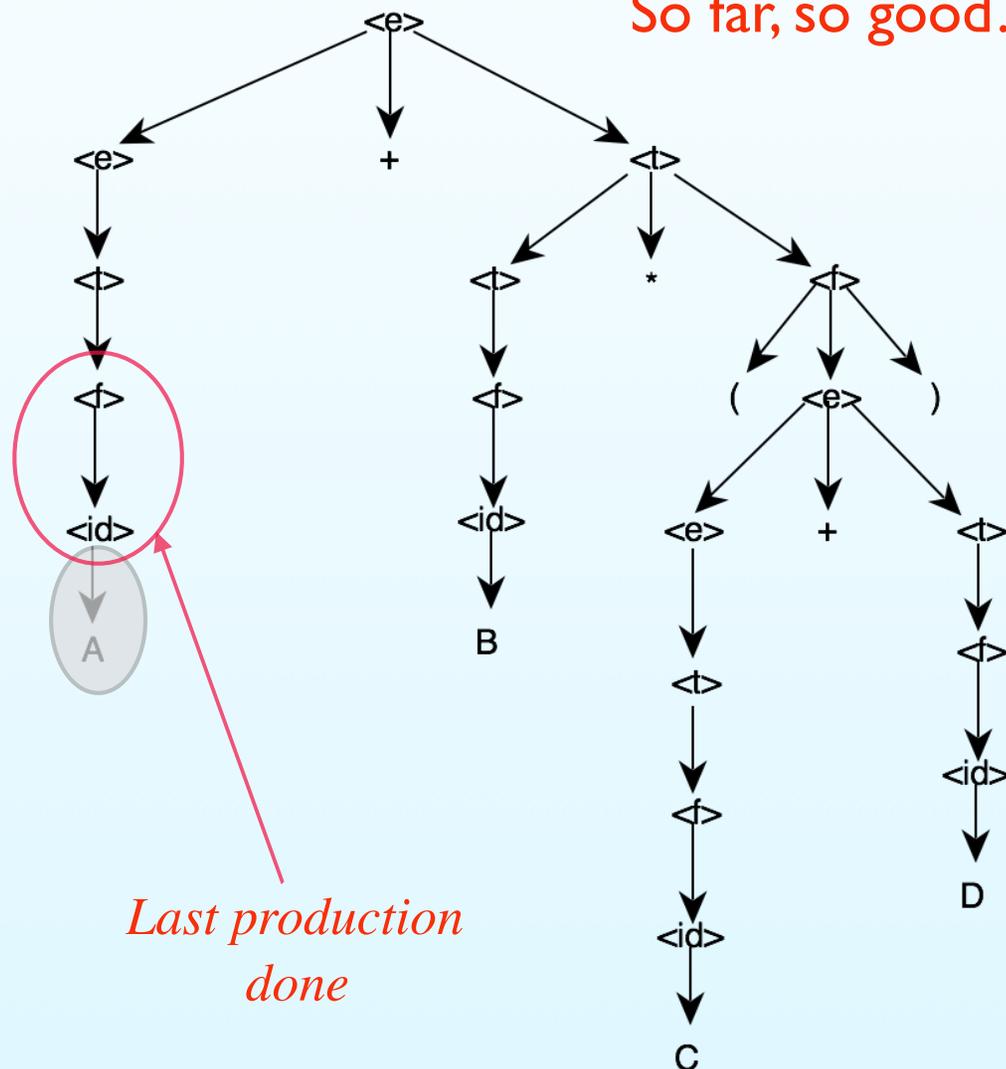
- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
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- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- $\langle t \rangle + B * (C + D)$
- $\langle f \rangle + B * (C + D)$
- $\langle id \rangle + B * (C + D)$
- ~~$A + B * (C + D)$~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Hypothesis: Handle is leftmost RHS?

So far, so good...



Last production done

Finding the handle

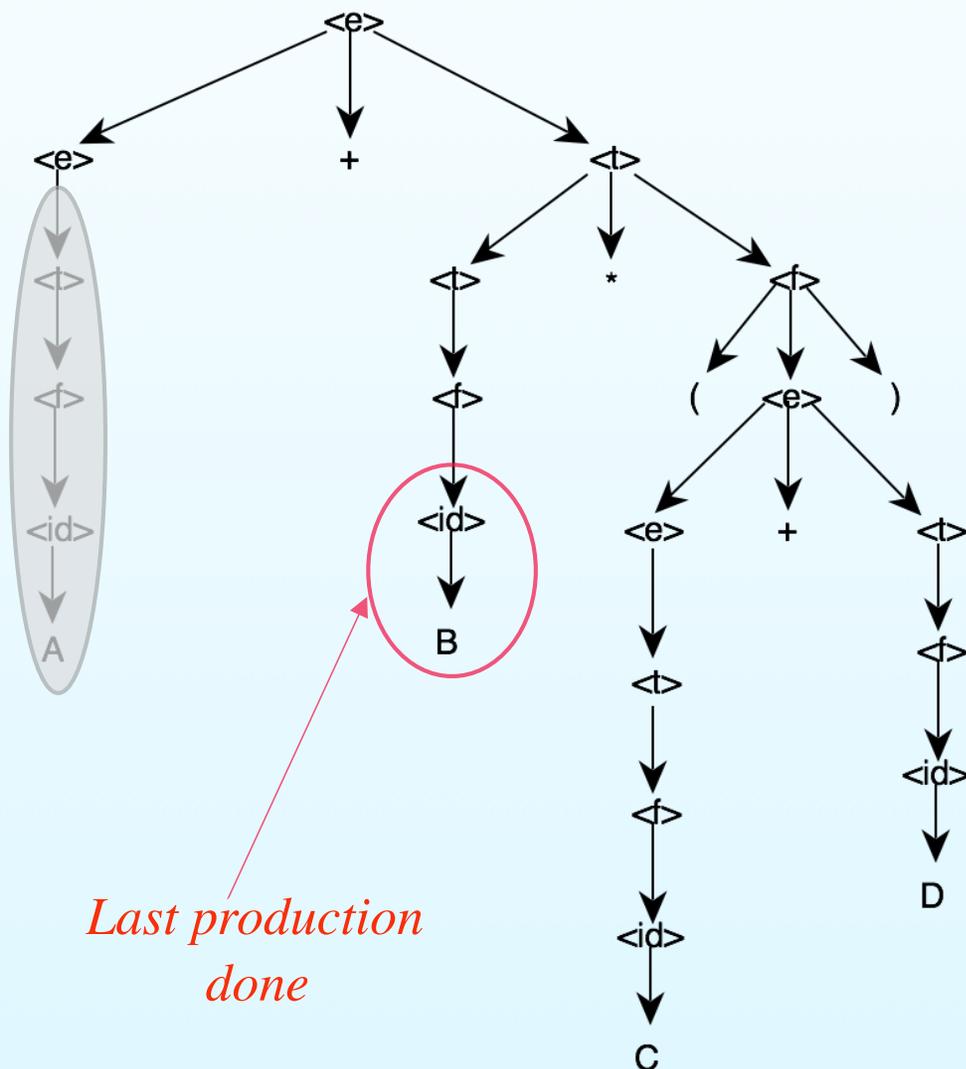
Parse: A + B * (C + D)

- <e>
- <e> + <t>
- <e> + <t> * <f>
- <e> + <t> * (<e>)
- <e> + <t> * (<e> + <t>)
- <e> + <t> * (<e> + <f>)
- <e> + <t> * (<e> + <id>)
- <e> + <t> * (<e> + D)
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- <e> + <t> * (<f> + D)
- <e> + <t> * (<id> + D)
- <e> + <t> * (C + D)
- <e> + <f> * (C + D)
- <e> + <id> * (C + D)
- <e> + **B** * (C + D)
- ~~<t> + B * (C + D)~~
- ~~<f> + B * (C + D)~~
- ~~<id> + B * (C + D)~~
- ~~A + B * (C + D)~~

Handle

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
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Parse: $A + B * (C + D)$

