COS 140: Foundations of Computer Science

Transport-Layer Protocols*

Fall 2018

*This lecture draws heavily from Kurose & Ross (2008): *Computer Networks: A Top-Down Approach*

Homework

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Reading: None
- Slides online
- Homework:
 - On Blackboard
 - Due 12/10



Protocols

Transport Layer

- Protocols
- Transport layer
- Ports & sockets
- Types

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Recall: Protocol is a description of a pattern of interaction between agents
- Most common: TCP/IP



Transport Layer

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Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP





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Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Recall: Layered *network stack*
 - Application: network applications (e.g., Web clients/servers), application protocols (e.g., HTTP)





- Protocols
- Transport layer
- Ports & sockets
- Types

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Recall: Layered *network stack*
 - Application: network applications (e.g., Web clients/servers), application protocols (e.g., HTTP)
 - Transport: Delivery of application messages between applications: error correction, reliable transport (some protocols), etc.



- Transport Layer
- Protocols
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- Types

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Reliable Transport Protocols

Example: TCP

- Recall: Layered *network stack*
 - Application: network applications (e.g., Web clients/servers), application protocols (e.g., HTTP)
 - Transport: Delivery of application messages between applications: error correction, reliable transport (some protocols), etc.
 - Network: Delivery of transport-layer messages: routing, etc.

- Transport Layer
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Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Recall: Layered *network stack*
 - Application: network applications (e.g., Web clients/servers), application protocols (e.g., HTTP)
 - Transport: Delivery of application messages between applications: error correction, reliable transport (some protocols), etc.
 - Network: Delivery of transport-layer messages: routing, etc.
 - Data link: Delivery of network-layer messages e.g., forwarding to next router/host (e.g., Ethernet, WiFi, link layers on routers)



- Transport Layer
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Reliable Transport Protocols

Example: TCP

- Recall: Layered *network stack*
 - Application: network applications (e.g., Web clients/servers), application protocols (e.g., HTTP)
 - Transport: Delivery of application messages between applications: error correction, reliable transport (some protocols), etc.
 - Network: Delivery of transport-layer messages: routing, etc.
 - Data link: Delivery of network-layer messages e.g., forwarding to next router/host (e.g., Ethernet, WiFi, link layers on routers)
 - Physical: Moving the bits (Ethernet has variants for different media)

Layered protocols

Transport Layer

- Protocols
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- Types

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- As message goes from top \longrightarrow bottom:
 - Broken into pieces
 - \circ $\,$ Each piece has its own header added at each level



Transport layer

Transport Layer

- Protocols
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- Ports & sockets
- Types

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Responsible for delivery of application-layer messages between processes
- Interacts with: application and network layers
 - Divides/encapsulates application messages as *segments*
- Possibly:
 - Error correction
 - Reliable delivery
 - Congestion control
- Connectionless (e.g., UDP) versus connection-oriented (e.g., TCP)



Ports and sockets

Transport Layer

- Protocols
- Transport layer
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Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Network stack manages set of *ports* on host:
 - Interface of host to outside world
 - Numbered (virtual) connection points
 - Some: well-known ports (e.g., email (25), Web (80))
 - Some: usable on fly by processes
- Sockets: virtual connections between application process and transport layer (and thus, a port)
- Each application process talks to (e.g.) TCP via a socket
- Can have multiple sockets attached to a port
- E.g.: Web server
- Transport layer: multiplexing/demultiplexing of sockets \leftrightarrow ports



Ports and sockets





Connection-oriented and connectionless protocols

Transport Layer

- Protocols
- Transport layer
- Ports & sockets
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Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

Other Transport Layer Topics

• Connectionless protocols:

- E.g., User Datagram Protocol (UDP)
- Send segments ("datagrams") between application layers
- No notion of a continuous connection think US Mail
- Maybe some error checking embedded
- No error correcting
- Not reliable

Connection-oriented and connectionless protocols

- Transport Layer
- Protocols
- Transport layer
- Ports & sockets
- Types

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

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- Connectionless protocols:
 - E.g., User Datagram Protocol (UDP)
 - Send segments ("datagrams") between application layers
 - No notion of a continuous connection think US Mail
 - Maybe some error checking embedded
 - No error correcting
 - Not reliable
- Connection-oriented protocols:
 - E.g., Transmission Control Protocol (TCP)
 - Connection: a virtual pipeline between applications
 - Think "phone", but no real connection
 - Reliable transport protocols

Unreliable Transport Protocols

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

Other Transport Layer Topics • Very simple

- Possibly segment application data
- Possibly add error-checking code (e.g., CRC)
- Just pass along to network layer



Reliable transport protocols

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

• Using a reliable channel

- Bit errors
- Bit error protocol
- Lossy channels
- Lossy protocol
- Pipelining
- Recovery in pipelining

Example: TCP

Other Transport Layer Topics • More complex

- Have to deal with:
 - bitwise errors
 - lost segments
 - out-of-order segments
- Can conceptualize as *state machines*
- We'll look at increasingly-complex variants



Using a reliable channel



Unreliable Transport Protocols

Reliable Transport

Protocols

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Example: TCP

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Using a channel with bit errors

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

• Using a reliable channel

• Bit errors

• Bit error protocol

• Lossy channels

Lossy protocol

• Pipelining

• Recovery in pipelining

Example: TCP

Other Transport Layer Topics • Problem: Few channels are reliable

Simple problem: errors in some bits

Detect this with: parity, cyclic redundancy checks (CRC)

The question is: what to do when detected?



