

COS 140: Foundations of Computer Science

Operating Systems

Fall 2018

Introduction	2
Announcements	2
Operating systems.	3
Why have an operating system?	4
Ways of viewing an OS	5
Service provider	6
Resource manager.	7
Virtual machine	8
Why Study OS in CS?	9
Kernel	10
Processes	11
Process state	12
Process synchronization	13
Process scheduling	14
Virtual Memory	15
Overview	15
File Systems	17
Miscellaneous	18

Announcements

- Reading: None for today!
- Homework: none

Copyright © 2002–2018 UMaine Computer Science Department – 2 / 18

Operating systems

- When you deal with a computer, you're really dealing with an *operating system*.
- Examples:
 - Windows (95/98, NT, 2000, XP, 7, 8, 10)
 - Macintosh OS X
 - UNIX – including Linux
 - iOS, Android, Palm OS, WebOS

Copyright © 2002–2018 UMaine Computer Science Department – 3 / 18

Why have an operating system?

- They provide services that all programs need
- They allow multiple programs to run at once
- They are the computers' housekeepers
- They facilitate portability of programs

Copyright © 2002–2018 UMaine Computer Science Department – 4 / 18

Ways of viewing an OS

- Service provider
- Resource manager
- Virtual machine

Copyright © 2002–2018 UMaine Computer Science Department – 5 / 18

Service provider

- Most programs need (e.g.):
 - input and output
 - access to memory
 - ways to handle errors
 - timers and alarms
- Could have each program do this itself
- Operating system centralizes these (and other) services
 - Programmers don't have to program them again
 - Programs don't have to include all the code to do these things

Copyright © 2002–2018 UMaine Computer Science Department – 6 / 18

Resource manager

- What *resources* does a computer have?
 - CPU, memory
 - disk (files, etc.)
 - I/O devices
- CPU must be managed to allow *multiprogramming* (*multitasking*)
- Why must other resources be managed:
 - ...in a single-user, no multitasking environment?
 - ...with multitasking?
 - ...in a multi-user, multitasking environment?

Copyright © 2002–2018 UMaine Computer Science Department – 7 / 18

Virtual machine

- To user, “computer” \equiv “operating system”: a *virtual* machine
- OS (and utilities) \Rightarrow environment for user
- Also for programs:
 - High-level, abstract view of hardware in terms of services OS provides: an *application programmer interface* (API)
 - Increases ease of programming
 - Increases portability – e.g., POSIX

Copyright © 2002–2018 UMaine Computer Science Department – 8 / 18

Why Study OS in CS?

- Why not just a technical school topic?
- Active area of research and development
- Much underlying theory, many fundamental concepts need to be mastered by OS designers

Copyright © 2002–2018 UMaine Computer Science Department – 9 / 18

Kernel

- The heart of the OS is the *kernel*
- Kernel does:
 - scheduling
 - resource management
- Other things often considered parts of the OS:
 - shells
 - libraries
 - utilities
- Kernel vs user mode

Copyright © 2002–2018 UMaine Computer Science Department – 10 / 18

Processes

- Process* \equiv program + its state of execution
- Almost all modern OS: multiprogramming, pseudo-parallelism
- System calls:
 - interface between processes and kernel
 - system calls + instructions = program's (process') API for computer

Copyright © 2002–2018 UMaine Computer Science Department – 11 / 18

Process state

- Why save a process' state?
 - So it can be interrupted, later resumed
- What is needed in the state?
 - program counter, stack pointer, CPU flags
 - where it is in memory, other memory-related information
 - bookkeeping information
- State of processes stored in *process table*
- Process control block (PCB)*

Copyright © 2002–2018 UMaine Computer Science Department – 12 / 18

Process synchronization

- Processes often share memory, or access to devices
- Race conditions
- Critical regions
- Need synchronization primitives, mechanisms

Copyright © 2002–2018 UMaine Computer Science Department – 13 / 18

Process scheduling

- If only 1 CPU, only one process can run at once
- When should a process give up the CPU?
- Scheduler* decides which one runs
- Simple scheduler: round-robin

Copyright © 2002–2018 UMaine Computer Science Department – 14 / 18

Virtual Memory

15 / 18

Virtual memory

- Size of memory limits number of processes that can run at once
- Also limits size of any given process
- Computers almost always can address more memory than they have
 - 32-bit computer – how much memory could it address?
 - 64-bit computer?
- Wouldn't it be nice to allow full *address space* to be used?
- Also nice if each process could think it had its own address space, starting at address 0

Copyright © 2002–2018 UMaine Computer Science Department – 15 / 18

Virtual memory

- Virtual memory gives each process the illusion of having its own, complete address space
- VM allows many more processes to run than will fit in memory
- Trick: use disk space instead of RAM (*swap space*, paging area)
- When a part of a process is not in use, *page it out* to disk
- When needed, page it in

Copyright © 2002–2018 UMaine Computer Science Department – 16 / 18

File Systems

17 / 18

File systems

- Another resource: disk space
- Files, directories (folders)
- Storing files
- Managing files
- Recovering from errors

Copyright © 2002–2018 UMaine Computer Science Department – 17 / 18

Other stuff

- Much more to OS than has been said
- Memory management, swap space management
- Input/output
- Network ↔ OS issues
- Handling deadlocks
- Computer security
- Distributed operating systems