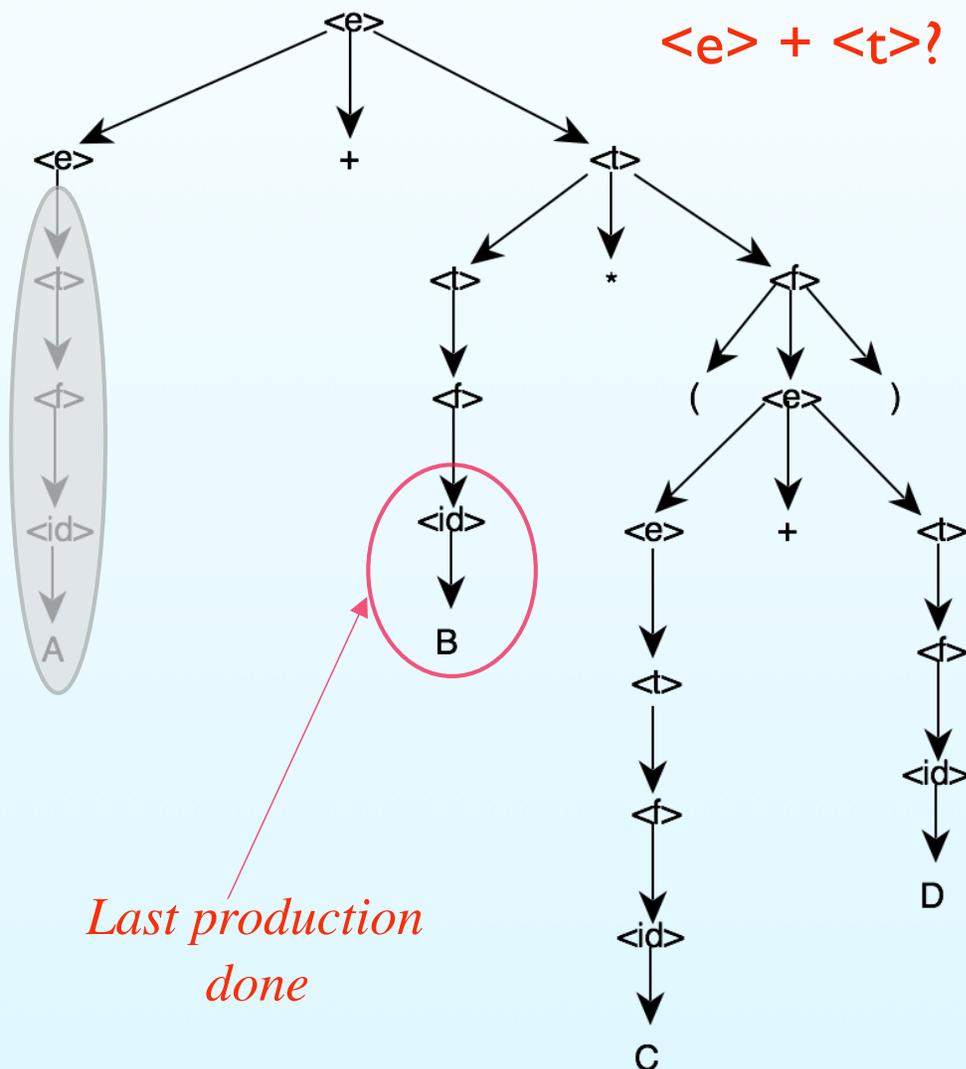
- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
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- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
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Hypothesis: Handle is leftmost RHS?

No...wouldn't this be $\langle e \rangle + \langle t \rangle$?



Last production done

Finding the handle

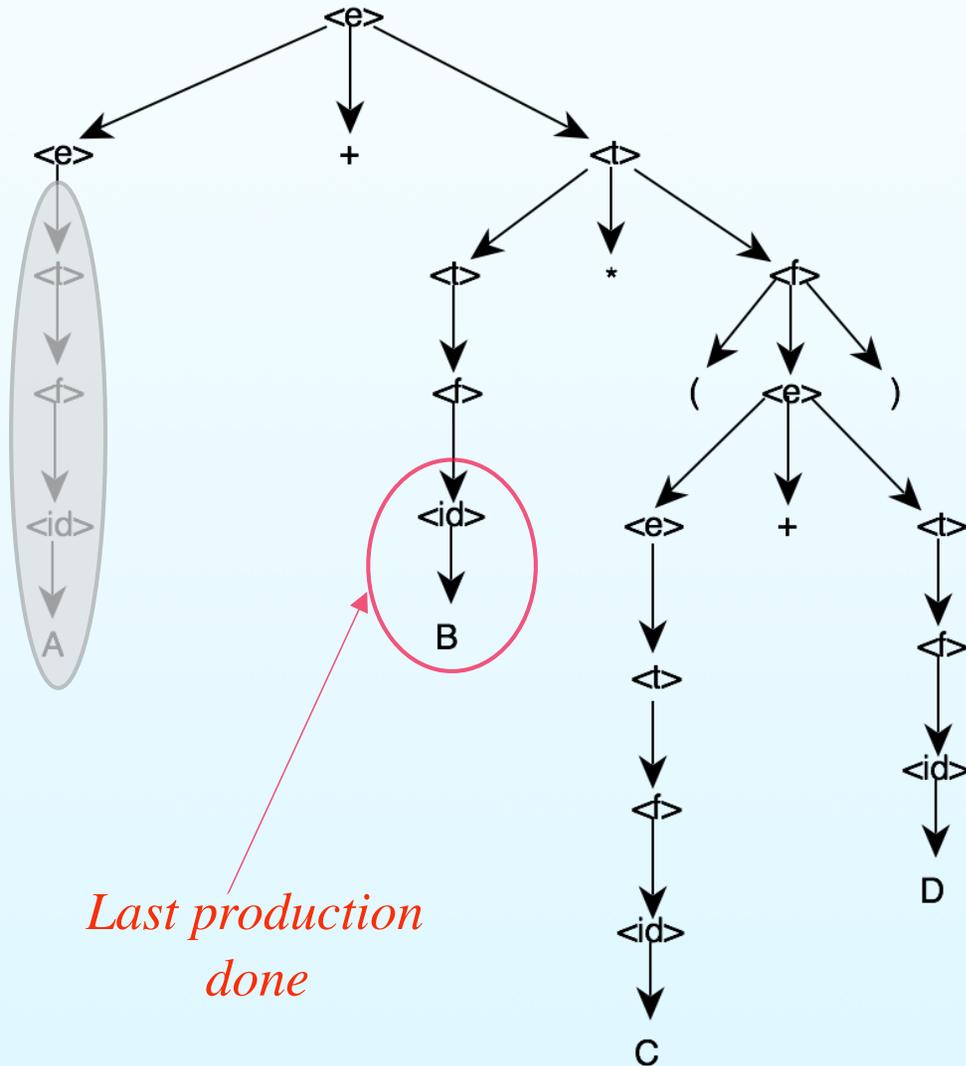
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Last production done

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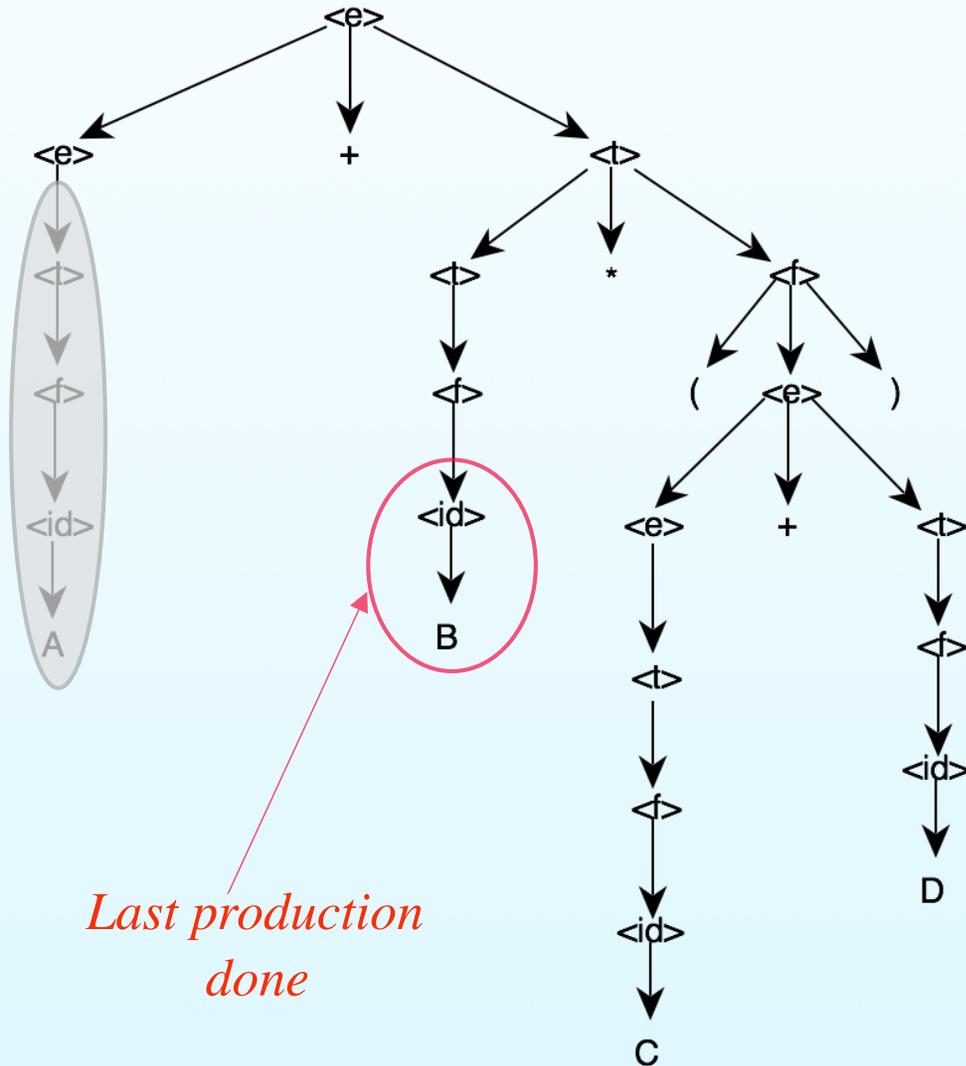
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Hypothesis: Handle is leftmost RHS?
What's different?