Using a channel with bit errors

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

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Example: TCP

- For this, need additional messages:
 - ACK: to acknowledge correct receipt
 - NAK: negative acknowledgment
- Recipient checks the packet, sends the appropriate message in reply
- Problem: ACKs and NAKs can be garbled, too!





Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

Using a reliable

channel

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Example: TCP

Other Transport Layer Topics Recipient:



Corrupt segment

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Transport Layer

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Example: TCP

Other Transport Layer Topics Recipient:

Send segment Wait for message Good segment Extract data Deliver data Create ACK segment Send segment

Corrupt segment

Create NAK segment

• Any problems with this?



Transport Layer

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Example: TCP

Other Transport Layer Topics Recipient:

Send segment Wait for message Good segment Extract data Deliver data Create ACK segment

Corrupt segment

Create NAK segment

• Any problems with this? Yes! What about errors in ACK/NAK?

Send segment



Transport Layer

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Example: TCP

- Have the sender number its segments
- Receiver can then determine if the packet received is a retransmission
- ACK/NAK don't need to say what they're ACK'ing (or not) since no messages are lost, garbled or okay, the last ACK/NAK was for its last message
- For this simple protocol, we only need two sequence numbers, 0 and 1 (a bit) – only one packet being dealt with at a time.



Transport Layer

Unreliable Transport Protocols **Reliable Transport**

Protocols

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Example: TCP





Transport Layer

• Recipient:



Unreliable Transport Protocols

Reliable Transport

Protocols

• Using a reliable channel



• Bit error protocol

- Lossy channels
- Lossy protocol
- Pipelining

• Recovery in pipelining

Example: TCP

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Using a "lossy" channel

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

• Using a reliable channel

• Bit errors

- Bit error protocol
- Lossy channels
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- Pipelining
- Recovery in pipelining

Example: TCP

- Even harder: What if the channel can lose some messages?
 - ACKs and NAKs could also be lost or garbled.
 - Need timers, now if haven't received an ACK after some time, retransmit
- How to choose the time?

Using a "lossy" channel

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

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Example: TCP

- Even harder: What if the channel can lose some messages?
 - ACKs and NAKs could also be lost or garbled.
 - Need timers, now if haven't received an ACK after some time, retransmit
- How to choose the time? At least round trip delay + some



Protocol for lossy channels

Transport Layer

• Sender:

Unreliable Transport Protocols

Reliable Transport

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Example: TCP







Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

• Using a reliable channel

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• Lossy channels

Lossy protocol

• Pipelining

• Recovery in pipelining

Example: TCP

Other Transport Layer Topics • Problem: reliable, but inefficient

- Suppose it takes 5 ms for message to propagate from sender to receiver, or vice versa a 1 Gbps channel, and time to actually put the message on the channel is negligible
- Time to send, say, a 1 KB message:
 - Send message: 5 ms
 - Send ACK: 5 ms
 - Total 10 ms/message
 - Transfer rate = 1 KB/10 ms = only 100 KB/s!

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

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• Bit error protocol

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Example: TCP

- Better idea: don't wait for ACK before sending other messages
 - Now can have n messages "in the pipe" at once.
 - How many, potentially?



Transport Layer

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Reliable Transport Protocols

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Example: TCP

- Better idea: don't wait for ACK before sending other messages
 Now can have n messages "in the pipe" at once.
- How many, potentially?
 - If 1 Gbps channel, 1 KB messages, then a message takes

$$\frac{8Kb}{1Gb/s} = \frac{2^{13}}{2^{30}}s = 2^{-17}s \approx 7.6\mu s$$



Transport Layer

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Other Transport Layer Topics

- Better idea: don't wait for ACK before sending other messages
 Now can have n messages "in the pipe" at once.
- How many, potentially?
 - If 1 Gbps channel, 1 KB messages, then a message takes

$$\frac{8Kb}{1Gb/s} = \frac{2^{13}}{2^{30}}s = 2^{-17}s \approx 7.6\mu s$$

• So there can be $\frac{5ms}{7.6\mu s} \approx 658$ messages in the pipe at once • In practice, have (far) fewer:

- *Window*: what can be sent before an ACK received
- Receive an ACK: slide window, can transmit more



Transport Layer

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Example: TCP

- In send-and-wait protocols, pretty clear what you're ACKing
 - What about pipelined protocols?
 - Have to mark ACK with what is being acknowledged
- What to do when one is missing/corrupt?