Last production done

Finding the handle

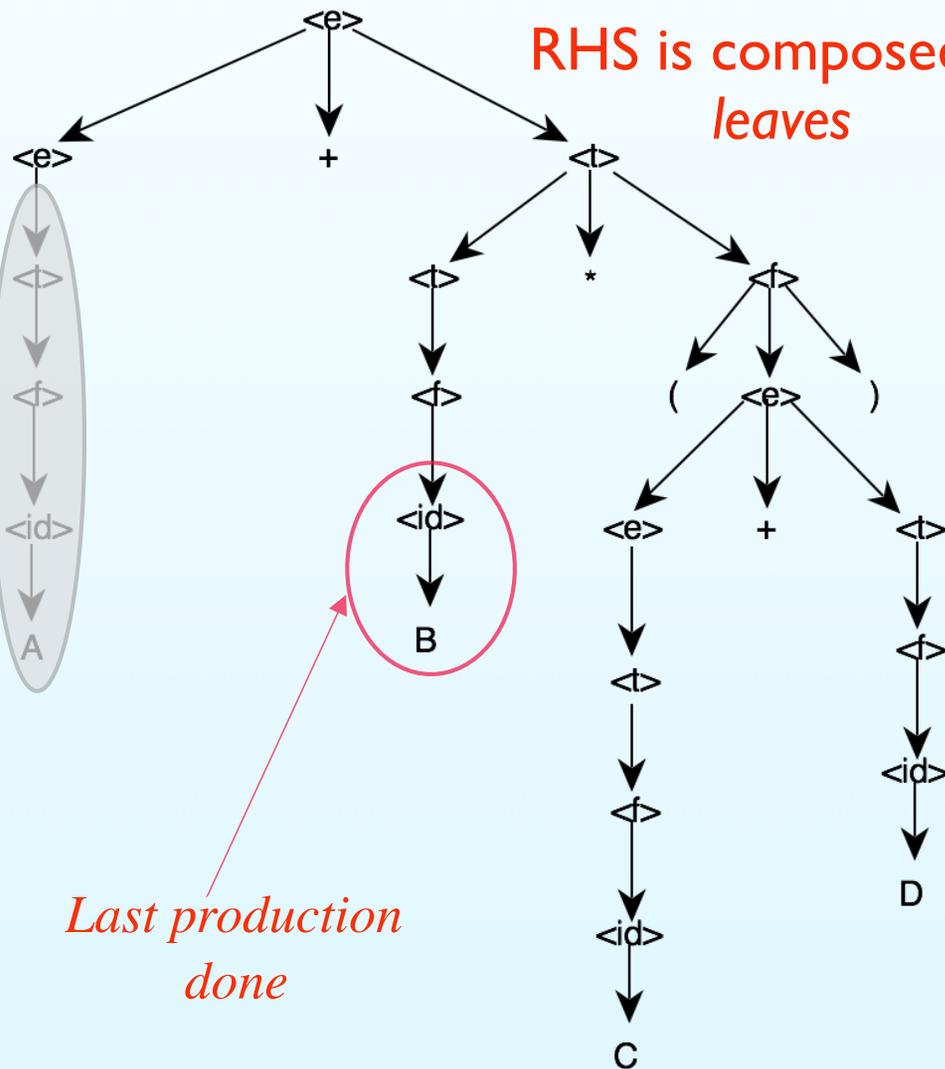
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Hypothesis: Handle is leftmost RHS?
What's different?
RHS is composed of leaves



Last production done

Definitions

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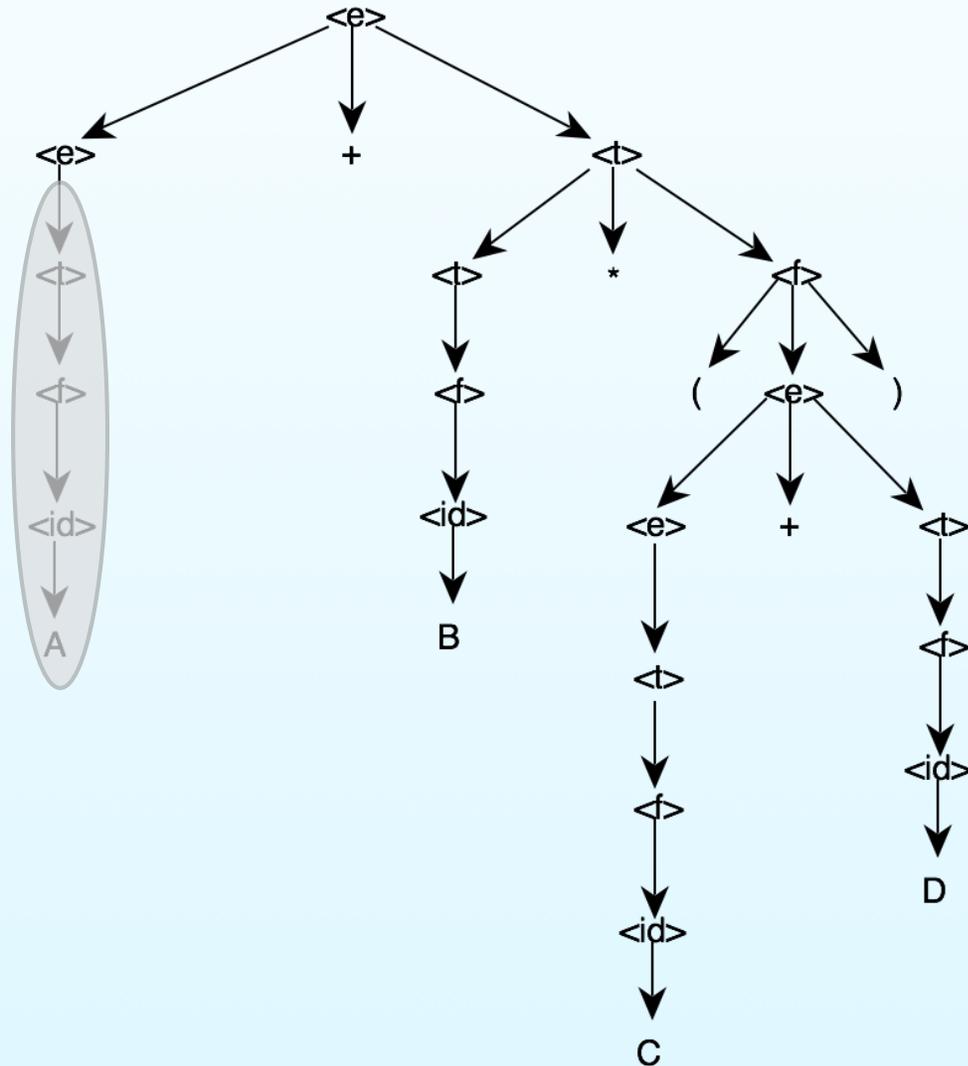
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Phrases:

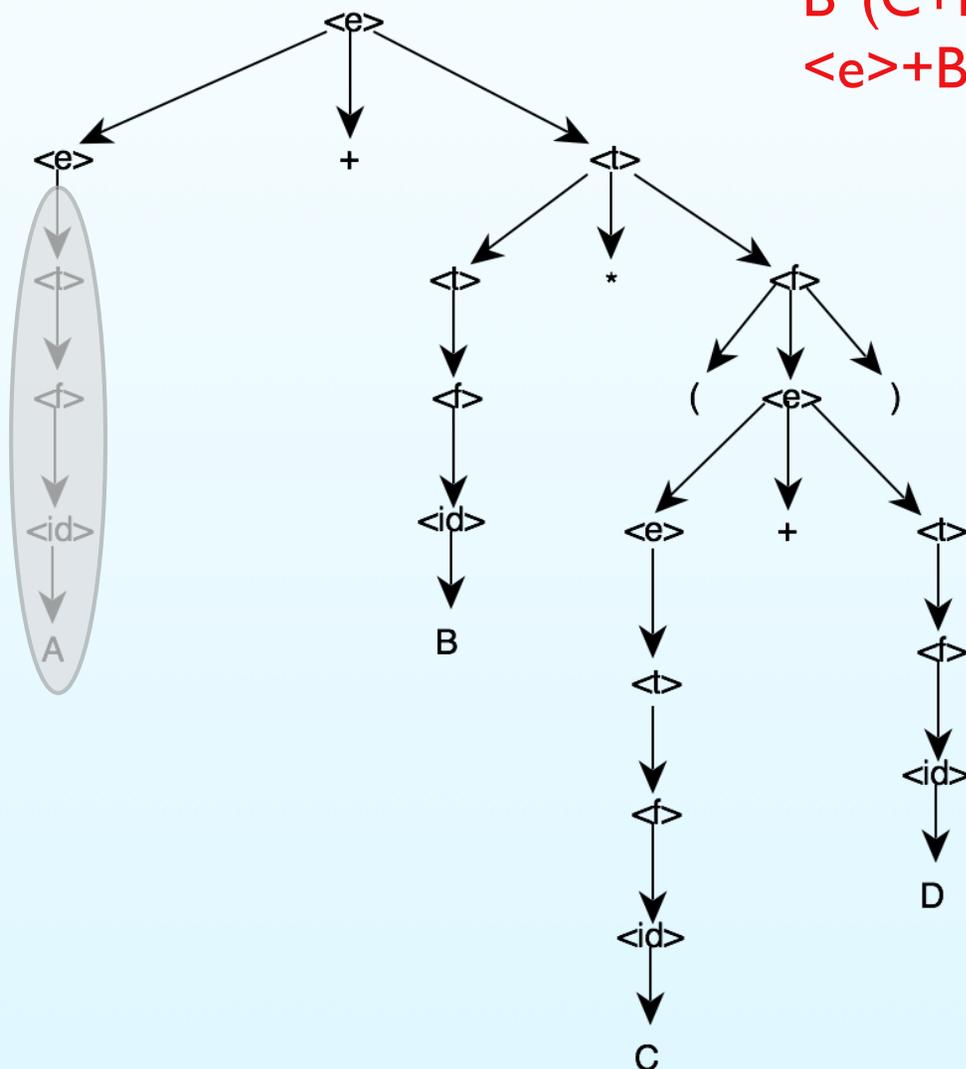
B, C, D

C+D

(C+D)

B*(C+D)

$\langle e \rangle + B * (C + D)$



Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Phrases:

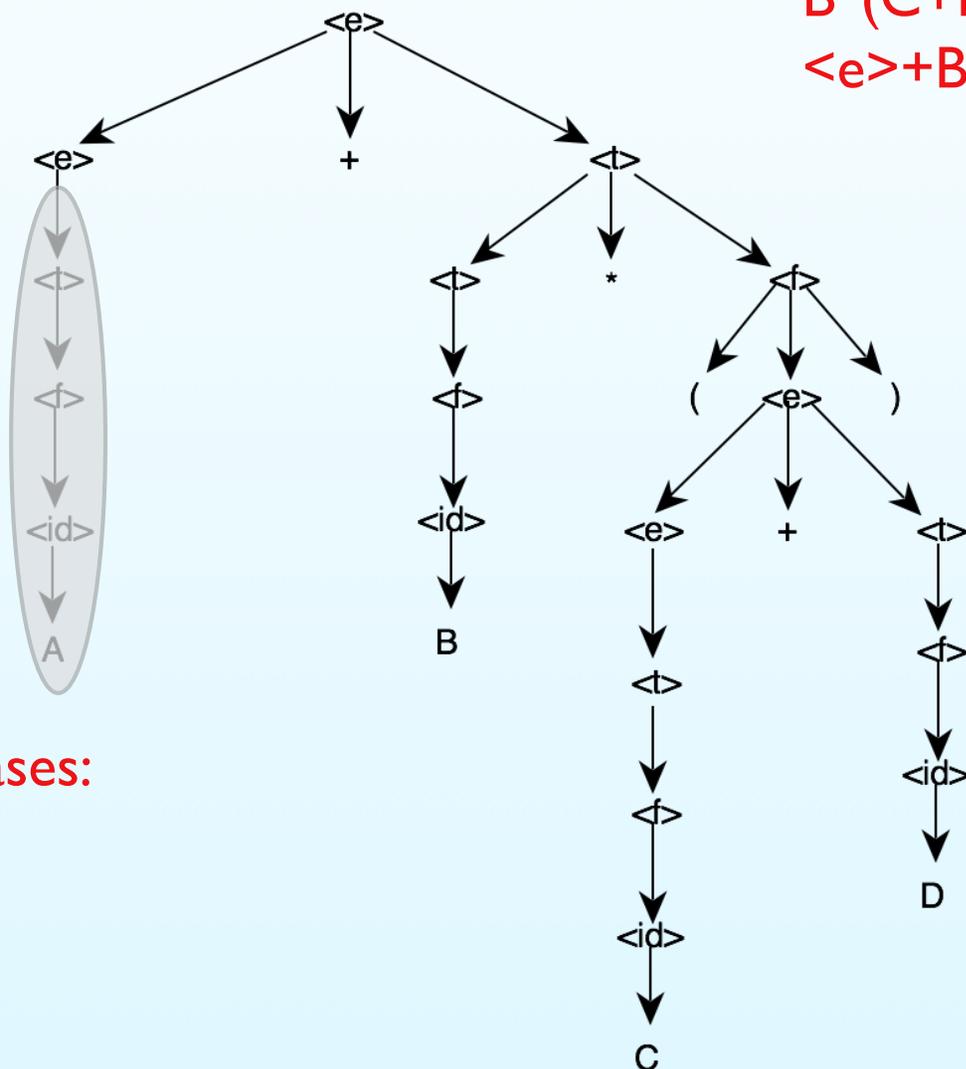
B, C, D

C+D

(C+D)

$B*(C+D)$

$\langle e \rangle + B*(C+D)$



Simple phrases:

B, C, D

Finding the handle

Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$
- $\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$
- $\langle e \rangle + \langle t \rangle * (C + D)$
- $\langle e \rangle + \langle f \rangle * (C + D)$
- $\langle e \rangle + \langle id \rangle * (C + D)$
- $\langle e \rangle + B * (C + D)$
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- ~~$A + B * (C + D)$~~

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Phrases:

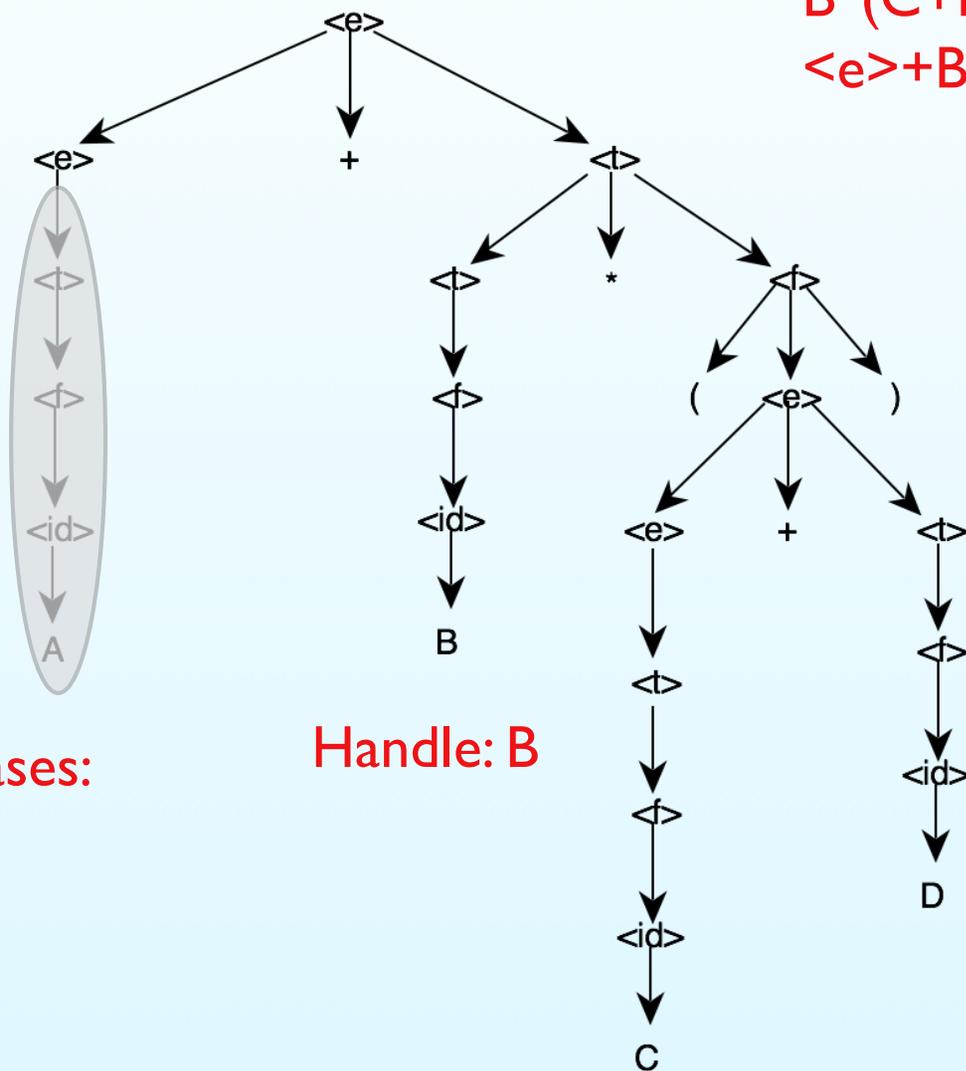
B, C, D

C+D

(C+D)