Transport Layer

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 - Go-Back-N: When missing one (corrupt or timeout), repeat it and all others after it

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Example: TCP

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- In send-and-wait protocols, pretty clear what you're ACKing
 - What about pipelined protocols?
 - Have to mark ACK with what is being acknowledged
- What to do when one is missing/corrupt?
 - Go-Back-N: When missing one (corrupt or timeout), repeat it and all others after it
 - Selective repeat: Just repeat the one missing
 - GBN simpler, not as efficient

Example: TCP

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Segments
- Sequence numbers
- Connection initiation
- Handling problems
- Handling problems

- Connection-oriented, reliable data transport protocol
- Can handle bit errors, lossy channels
- Full-duplex



TCP segment structure

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Segments
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- Header + data
- Source port number, destination port number
- Sequence number
- ACK number
- Checksum
- Header length field, flags, options, some other stuff



Sequence numbers

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Segments
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- Each half of the conversation is considered an ordered sequence of bytes
- Sequence number of a segment is the byte number of the first byte in the segment not the segment number!
- ACK number: The next byte expected from the sender
- These are *cumulative acknowledgments*



Connection initiation: "three-way handshake"

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

• Segments

- Sequence numbers
- Connection initiation
- Handling problems
- Handling problems

- First: Client sends a special segment ("SYN segment") to request connection
 - No application data contained
 - SYN bit in header = 1
 - Random sequence number
- Second: Server sets up its side of the connection and sends message 2 ("SYNACK segment")
 - Allocates buffers, variables
 - Response segment: no application data
 - SYN = 1, ACK = client sequence number + 1, random sequence number

Connection initiation: "three-way handshake"

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Segments
- Sequence numbers
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- Handling problems

- Third: Client sets up its side, sends another message
 - Allocates client-side buffers, variables
 - Segment has SYN = 0, server's sequence number +1 as ACK
 - Can carry application data (*payload*)



Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

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Transport Layer

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Transport Layer

Unreliable Transport Protocols

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- Handling problems

- Corrupt segment received (bitwise error) don't ACK
- Receive duplicate segment -



Transport Layer

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Example: TCP

- Segments
- Sequence numbers
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- Handling problems

- Corrupt segment received (bitwise error) don't ACK
- Receive duplicate segment
 - Why? Lost ACK



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- Segments
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- Corrupt segment received (bitwise error) don't ACK
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 - Why? Lost ACK
 - Just discard data, re-ACK



Transport Layer

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Example: TCP

- Segments
- Sequence numbers
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- Handling problems
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Other Transport Layer Topics

- Corrupt segment received (bitwise error) don't ACK
- Receive duplicate segment
 - Why? Lost ACK
 - Just discard data, re-ACK
- Timeout –

Transport Layer

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- Timeout
 - Why? Lost segment



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- Timeout
 - Why? Lost segment
 - Single timer for all messages to reduce overhead



Transport Layer

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- Receive duplicate segment
 - Why? Lost ACK
 - Just discard data, re-ACK
- Timeout
 - Why? Lost segment
 - Single timer for all messages to reduce overhead
 - Retransmit segment



Transport Layer

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Transport Layer

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- Connection initiation
- Handling problems
- Handling problems

- ACK received for segment *after* one it's expecting an ACK for
 - Not really a problem
 - Cumulative acknowledgment, so previous segments received, too



Transport Layer

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Reliable Transport Protocols

Example: TCP

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- Sequence numbers
- Connection initiation
- Handling problems
- Handling problems





Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- Segments
- Sequence numbers
- Connection initiation
- Handling problems
- Handling problems

- Segment received out of order
 - One reason: a segment was lost



Transport Layer

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Example: TCP

- Segments
- Sequence numbers
- Connection initiation
- Handling problems
- Handling problems

- Segment received out of order
 - One reason: a segment was lost
 - For this, send ACK, but with next byte expected being the missing segment



Transport Layer

Unreliable Transport Protocols

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Example: TCP

• Segments

- Sequence numbers
- Connection initiation
- Handling problems
- Handling problems

- Segment received out of order
 - One reason: a segment was lost
 - For this, send ACK, but with next byte expected being the missing segment
 - Another reason: segment tied up in network



Transport Layer

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Example: TCP

• Segments

- Sequence numbers
- Connection initiation
- Handling problems
- Handling problems

- Segment received out of order
 - One reason: a segment was lost
 - For this, send ACK, but with next byte expected being the missing segment
 - Another reason: segment tied up in network
 - For this, ACK with next real byte expected, necessarily one right after this one



Other transport layer topics

Transport Layer

Unreliable Transport Protocols

Reliable Transport Protocols

Example: TCP

- How to choose timeouts
- Flow control
- Congestion control
- When to use UDP vs TCP