$B*(C+D)$

$\langle e \rangle + B*(C+D)$



Simple phrases:
B, C, D

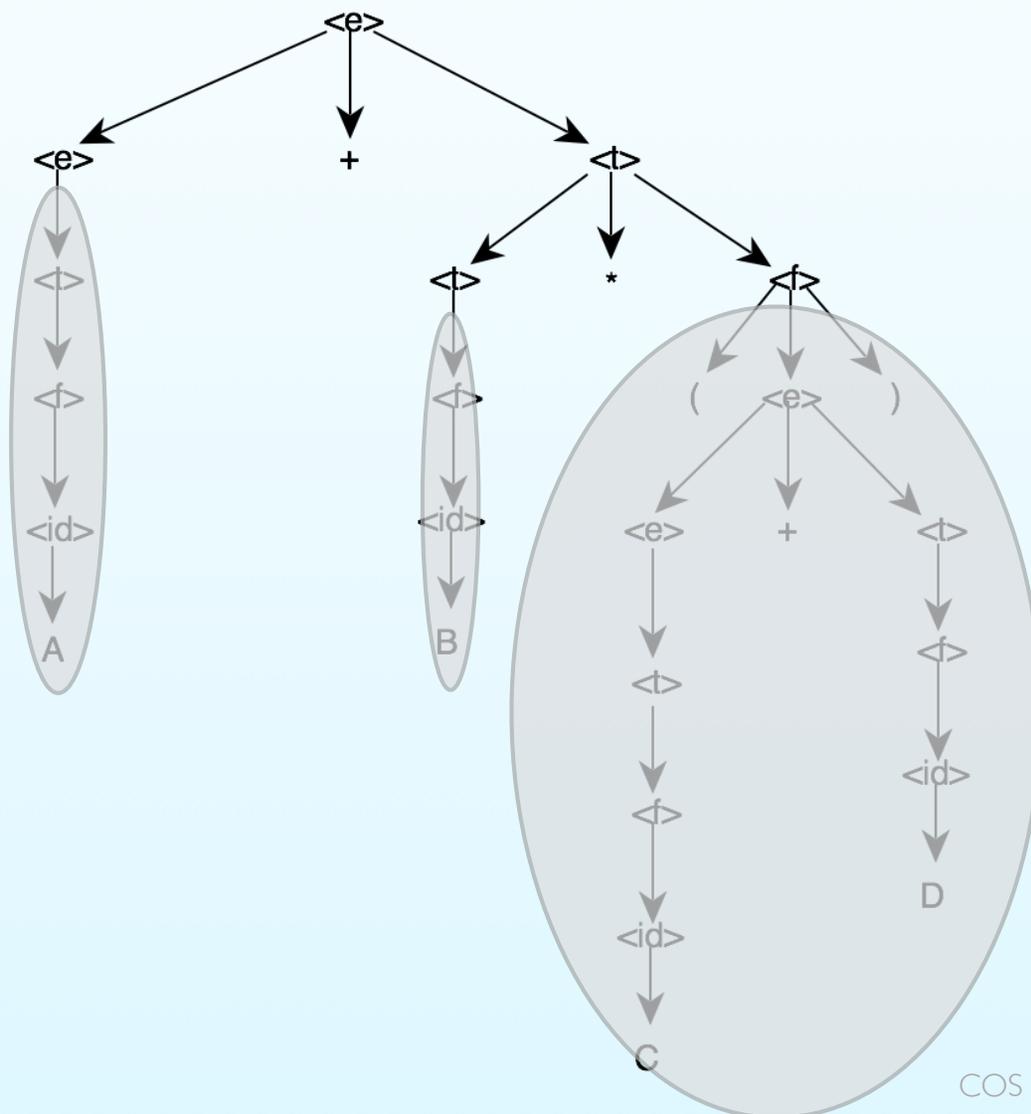
Handle: B

Finding the handle

Parse: A + B * (C + D)

- <e>
- <e> + <t>
- <e> + <t> * <f>
- ~~<e> + <t> * (<e>)~~
- ~~<e> + <t> * (<e> + <t>)~~
- ~~<e> + <t> * (<e> + <f>)~~
- ~~<e> + <t> * (<e> + <id>)~~
- ~~<e> + <t> * (<e> + D)~~
- ~~<e> + <t> * (<t> + D)~~
- ~~<e> + <t> * (<f> + D)~~
- ~~<e> + <t> * (<id> + D)~~
- ~~<e> + <t> * (C + D)~~
- ~~<e> + <f> * (C + D)~~
- ~~<e> + <id> * (C + D)~~
- ~~<e> + B * (C + D)~~
- ~~<t> + B * (C + D)~~
- ~~<f> + B * (C + D)~~
- ~~<id> + B * (C + D)~~
- A + B * (C + D)

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$



Finding the handle

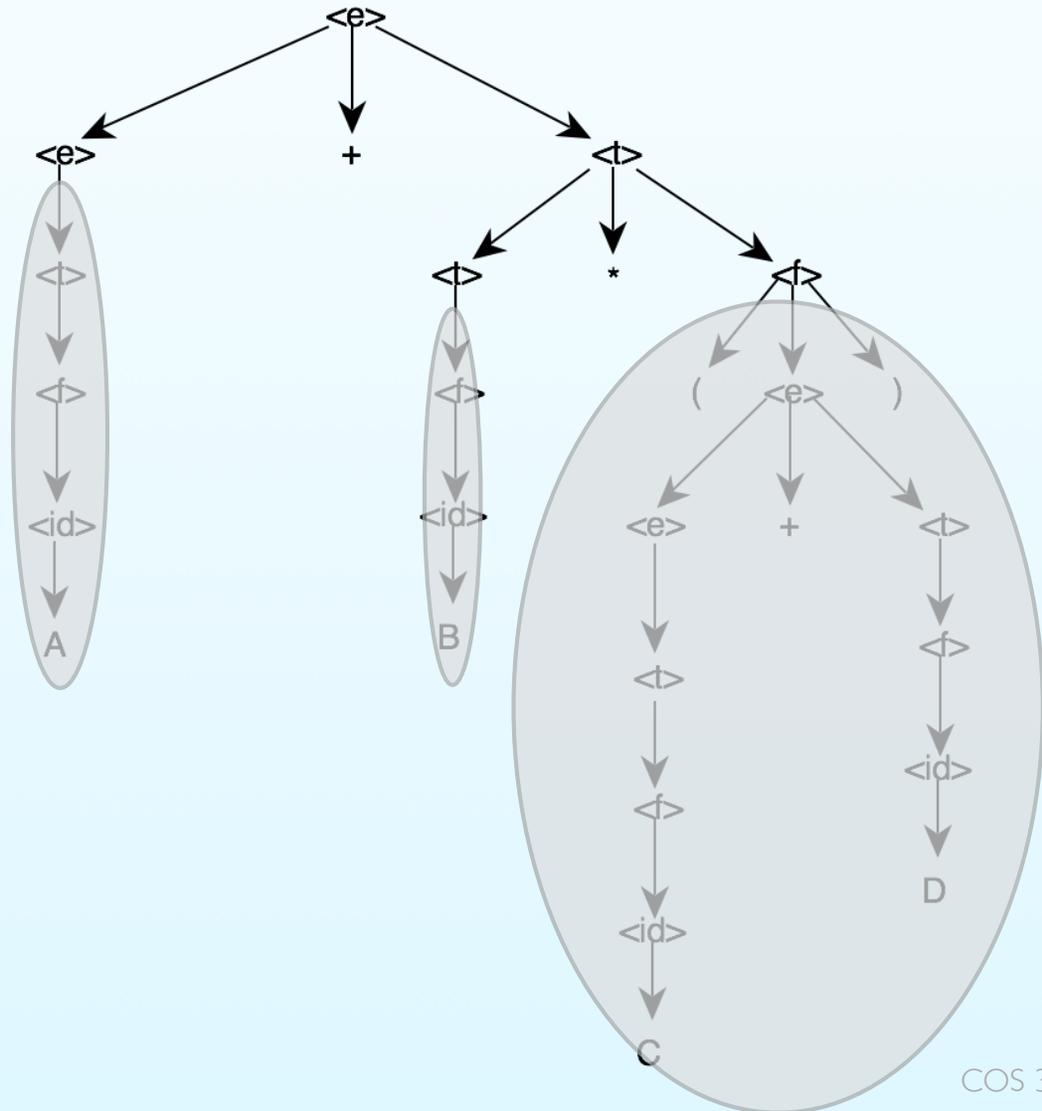
Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Phrases:

$\langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$



Finding the handle

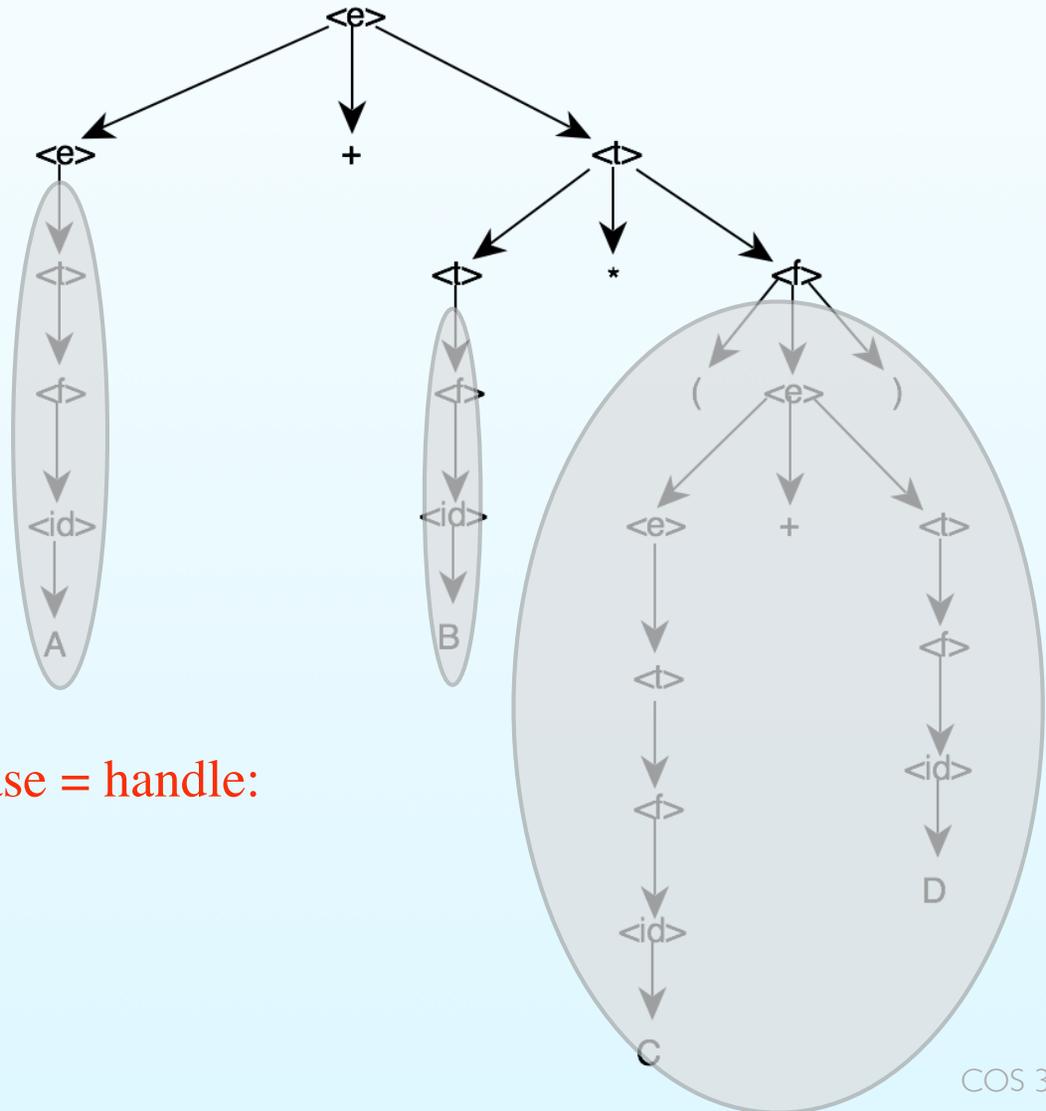
Parse: $A + B * (C + D)$

- $\langle e \rangle$
- $\langle e \rangle + \langle t \rangle$
- $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle e \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle t \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle f \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (\langle id \rangle + D)$~~
- ~~$\langle e \rangle + \langle t \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle f \rangle * (C + D)$~~
- ~~$\langle e \rangle + \langle id \rangle * (C + D)$~~
- ~~$\langle e \rangle + B * (C + D)$~~
- ~~$\langle t \rangle + B * (C + D)$~~
- ~~$\langle f \rangle + B * (C + D)$~~
- ~~$\langle id \rangle + B * (C + D)$~~
- $A + B * (C + D)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Phrases:

$\langle t \rangle * \langle f \rangle$
 $\langle e \rangle + \langle t \rangle * \langle f \rangle$



Simple phrase = handle:

$\langle t \rangle * \langle f \rangle$

Finding handles

- Consider: $\langle e \rangle + \langle t \rangle * \langle f \rangle$
- If have parse tree, know $\langle t \rangle * \langle f \rangle$ is simple phrase
- But if had parse tree...wouldn't have to parse!
- How to know $\langle e \rangle + \langle t \rangle$ isn't handle?
- Suppose we chose it as handle
 - $\Rightarrow \langle e \rangle * \langle f \rangle$
 - No rule RHS begins $\langle e \rangle *$
 - So $\langle e \rangle + \langle t \rangle$ isn't the handle
- *Lookahead*

LR(1)

- How far to look ahead?
- Want minimal amount: ideally 1 token
- Grammars that support this: *LR(1) grammars*
 - L: left-to-right (& bottom-up) parsing
 - R: rightmost derivation
 - (1): single token lookahead

Bottom-up LR(1) parsing

- Start with leftmost token
- If most recently-read leftmost tokens are handle
 - Replace with LHS
- If not, then read next token
- Handle:
 - if most recently-read
 - if lookahead token doesn't conflict with potential LHS

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle \text{id} \rangle + \langle \text{id} \rangle * (\langle \text{id} \rangle + \langle \text{id} \rangle)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle \text{id} \rangle$

Parsing example

Sentential form:

Next Rest

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle \text{id} \rangle + \langle \text{id} \rangle * (\langle \text{id} \rangle + \langle \text{id} \rangle)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle \text{id} \rangle$

Parsing example

Sentential form:

<u>Next</u>	<u>Rest</u>
<id>	+ <id> * (<id> + <id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
<id>	<id>	+ <id> * (<id> + <id>)
<id> +	+	<id> * (<id> + <id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)
	don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *	
<e> + <t> *	(<id> + <id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> | <t>
 <t> \rightarrow <t> * <f> | <f>
 <f> \rightarrow (<e>) | <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)
	don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *	
<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> | <t>
 <t> \rightarrow <t> * <f> | <f>
 <f> \rightarrow (<e>) | <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)
	don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *	
<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> | <t>
 <t> \rightarrow <t> * <f> | <f>
 <f> \rightarrow (<e>) | <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)
	don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *	
<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> | <t>
 <t> \rightarrow <t> * <f> | <f>
 <f> \rightarrow (<e>) | <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)
don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *		
<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> | <t>
 <t> \rightarrow <t> * <f> | <f>
 <f> \rightarrow (<e>) | <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)
don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *		
<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)
<e> + <t> * (<e>	+	<id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> | <t>
 <t> \rightarrow <t> * <f> | <f>
 <f> \rightarrow (<e>) | <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *

<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)
<e> + <t> * (<e>	+	<id>)
<e> + <t> * (<e> +	<id>)

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *

<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)
<e> + <t> * (<e>	+	<id>)
<e> + <t> * (<e> +	<id>)
<e> + <t> * (<e> + <id>)	

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> | <t>
 <t> \rightarrow <t> * <f> | <f>
 <f> \rightarrow (<e>) | <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *

<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)
<e> + <t> * (<e>	+	<id>)
<e> + <t> * (<e> +	<id>)
<e> + <t> * (<e> + <id>)	
<e> + <t> * (<e> + <f>)	

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *

<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)
<e> + <t> * (<e>	+	<id>)
<e> + <t> * (<e> +	<id>)
<e> + <t> * (<e> + <id>)	
<e> + <t> * (<e> + <f>)	
<e> + <t> * (<e> + <t>)	

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> | <t>
 <t> \rightarrow <t> * <f> | <f>
 <f> \rightarrow (<e>) | <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *

<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)
<e> + <t> * (<e>	+	<id>)
<e> + <t> * (<e> +	<id>)
<e> + <t> * (<e> + <id>)	
<e> + <t> * (<e> + <f>)	
<e> + <t> * (<e> + <t>)	
<e> + <t> * (<e>)	

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> | <t>
 <t> \rightarrow <t> * <f> | <f>
 <f> \rightarrow (<e>) | <id>

Parsing example

Sentential form:

Next

Rest

	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *

<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)
<e> + <t> * (<e>	+	<id>)
<e> + <t> * (<e> +	<id>)
<e> + <t> * (<e> + <id>)	
<e> + <t> * (<e> + <f>)	
<e> + <t> * (<e> + <t>)	
<e> + <t> * (<e>)	
<e> + <t> * (<e>)	

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

Next

Rest

	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *

<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)
<e> + <t> * (<e>	+	<id>)
<e> + <t> * (<e> +	<id>)
<e> + <t> * (<e> + <id>)	
<e> + <t> * (<e> + <f>)	
<e> + <t> * (<e> + <t>)	
<e> + <t> * (<e>)	
<e> + <t> * (<e>)	
<e> + <t> * <f>		

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *

<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)
<e> + <t> * (<e>	+	<id>)
<e> + <t> * (<e> +	<id>)
<e> + <t> * (<e> + <id>)	
<e> + <t> * (<e> + <f>)	
<e> + <t> * (<e> + <t>)	
<e> + <t> * (<e>)	
<e> + <t> * (<e>)		
<e> + <t> * <f>		
<e> + <t>		

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing example

Sentential form:

	<u>Next</u>	<u>Rest</u>
	<id>	+ <id> * (<id> + <id>)
<id>	+	<id> * (<id> + <id>)
<f>	+	<id> * (<id> + <id>)
<t>	+	<id> * (<id> + <id>)
<e>	+	<id> * (<id> + <id>)
<e> +	<id>	* (<id> + <id>)
<e> + <id>	*	(<id> + <id>)
<e> + <f>	*	(<id> + <id>)
<e> + <t>	*	(<id> + <id>)

don't replace <t> or <e> + <t> w/ <e>: no rule RHS starting w/ <e> *

<e> + <t> *	(<id> + <id>)
<e> + <t> * (<id>	+ <id>)
<e> + <t> * (<id>	+	<id>)
<e> + <t> * (<f>	+	<id>)
<e> + <t> * (<t>	+	<id>)
<e> + <t> * (<e>	+	<id>)
<e> + <t> * (<e> +	<id>)
<e> + <t> * (<e> + <id>)	
<e> + <t> * (<e> + <f>)	
<e> + <t> * (<e> + <t>)	
<e> + <t> * (<e>)	
<e> + <t> * (<e>)		
<e> + <t> * <f>		
<e> + <t>		
<e>		

Parse: A + B * (C + D)

Assume lex. analy. \Rightarrow <id> + <id> * (<id> + <id>)

<e> \rightarrow <e> + <t> <t>
<t> \rightarrow <t> * <f> <f>
<f> \rightarrow (<e>) <id>

Parsing

- Intuitively, this is simple...
- ...but how to implement?

Shift-reduce parsing

- *Shift-reduce parsing:*
 - Bottom-up parsing of LR(1) grammars
 - Automates the process of finding handles
- CFG needs *pushdown automaton* \Rightarrow *stack*
 - Top-down, recursive-descent parser: runtime stack
 - Shift-reduce parser \Rightarrow explicit stack
- Stack:
 - Contains parsing *state*
 - Current right sentential form
- Only 2 actions:
 - **shift** a token from input \rightarrow stack
 - **reduce** top-of-stack (n frames) \rightarrow LHS
- Reduce only done when top n frames = handle

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle \text{id} \rangle + \langle \text{id} \rangle * (\langle \text{id} \rangle + \langle \text{id} \rangle)$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle \text{id} \rangle$

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Input

Next Action

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Input

$\mid \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Next Action

shift

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle id \rangle$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Input

$\mid \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Next Action

shift

reduce

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle id \rangle$

$\langle f \rangle$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Input

$\mid \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Next Action

shift

reduce

reduce

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle id \rangle$

$\langle f \rangle$

$\langle t \rangle$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Input

$\mid \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Next Action

shift

reduce

reduce

reduce

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle id \rangle$
 $\langle f \rangle$
 $\langle t \rangle$
 $\langle e \rangle$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Input

$ \langle id \rangle$	$+$	$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$
$\langle id \rangle $	$+$	$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$
$\langle id \rangle $	$+$	$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$
$\langle id \rangle $	$+$	$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$
$\langle id \rangle $	$+$	$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$

Next Action

shift
reduce
reduce
reduce
shift

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle id \rangle$
 $\langle f \rangle$
 $\langle t \rangle$
 $\langle e \rangle$
 $\langle e \rangle +$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Input

$ \langle id \rangle$	$+$	$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$
$\langle id \rangle $	$+$	$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$
$\langle id \rangle $	$+$	$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$
$\langle id \rangle $	$+$	$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$
$\langle id \rangle $	$+$	$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$
$\langle id \rangle + $		$\langle id \rangle$	$*$	$($	$\langle id \rangle$	$+$	$\langle id \rangle$	$)$

Next Action

shift
reduce
reduce
reduce
shift
shift

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle id \rangle$
 $\langle f \rangle$
 $\langle t \rangle$
 $\langle e \rangle$
 $\langle e \rangle +$
 $\langle e \rangle + \langle id \rangle$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Input

$| \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + | \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$

Next Action

shift
reduce
reduce
reduce
shift
shift
reduce

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle id \rangle$

$\langle f \rangle$

$\langle t \rangle$

$\langle e \rangle$

$\langle e \rangle +$

$\langle e \rangle + \langle id \rangle$

$\langle e \rangle + \langle f \rangle$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Input

$\mid \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle + \mid \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle + \langle id \rangle \mid * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle + \langle id \rangle \mid * (\langle id \rangle + \langle id \rangle)$

Next Action

shift

reduce

reduce

reduce

shift

shift

reduce

reduce

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle id \rangle$

$\langle f \rangle$

$\langle t \rangle$

$\langle e \rangle$

$\langle e \rangle +$

$\langle e \rangle + \langle id \rangle$

$\langle e \rangle + \langle f \rangle$

$\langle e \rangle + \langle t \rangle$

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$

$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$

$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Input

$\mid \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle \mid + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle + \mid \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle + \langle id \rangle \mid * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle + \langle id \rangle \mid * (\langle id \rangle + \langle id \rangle)$

$\langle id \rangle + \langle id \rangle \mid * (\langle id \rangle + \langle id \rangle)$

Next Action

shift

reduce

reduce

reduce

shift

shift

reduce

reduce

shift

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

$\langle id \rangle$
 $\langle f \rangle$
 $\langle t \rangle$
 $\langle e \rangle$
 $\langle e \rangle +$
 $\langle e \rangle + \langle id \rangle$
 $\langle e \rangle + \langle f \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle *$

Input

$| \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + | \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle * | (\langle id \rangle + \langle id \rangle)$

Next Action

shift
 reduce
 reduce
 reduce
 shift
 shift
 reduce
 reduce
 shift
 shift

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

$\langle id \rangle$
 $\langle f \rangle$
 $\langle t \rangle$
 $\langle e \rangle$
 $\langle e \rangle +$
 $\langle e \rangle + \langle id \rangle$
 $\langle e \rangle + \langle f \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle *$
 $\langle e \rangle + \langle t \rangle * ($

Input

$| \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + | \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle * | (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle * (| \langle id \rangle + \langle id \rangle)$

Next Action

shift
reduce
reduce
reduce
shift
shift
reduce
reduce
shift
shift
shift
shift

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

$\langle id \rangle$
 $\langle f \rangle$
 $\langle t \rangle$
 $\langle e \rangle$
 $\langle e \rangle +$
 $\langle e \rangle + \langle id \rangle$
 $\langle e \rangle + \langle f \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle *$
 $\langle e \rangle + \langle t \rangle * ($
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle$

Input

$| \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + | \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle * | (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle * (| \langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle * (\langle id \rangle | + \langle id \rangle)$

Next Action

shift
reduce
reduce
reduce
shift
shift
reduce
reduce
shift
shift
shift
shift
reduce

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
 $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
 $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

$\langle id \rangle$
 $\langle f \rangle$
 $\langle t \rangle$
 $\langle e \rangle$
 $\langle e \rangle +$
 $\langle e \rangle + \langle id \rangle$
 $\langle e \rangle + \langle f \rangle$
 $\langle e \rangle + \langle t \rangle$
 $\langle e \rangle + \langle t \rangle *$
 $\langle e \rangle + \langle t \rangle * ($
 $\langle e \rangle + \langle t \rangle * (\langle id \rangle$
 $\langle e \rangle + \langle t \rangle * (\langle f \rangle$

Input

$| \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle | + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + | \langle id \rangle * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle | * (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle * | (\langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle * (| \langle id \rangle + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle * (\langle id \rangle | + \langle id \rangle)$
 $\langle id \rangle + \langle id \rangle * (\langle id \rangle | + \langle id \rangle)$

Next Action

shift
reduce
reduce
reduce
shift
shift
shift
shift
shift
reduce
reduce

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

<u>Stack</u>	<u>Input</u>	<u>Next Action</u>
	$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack	Input	Next Action
	$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack	Input	Next Action
	$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack	Input	Next Action
	$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack	Input	Next Action
	$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

<u>Stack</u>	<u>Input</u>	<u>Next Action</u>
	$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

<u>Stack</u>	<u>Input</u>	<u>Next Action</u>
	$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack	Input	Next Action
	$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle) $	reduce

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack	<div style="border: 1px solid black; padding: 5px; display: inline-block;"> $\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$ $\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$ $\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$ </div>	Input	Next Action
		$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle +$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle) $	reduce
$\langle e \rangle + \langle t \rangle * \langle f \rangle$		$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle) $	reduce

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack	Input	Next Action
	$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle) $	reduce
$\langle e \rangle + \langle t \rangle * \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle) $	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle) $	reduce

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Parsing example

Parse: $A + B * (C + D)$

Assume lex. analy. \Rightarrow $\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$

Stack	Input	Next Action
	$ \langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle *$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * ($	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle +$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle id \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	reduce
$\langle e \rangle + \langle t \rangle * (\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle)$	shift
$\langle e \rangle + \langle t \rangle * (\langle e \rangle)$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle) $	reduce
$\langle e \rangle + \langle t \rangle * \langle f \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle) $	reduce
$\langle e \rangle + \langle t \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle) $	reduce
$\langle e \rangle$	$\langle id \rangle + \langle id \rangle * (\langle id \rangle + \langle id \rangle) $	DONE

$\langle e \rangle \rightarrow \langle e \rangle + \langle t \rangle \mid \langle t \rangle$
$\langle t \rangle \rightarrow \langle t \rangle * \langle f \rangle \mid \langle f \rangle$
$\langle f \rangle \rightarrow (\langle e \rangle) \mid \langle id \rangle$

Shift-reduce parser

- Consider all of stack + 1 character (for LR(1)) to determine what to do
- Problem:
 - read input \Rightarrow stack grows \Rightarrow slower
- Solution:
 - At any time, parser is in particular *state*
 - Can think of it as a state machine
 - Table-based implementation
 - Index: state + lookahead
 - Specifies what to do, next state

LR parsers

- *LR parser*: table-based shift-reduce parser
 - Small program
 - Most knowledge of parsing encoded in *parsing table*
- Parsing table:
 - built from grammar
 - producing table can be difficult
 - usually by automated tool (e.g., yacc)
- First LR (“canonical LR”) ← Don Knuth (1965)
 - simpler LR algorithms used to simplify table-building (e.g., SLR, LALR)...
 - ...but work on smaller classes of grammars

Advantages of LR parsers

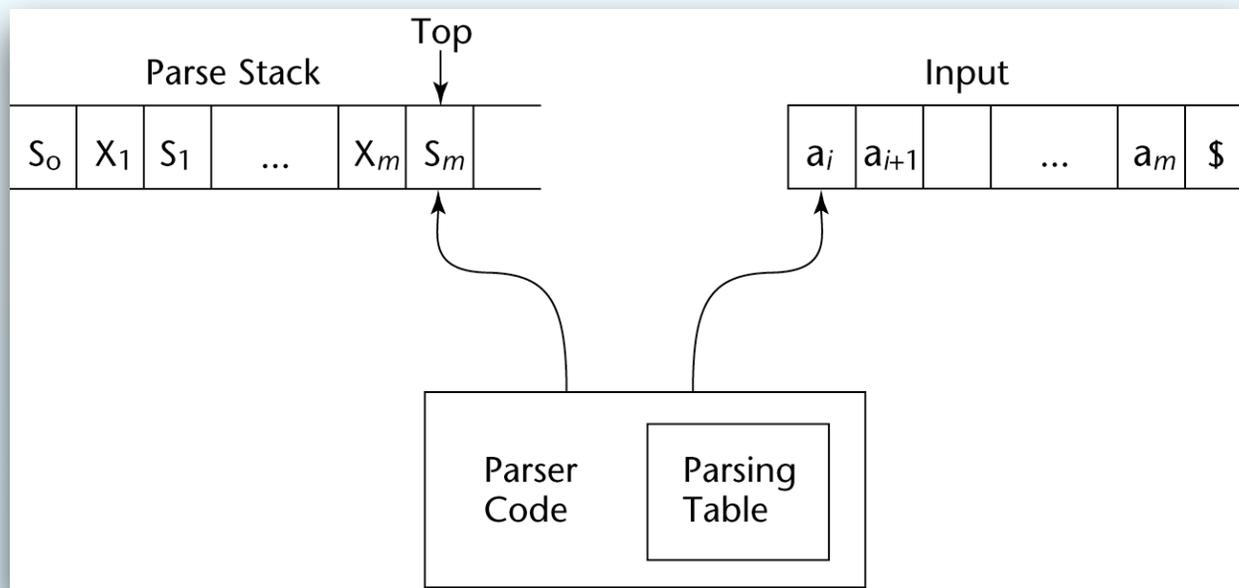
- Can be used for all programming languages
- Detect syntax errors as soon as possible (reading left-to-right)
- LR grammar class \supset LL grammar class (e.g., LR includes many left recursive grammars)

LR parsing

- Knuth's insight: entire history of parse is available (in stack) for parser to use to make decision
- There is a finite, small set of different parse situations that have occurred
- So the parse stack can contain the parser state

LR parsers

- Uses a stack + input string
- At any given time, parser is considered to be in a particular, named *state*
- Stack alternates token/state



LR parsers

- Parsing table – really two tables:
 - **Action** table:
 - state + next token \Rightarrow next action
 - actions: shift & go to state; reduce by a rule; accept; error
 - **Goto** table: state + token on TOS \Rightarrow next state
- Built by (e.g.) yacc or bison
- On error, call error-handling routine

LR parsing: example

1. $E \rightarrow E + T$

2. $E \rightarrow T$

3. $T \rightarrow T * F$

4. $T \rightarrow F$

5. $F \rightarrow (E)$

6. $F \rightarrow id$

	Action						Goto		
State	id	+	*	()	\$	E	T	F
0	S5		S4				1	2	3
1		S6				accept			
2		R2	S7		R2	R2			
3		R4	R4		R4	R4			
4	S5			S4			8	2	3
5		R6	R6		R6	R6			
6	S5			S4				9	3
7	S5			S4					10
8		S6			S11				
9		R1	S7		R1	R1			
10		R3	R3		R3	R3			
11		R5	R5		R5	R5			

LR parsing: example

Stack	Input	Action									
0	id+id*id\$	Shift 5									
0id5			Action							Goto	
0F3		State	id	+	*	()	\$	E	T	
0T2		0	S5		S4				1	2	
0EI		1		S6				accept			
0EI+6		2		R2	S7		R2	R2			
0EI+6id5		3		R4	R4		R4	R4			
0EI+6F3		4	S5			S4			8	2	
0EI+6T9		5		R6	R6		R6	R6			
0EI+6T9*7		6	S5			S4				9	
0EI+6T9*7id5		7	S5			S4					
0EI+6T9*7F10		8		S6			S11				
0EI+6T9		9		R1	S7		R1	R1			
0EI		10		R3	R3		R3	R3			
		11		R5	R5		R5	R5			

LR parsing: example

Which grammar rule to use

Stack		Input						Action			
0		id+id*id\$						Shift 5			
0id5		+id*id\$						Reduce 6, goto 3 (0/F)			
0F3		Action						Goto			
0T2		State	id	+	*	()	\$	E	T	F
0EI		0	S5		S4				1	2	3
0EI+6		1		S6				accept			
0EI+6id5		2		R2	S7		R2	R2			
0EI+6F3		3		R4	R4		R4	R4			
0EI+6T9		4	S5			S4			8	2	3
0EI+6T9*7		5		R6	R6		R6	R6			
0EI+6T9*7id		6	S5			S4				9	3
0EI+6T9*7F3		7	S5			S4					10
0EI+6T9		8		S6			S11				
0EI		9		R1	S7		R1	R1			
		10		R3	R3		R3	R3			
		11		R5	R5		R5	R5			

LR parsing: example

Stack	Input	Action
0	id+id*id\$	Shift 5
0id5	+id*id\$	Reduce 6, goto 3 (0/F)
0F3	+id*id\$	R4, goto 2 (0/T)
0T2	+id*id\$	R2, goto 1 (0/E)

State	Action							Goto		
	id	+	*	()	\$	E	T	F	
0	S5		S4				1	2	3	
1		S6				accept				
2		R2	S7		R2	R2				
3		R4	R4		R4	R4				
4	S5			S4			8	2	3	
5		R6	R6		R6	R6				
6	S5			S4				9	3	
7	S5			S4					10	
8		S6			S11					
9		R1	S7		R1	R1				
10		R3	R3		R3	R3				
11										

LR parsing: example

Stack	Input	Action
0	id+id*id\$	Shift 5
0id5	+id*id\$	Reduce 6, goto 3 (0/F)
0F3	+id*id\$	R4, goto 2 (0/T)
0T2	+id*id\$	R2, goto 1 (0/E)
0E1	+id*id\$	Shift 6
0E1+6	id*id\$	Shift 5
0E1+6id5	*id\$	R6, goto 3 (6/F)

Stack	State	Action						Goto		
		id	+	*	()	\$	E	T	F
0E1+6F3	0	S5		S4				1	2	3
0E1+6T9	1		S6			accept				
0E1+6T9*7	2		R2	S7		R2	R2			
0E1+6T9*7id	3		R4	R4		R4	R4			
0E1+6T9*7FI	4	S5			S4			8	2	3
0E1+6T9	5		R6	R6		R6	R6			
0E1	6	S5			S4				9	3
	7	S5			S4					10
	8		S6			S11				

State	Action						Goto		
	id	+	*	()	\$	E	T	F
0	S5		S4				1	2	3
1		S6				accept			
2		R2	S7		R2	R2			
3		R4	R4		R4	R4			
4	S5			S4			8	2	3
5		R6	R6		R6	R6			
6	S5			S4				9	3
7	S5			S4					10
8		S6			S11				
9		R1	S7		R1	R1			
10		R3	R3		R3	R3			
11		R5	R5		R5	R5			

Stack		
0		
0id5		
0F3		
0T2		
0E1		
0E1+6		
0E1+6id5	^id\$	R6, goto 3 (6/F)
0E1+6F3	*id\$	R4, goto 9 (6/T)
0E1+6T9	*id\$	Shift 7
0E1+6T9*7	id\$	Shift 5
0E1+6T9*7id5	\$	R6, goto 10 (7/F)
0E1+6T9*7F10	\$	R3, goto 9 (6/T)
0E1+6T9	\$	R1, goto 1 (0/E)
0E1	R	Accept

LR parsing: For you

1. $E \rightarrow E + T$

2. $E \rightarrow T$

3. $T \rightarrow T * F$

4. $T \rightarrow F$

5. $F \rightarrow (E)$

6. $F \rightarrow id$

Use table to parse:
 $id * (id + id)$

	Action						Goto		
State	id	+	*	()	\$	E	T	F
0	S5		S4				1	2	3
1		S6				accept			
2		R2	S7		R2	R2			
3		R4	R4		R4	R4			
4	S5			S4			8	2	3
5		R6	R6		R6	R6			
6	S5			S4				9	3
7	S5			S4					10
8		S6			S11				
9		R1	S7		R1	R1			
10		R3	R3		R3	R3			
11		R5	R5		R5	R5			